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Robust Header Compression (ROHC):  
Framework and four profiles: RTP, UDP, ESP, and uncompressed

Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

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Abstract

This document specifies a highly robust and efficient header compression scheme for RTP/UDP/IP (Real-Time Transport Protocol, User Datagram Protocol, Internet Protocol), UDP/IP, and ESP/IP (Encapsulating Security Payload) headers.

Existing header compression schemes do not work well when used over links with significant error rates and long round-trip times. For many bandwidth limited links where header compression is essential, such characteristics are common.

This is done in a framework designed to be extensible. For example, a scheme for compressing TCP/IP headers will be simple to add, and is in development. Headers specific to Mobile IPv4 are not subject to special treatment, but are expected to be compressed sufficiently well by the provided methods for compression of sequences of extension headers and tunneling headers. For the most part, the same will apply to work in progress on Mobile IPv6, but future work might be required to handle some extension headers, when a standards track Mobile IPv6 has been completed.

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1. Introduction

During the last five years, two communication technologies in particular have become commonly used by the general public: cellular telephony and the Internet. Cellular telephony has provided its users with the revolutionary possibility of always being reachable with reasonable service quality no matter where they are. The main service provided by the dedicated terminals has been speech. The Internet, on the other hand, has from the beginning been designed for multiple services and its flexibility for all kinds of usage has been one of its strengths. Internet terminals have usually been general-purpose and have been attached over fixed connections. The experienced quality of some services (such as Internet telephony) has sometimes been low.

Today, IP telephony is gaining momentum thanks to improved technical solutions. It seems reasonable to believe that in the years to come, IP will become a commonly used way to carry telephony. Some future cellular telephony links might also be based on IP and IP telephony. Cellular phones may have become more general-purpose, and may have IP stacks supporting not only audio and video, but also web browsing, email, gaming, etc.

One of the scenarios we are envisioning might then be the one in Figure 1.1, where two mobile terminals are communicating with each other. Both are connected to base stations over cellular links, and the base stations are connected to each other through a wired (or possibly wireless) network. Instead of two mobile terminals, there could of course be one mobile and one wired terminal, but the case with two cellular links is technically more demanding.

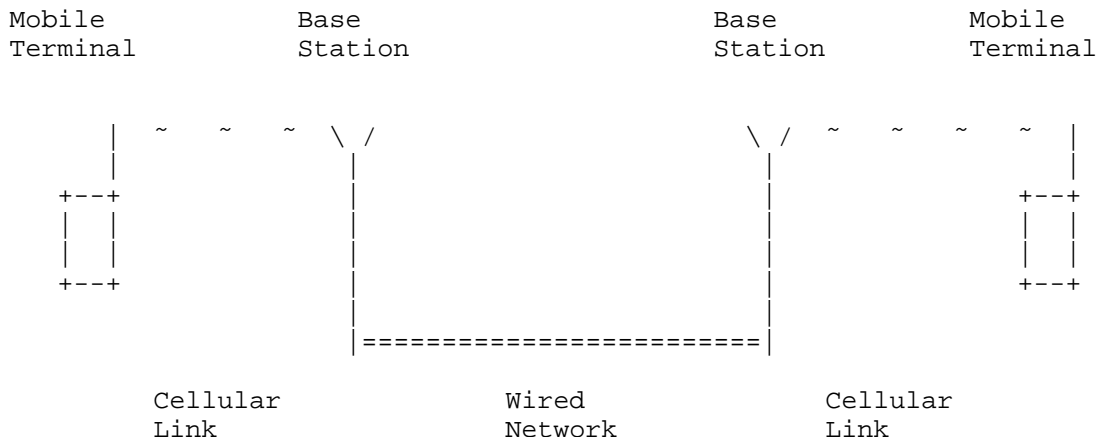


Figure 1.1 : Scenario for IP telephony over cellular links

It is obvious that the wired network can be IP-based. With the cellular links, the situation is less clear. IP could be terminated in the fixed network, and special solutions implemented for each supported service over the cellular link. However, this would limit the flexibility of the services supported. If technically and economically feasible, a solution with pure IP all the way from terminal to terminal would have certain advantages. However, to make this a viable alternative, a number of problems have to be addressed, in particular problems regarding bandwidth efficiency.

For cellular phone systems, it is of vital importance to use the scarce radio resources in an efficient way. A sufficient number of users per cell is crucial, otherwise deployment costs will be prohibitive. The quality of the voice service should also be as good as in today's cellular systems. It is likely that even with support for new services, lower quality of the voice service is acceptable only if costs are significantly reduced.

A problem with IP over cellular links when used for interactive voice conversations is the large header overhead. Speech data for IP telephony will most likely be carried by RTP [RTP]. A packet will then, in addition to link layer framing, have an IP [IPv4] header (20 octets), a UDP [UDP] header (8 octets), and an RTP header (12 octets) for a total of 40 octets. With IPv6 [IPv6], the IP header is 40 octets for a total of 60 octets. The size of the payload depends on the speech coding and frame sizes being used and may be as low as 15-20 octets.

From these numbers, the need for reducing header sizes for efficiency reasons is obvious. However, cellular links have characteristics that make header compression as defined in [IPHC,CRTP] perform less than well. The most important characteristic is the lossy behavior of cellular links, where a bit error rate (BER) as high as  $1e-3$  must be accepted to keep the radio resources efficiently utilized. In severe operating situations, the BER can be as high as  $1e-2$ . The other problematic characteristic is the long round-trip time (RTT) of the cellular link, which can be as high as 100-200 milliseconds. An additional problem is that the residual BER is nontrivial, i.e., lower layers can sometimes deliver frames containing undetected errors. A viable header compression scheme for cellular links must be able to handle loss on the link between the compression and decompression point as well as loss before the compression point.

Bandwidth is the most costly resource in cellular links. Processing power is very cheap in comparison. Implementation or computational simplicity of a header compression scheme is therefore of less importance than its compression ratio and robustness.

## 2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119.

### BER

Bit Error Rate. Cellular radio links can have a fairly high BER. In this document BER is usually given as a probability, but one also needs to consider the error distribution as bit errors are not independent.



#### Cellular links

Wireless links between mobile terminals and base stations.

#### Compression efficiency

The performance of a header compression scheme can be described with three parameters: compression efficiency, robustness and compression transparency. The compression efficiency is determined by how much the header sizes are reduced by the compression scheme.

#### Compression transparency

The performance of a header compression scheme can be described with three parameters: compression efficiency, robustness, and compression transparency. The compression transparency is a measure of the extent to which the scheme ensures that the decompressed headers are semantically identical to the original headers. If all decompressed headers are semantically identical to the corresponding original headers, the transparency is 100 percent. Compression transparency is high when damage propagation is low.

#### Context

The context of the compressor is the state it uses to compress a header. The context of the decompressor is the state it uses to decompress a header. Either of these or the two in combination are usually referred to as "context", when it is clear which is intended. The context contains relevant information from previous headers in the packet stream, such as static fields and possible reference values for compression and decompression. Moreover, additional information describing the packet stream is also part of the context, for example information about how the IP Identifier field changes and the typical inter-packet increase in sequence numbers or timestamps.

#### Context damage

When the context of the decompressor is not consistent with the context of the compressor, decompression may fail to reproduce the original header. This situation can occur when the context of the decompressor has not been initialized properly or when packets have been lost or damaged between compressor and decompressor.

Packets which cannot be decompressed due to inconsistent contexts are said to be lost due to context damage. Packets that are decompressed but contain errors due to inconsistent contexts are said to be damaged due to context damage.

#### Context repair mechanism

Context repair mechanisms are mechanisms that bring the contexts in sync when they were not. This is needed to avoid excessive loss due to context damage. Examples are the context request mechanism of CRTP, the NACK mechanisms of O- and R-mode, and the periodic refreshes of U-mode.

Note that there are also mechanisms that prevent (some) context inconsistencies from occurring, for example the ACK-based updates of the context in R-mode, the repetitions after change in U- and O-mode, and the CRCs which protect context updating information.

#### CRC-DYNAMIC

Opposite of CRC-STATIC.

#### CRC-STATIC

A CRC over the original header is the primary mechanism used by ROHC to detect incorrect decompression. In order to decrease computational complexity, the fields of the header are conceptually rearranged when the CRC is computed, so that it is first computed over octets which are static (called CRC-STATIC in this document) and then over octets whose values are expected to change between packets (CRC-DYNAMIC). In this manner, the intermediate result of the CRC computation, after it has covered the CRC-STATIC fields, can be reused for several packets. The restarted CRC computation only covers the CRC-DYNAMIC octets. See section 5.9.

#### Damage propagation

Delivery of incorrect decompressed headers, due to errors in (i.e., loss of or damage to) previous header(s) or feedback.

#### Loss propagation

Loss of headers, due to errors in (i.e., loss of or damage to) previous header(s) or feedback.

#### Error detection

Detection of errors. If error detection is not perfect, there will be residual errors.

#### Error propagation

Damage propagation or loss propagation.

#### Header compression profile

A header compression profile is a specification of how to compress the headers of a certain kind of packet stream over a certain kind of link. Compression profiles provide the details of the header compression framework introduced in this document. The profile concept makes use of profile identifiers to separate different profiles which are used when setting up the compression scheme. All variations and parameters of the header compression scheme that are not part of the context state are handled by different profile identifiers.

#### Packet

Generally, a unit of transmission and reception (protocol data unit). Specifically, when contrasted with "frame", the packet compressed and then decompressed by ROHC. Also called "uncompressed packet".

#### Packet Stream

A sequence of packets where the field values and change patterns of field values are such that the headers can be compressed using the same context.

#### Pre-HC links

The Pre-HC links are all links that a packet has traversed before the header compression point. If we consider a path with cellular links as first and last hops, the Pre-HC links for the compressor at the last link are the first cellular link plus the wired links in between.

#### Residual error

Error introduced during transmission and not detected by lower-layer error detection schemes.

### Robustness

The performance of a header compression scheme can be described with three parameters: compression efficiency, robustness, and compression transparency. A robust scheme tolerates loss and residual errors on the link over which header compression takes place without losing additional packets or introducing additional errors in decompressed headers.

### RTT

The RTT (round-trip time) is the time elapsing from the moment the compressor sends a packet until it receives feedback related to that packet (when such feedback is sent).

### Spectrum efficiency

Radio resources are limited and expensive. Therefore they must be used efficiently to make the system economically feasible. In cellular systems this is achieved by maximizing the number of users served within each cell, while the quality of the provided services is kept at an acceptable level. A consequence of efficient spectrum use is a high rate of errors (frame loss and residual bit errors), even after channel coding with error correction.

### String

A sequence of headers in which the values of all fields being compressed change according to a pattern which is fixed with respect to a sequence number. Each header in a string can be compressed by representing it with a ROHC header which essentially only carries an encoded sequence number. Fields not being compressed (e.g., random IP-ID, UDP Checksum) are irrelevant to this definition.

### Timestamp stride

The timestamp stride (TS\_STRIDE) is the expected increase in the timestamp value between two RTP packets with consecutive sequence numbers.

## 2.1. Acronyms

This section lists most acronyms used for reference.

AH Authentication Header.  
CID Context Identifier.  
CRC Cyclic Redundancy Check. Error detection mechanism.  
CRTP Compressed RTP. RFC 2508.  
CTCP Compressed TCP. Also called VJ header compression. RFC 1144.  
ESP Encapsulating Security Payload.  
FC Full Context state (decompressor).  
FO First Order state (compressor).  
GRE Generic Routing Encapsulation. RFC 2784, RFC 2890.  
HC Header Compression.  
IPHC IP Header Compression. RFC 2507.  
IPX Flag in Extension 2.  
IR Initiation and Refresh state (compressor). Also IR packet.  
IR-DYN IR-DYN packet.  
LSB Least Significant Bits.  
MRRU Maximum Reconstructed Reception Unit.  
MTU Maximum Transmission Unit.  
MSB Most Significant Bits.  
NBO Flag indicating whether the IP-ID is in Network Byte Order.  
NC No Context state (decompressor).  
O-mode Bidirectional Optimistic mode.  
PPP Point-to-Point Protocol.  
R-mode Bidirectional Reliable mode.  
RND Flag indicating whether the IP-ID behaves randomly.  
ROHC RObust Header Compression.  
RTCP Real-Time Control Protocol. See RTP.  
RTP Real-Time Protocol. RFC 1889.  
RTT Round Trip Time (see section 2).  
SC Static Context state (decompressor).  
SN (compressed) Sequence Number. Usually RTP Sequence Number.  
SO Second Order state (compressor).  
SPI Security Parameters Index.  
SSRC Sending source. Field in RTP header.  
CSRC Contributing source. Optional list of CSRCs in RTP header.  
TC Traffic Class. Octet in IPv6 header. See also TOS.  
TOS Type Of Service. Octet in IPv4 header. See also TC.  
TS (compressed) RTP Timestamp.  
U-mode Unidirectional mode.  
W-LSB Window based LSB encoding. See section 4.5.2.

### 3. Background

This chapter provides a background to the subject of header compression. The fundamental ideas are described together with existing header compression schemes. Their drawbacks and requirements are then discussed, providing motivation for new header compression solutions.

#### 3.1. Header compression fundamentals

The main reason why header compression can be done at all is the fact that there is significant redundancy between header fields, both within the same packet header but in particular between consecutive packets belonging to the same packet stream. By sending static field information only initially and utilizing dependencies and predictability for other fields, the header size can be significantly reduced for most packets.

Relevant information from past packets is maintained in a context. The context information is used to compress (decompress) subsequent packets. The compressor and decompressor update their contexts upon certain events. Impairment events may lead to inconsistencies between the contexts of the compressor and decompressor, which in turn may cause incorrect decompression. A robust header compression scheme needs mechanisms for avoiding context inconsistencies and also needs mechanisms for making the contexts consistent when they were not.

#### 3.2. Existing header compression schemes

The original header compression scheme, CTCP [VJHC], was invented by Van Jacobson. CTCP compresses the 40 octet IP+TCP header to 4 octets. The CTCP compressor detects transport-level retransmissions and sends a header that updates the context completely when they occur. This repair mechanism does not require any explicit signaling between compressor and decompressor.

A general IP header compression scheme, IP header compression [IPHC], improves somewhat on CTCP and can compress arbitrary IP, TCP, and UDP headers. When compressing non-TCP headers, IPHC does not use delta encoding and is robust. When compressing TCP, the repair mechanism of CTCP is augmented with a link-level nacking scheme which speeds up the repair. IPHC does not compress RTP headers.

CRTP [CRTP, IPHC] by Casner and Jacobson is a header compression scheme that compresses 40 octets of IPv4/UDP/RTP headers to a minimum of 2 octets when the UDP Checksum is not enabled. If the UDP Checksum is enabled, the minimum CRTP header is 4 octets. CRTP

cannot use the same repair mechanism as CTCP since UDP/RTP does not retransmit. Instead, CRTP uses explicit signaling messages from decompressor to compressor, called CONTEXT\_STATE messages, to indicate that the context is out of sync. The link round-trip time will thus limit the speed of this context repair mechanism.

On lossy links with long round-trip times, such as most cellular links, CRTP does not perform well. Each lost packet over the link causes several subsequent packets to be lost since the context is out of sync during at least one link round-trip time. This behavior is documented in [CRTPC]. For voice conversations such long loss events will degrade the voice quality. Moreover, bandwidth is wasted by the large headers sent by CRTP when updating the context. [CRTPC] found that CRTP did not perform well enough for a lossy cellular link. It is clear that CRTP alone is not a viable header compression scheme for IP telephony over cellular links.

To avoid losing packets due to the context being out of sync, CRTP decompressors can attempt to repair the context locally by using a mechanism known as TWICE. Each CRTP packet contains a counter which is incremented by one for each packet sent out by the CRTP compressor. If the counter increases by more than one, at least one packet was lost over the link. The decompressor then attempts to repair the context by guessing how the lost packet(s) would have updated it. The guess is then verified by decompressing the packet and checking the UDP Checksum -- if it succeeds, the repair is deemed successful and the packet can be forwarded or delivered. TWICE derives its name from the observation that when the compressed packet stream is regular, the correct guess is to apply the update in the current packet twice. [CRTPC] found that even with TWICE, CRTP doubled the number of lost packets. TWICE improves CRTP performance significantly. However, there are several problems with using TWICE:

1) It becomes mandatory to use the UDP Checksum:

- the minimal compressed header size increases by 100% to 4 octets.
- most speech codecs developed for cellular links tolerate errors in the encoded data. Such codecs will not want to enable the UDP Checksum, since they do want damaged packets to be delivered.
- errors in the payload will make the UDP Checksum fail when the guess is correct (and might make it succeed when the guess is wrong).

- 2) Loss in an RTP stream that occurs before the compression point will make updates in CRTP headers less regular. Simple-minded versions of TWICE will then perform badly. More sophisticated versions would need more repair attempts to succeed.

### 3.3. Requirements on a new header compression scheme

The major problem with CRTP is that it is not sufficiently robust against packets being damaged between compressor and decompressor. A viable header compression scheme must be less fragile. This increased robustness must be obtained without increasing the compressed header size; a larger header would make IP telephony over cellular links economically unattractive.

A major cause of the bad performance of CRTP over cellular links is the long link round-trip time, during which many packets are lost when the context is out of sync. This problem can be attacked directly by finding ways to reduce the link round-trip time. Future generations of cellular technologies may indeed achieve lower link round-trip times. However, these will probably always be fairly high. The benefits in terms of lower loss and smaller bandwidth demands if the context can be repaired locally will be present even if the link round-trip time is decreased. A reliable way to detect a successful context repair is then needed.

One might argue that a better way to solve the problem is to improve the cellular link so that packet loss is less likely to occur. Such modifications do not appear to come for free, however. If links were made (almost) error free, the system might not be able to support a sufficiently large number of users per cell and might thus be economically infeasible.

One might also argue that the speech codecs should be able to deal with the kind of packet loss induced by CRTP, in particular since the speech codecs probably must be able to deal with packet loss anyway if the RTP stream crosses the Internet. While the latter is true, the kind of loss induced by CRTP is difficult to deal with. It is usually not possible to completely hide a loss event where well over 100 ms worth of sound is completely lost. If such loss occurs frequently at both ends of the end-to-end path, the speech quality will suffer.

A detailed description of the requirements specified for ROHC may be found in [REQ].



### 3.4. Classification of header fields

As mentioned earlier, header compression is possible due to the fact that there is much redundancy between header field values within packets, but especially between consecutive packets. To utilize these properties for header compression, it is important to understand the change patterns of the various header fields.

All header fields have been classified in detail in appendix A. The fields are first classified at a high level and then some of them are studied more in detail. Finally, the appendix concludes with recommendations on how the various fields should be handled by header compression algorithms. The main conclusion that can be drawn is that most of the header fields can easily be compressed away since they never or seldom change. Only 5 fields, with a combined size of about 10 octets, need more sophisticated mechanisms. These fields are:

- IPv4 Identification (16 bits)      - IP-ID
- UDP Checksum (16 bits)
- RTP Marker (1 bit)                      - M-bit
- RTP Sequence Number (16 bits)      - SN
- RTP Timestamp (32 bits)              - TS

The analysis in Appendix A reveals that the values of the TS and IP-ID fields can usually be predicted from the RTP Sequence Number, which increments by one for each packet emitted by an RTP source. The M-bit is also usually the same, but needs to be communicated explicitly occasionally. The UDP Checksum should not be predicted and is sent as-is when enabled.

The way ROHC RTP compression operates, then, is to first establish functions from SN to the other fields, and then reliably communicate the SN. Whenever a function from SN to another field changes, i.e., the existing function gives a result which is different from the field in the header to be compressed, additional information is sent to update the parameters of that function.

Headers specific to Mobile IP (for IPv4 or IPv6) do not receive any special treatment in this document. They are compressible, however, and it is expected that the compression efficiency for Mobile IP headers will be good enough due to the handling of extension header lists and tunneling headers. It would be relatively painless to introduce a new ROHC profile with special treatment for Mobile IPv6 specific headers should the completed work on the Mobile IPv6 protocols (work in progress in the IETF) make that necessary.

## 4. Header compression framework

### 4.1. Operating assumptions

Cellular links, which are a primary target for ROHC, have a number of characteristics that are described briefly here. ROHC requires functionality from lower layers that is outlined here and more thoroughly described in the lower layer guidelines document [LLG].

#### Channels

ROHC header-compressed packets flow on channels. Unlike many fixed links, some cellular radio links can have several channels connecting the same pair of nodes. Each channel can have different characteristics in terms of error rate, bandwidth, etc.

#### Context identifiers

On some channels, the ability to transport multiple packet streams is required. It can also be feasible to have channels dedicated to individual packet streams. Therefore, ROHC uses a distinct context identifier space per channel and can eliminate context identifiers completely for one of the streams when few streams share a channel.

#### Packet type indication

Packet type indication is done in the header compression scheme itself. Unless the link already has a way of indicating packet types which can be used, such as PPP, this provides smaller compressed headers overall. It may also be less difficult to allocate a single packet type, rather than many, in order to run ROHC over links such as PPP.

#### Reordering

The channel between compressor and decompressor is required to maintain packet ordering, i.e., the decompressor must receive packets in the same order as the compressor sent them. (Reordering before the compression point, however, is dealt with, i.e., there is no assumption that the compressor will only receive packets in sequence.)

#### Duplication

The channel between compressor and decompressor is required to not duplicate packets. (Duplication before the compression point, however, is dealt with, i.e., there is no assumption that the compressor will receive only one copy of each packet.)

#### Packet length

ROHC is designed under the assumption that lower layers indicate the length of a compressed packet. ROHC packets do not contain length information for the payload.

#### Framing

The link layer must provide framing that makes it possible to distinguish frame boundaries and individual frames.

#### Error detection/protection

The ROHC scheme has been designed to cope with residual errors in the headers delivered to the decompressor. CRCs and sanity checks are used to prevent or reduce damage propagation. However, it is RECOMMENDED that lower layers deploy error detection for ROHC headers and do not deliver ROHC headers with high residual error rates.

Without giving a hard limit on the residual error rate acceptable to ROHC, it is noted that for a residual bit error rate of at most  $1E-5$ , the ROHC scheme has been designed not to increase the number of damaged headers, i.e., the number of damaged headers due to damage propagation is designed to be less than the number of damaged headers caught by the ROHC error detection scheme.

#### Negotiation

In addition to the packet handling mechanisms above, the link layer MUST provide a way to negotiate header compression parameters, see also section 5.1.1. (For unidirectional links, this negotiation may be performed out-of-band or even a priori.)

### 4.2. Dynamicity

The ROHC protocol achieves its compression gain by establishing state information at both ends of the link, i.e., at the compressor and at the decompressor. Different parts of the state are established at different times and with different frequency; hence, it can be said that some of the state information is more dynamic than the rest.

Some state information is established at the time a channel is established; ROHC assumes the existence of an out-of-band negotiation protocol (such as PPP), or predefined channel state (most useful for unidirectional links). In both cases, we speak of "negotiated channel state". ROHC does not assume that this state can change dynamically during the channel lifetime (and does not explicitly support such changes, although some changes may be innocuous from a protocol point of view). An example of negotiated channel state is the highest context ID number to be used by the compressor (MAX\_CID).

Other state information is associated with the individual packet streams in the channel; this state is said to be part of the context. Using context identifiers (CIDs), multiple packet streams with different contexts can share a channel. The negotiated channel state indicates the highest context identifier to be used, as well as the selection of one of two ways to indicate the CID in the compressed header.

It is up to the compressor to decide which packets to associate with a context (or, equivalently, which packets constitute a single stream); however, ROHC is efficient only when all packets of a stream share certain properties, such as having the same values for fields that are described as "static" in this document (e.g., the IP addresses, port numbers, and RTP parameters such as the payload type). The efficiency of ROHC RTP also depends on the compressor seeing most RTP Sequence Numbers.

Streams need not share all characteristics important for compression. ROHC has a notion of compression profiles: a compression profile denotes a predefined set of such characteristics. To provide extensibility, the negotiated channel state includes the set of profiles acceptable to the decompressor. The context state includes the profile currently in use for the context.

Other elements of the context state may include the current values of all header fields (from these one can deduce whether an IPv4 header is present in the header chain, and whether UDP Checksums are enabled), as well as additional compression context that is not part of an uncompressed header, e.g., TS\_STRIDE, IP-ID characteristics (incrementing as a 16-bit value in network byte order? random?), a number of old reference headers, and the compressor/decompressor state machines (see next section).

This document actually defines four ROHC profiles: One uncompressed profile, the main ROHC RTP compression profile, and two variants of this profile for compression of packets with header chains that end

in UDP and ESP, respectively, but where RTP compression is not applicable. The descriptive text in the rest of this section is referring to the main ROHC RTP compression profile.

4.3. Compression and decompression states

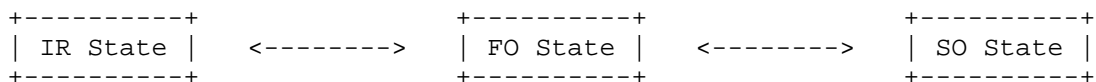
Header compression with ROHC can be characterized as an interaction between two state machines, one compressor machine and one decompressor machine, each instantiated once per context. The compressor and the decompressor have three states each, which in many ways are related to each other even if the meaning of the states are slightly different for the two parties. Both machines start in the lowest compression state and transit gradually to higher states.

Transitions need not be synchronized between the two machines. In normal operation it is only the compressor that temporarily transits back to lower states. The decompressor will transit back only when context damage is detected.

Subsequent sections present an overview of the state machines and their corresponding states, respectively, starting with the compressor.

4.3.1. Compressor states

For ROHC compression, the three compressor states are the Initialization and Refresh (IR), First Order (FO), and Second Order (SO) states. The compressor starts in the lowest compression state (IR) and transits gradually to higher compression states. The compressor will always operate in the highest possible compression state, under the constraint that the compressor is sufficiently confident that the decompressor has the information necessary to decompress a header compressed according to that state.



Decisions about transitions between the various compression states are taken by the compressor on the basis of:

- variations in packet headers
- positive feedback from decompressor (Acknowledgments -- ACKs)
- negative feedback from decompressor (Negative ACKs -- NACKs)
- periodic timeouts (when operating in unidirectional mode, i.e., over simplex channels or when feedback is not enabled)

How transitions are performed is explained in detail in chapter 5 for each mode of operation.

#### 4.3.1.1. Initialization and Refresh (IR) State

The purpose of the IR state is to initialize the static parts of the context at the decompressor or to recover after failure. In this state, the compressor sends complete header information. This includes all static and nonstatic fields in uncompressed form plus some additional information.

The compressor stays in the IR state until it is fairly confident that the decompressor has received the static information correctly.

#### 4.3.1.2. First Order (FO) State

The purpose of the FO state is to efficiently communicate irregularities in the packet stream. When operating in this state, the compressor rarely sends information about all dynamic fields, and the information sent is usually compressed at least partially. Only a few static fields can be updated. The difference between IR and FO should therefore be clear.

The compressor enters this state from the IR state, and from the SO state whenever the headers of the packet stream do not conform to their previous pattern. It stays in the FO state until it is confident that the decompressor has acquired all the parameters of the new pattern. Changes in fields that are always irregular are communicated in all packets and are therefore part of what is a uniform pattern.

Some or all packets sent in the FO state carry context updating information. It is very important to detect corruption of such packets to avoid erroneous updates and context inconsistencies.

#### 4.3.1.3. Second Order (SO) State

This is the state where compression is optimal. The compressor enters the SO state when the header to be compressed is completely predictable given the SN (RTP Sequence Number) and the compressor is sufficiently confident that the decompressor has acquired all parameters of the functions from SN to other fields. Correct decompression of packets sent in the SO state only hinges on correct decompression of the SN. However, successful decompression also requires that the information sent in the preceding FO state packets has been successfully received by the decompressor.

The compressor leaves this state and goes back to the FO state when the header no longer conforms to the uniform pattern and cannot be independently compressed on the basis of previous context information.

4.3.2. Decompressor states

The decompressor starts in its lowest compression state, "No Context" and gradually transits to higher states. The decompressor state machine normally never leaves the "Full Context" state once it has entered this state.



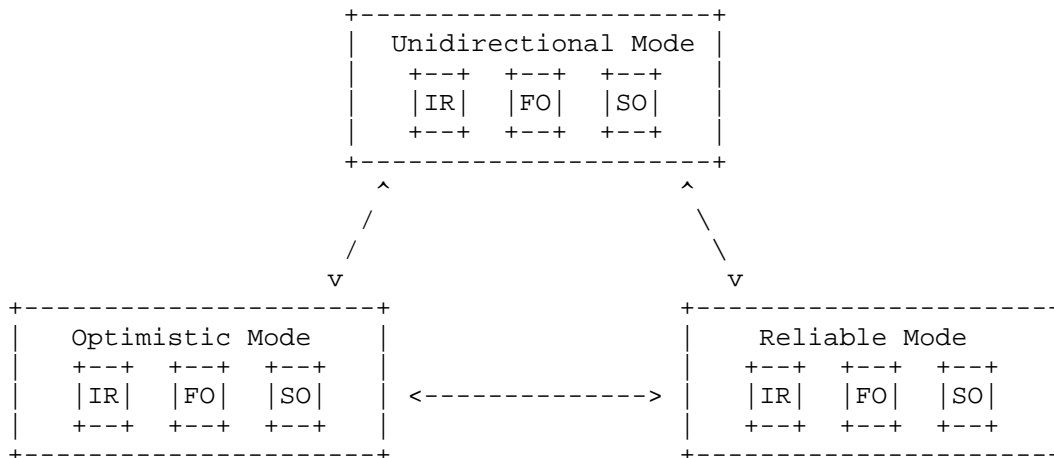
Initially, while working in the "No Context" state, the decompressor has not yet successfully decompressed a packet. Once a packet has been decompressed correctly (for example, upon reception of an initialization packet with static and dynamic information), the decompressor can transit all the way to the "Full Context" state, and only upon repeated failures will it transit back to lower states. However, when that happens it first transits back to the "Static Context" state. There, reception of any packet sent in the FO state is normally sufficient to enable transition to the "Full Context" state again. Only when decompression of several packets sent in the FO state fails in the "Static Context" state will the decompressor go all the way back to the "No Context" state.

When state transitions are performed is explained in detail in chapter 5.

4.4. Modes of operation

The ROHC scheme has three modes of operation, called Unidirectional, Bidirectional Optimistic, and Bidirectional Reliable mode.

It is important to understand the difference between states, as described in the previous chapter, and modes. These abstractions are orthogonal to each other. The state abstraction is the same for all modes of operation, while the mode controls the logic of state transitions and what actions to perform in each state.



The optimal mode to operate in depends on the characteristics of the environment of the compression protocol, such as feedback abilities, error probabilities and distributions, effects of header size variation, etc. All ROHC implementations MUST implement and support all three modes of operation. The three modes are briefly described in the following subsections.

Detailed descriptions of the three modes of operation regarding compression and decompression logic are given in chapter 5. The mode transition mechanisms, too, are described in chapter 5.

#### 4.4.1. Unidirectional mode -- U-mode

When in the Unidirectional mode of operation, packets are sent in one direction only: from compressor to decompressor. This mode therefore makes ROHC usable over links where a return path from decompressor to compressor is unavailable or undesirable.

In U-mode, transitions between compressor states are performed only on account of periodic timeouts and irregularities in the header field change patterns in the compressed packet stream. Due to the periodic refreshes and the lack of feedback for initiation of error recovery, compression in the Unidirectional mode will be less efficient and have a slightly higher probability of loss propagation compared to any of the Bidirectional modes.

Compression with ROHC MUST start in the Unidirectional mode. Transition to any of the Bidirectional modes can be performed as soon as a packet has reached the decompressor and it has replied with a feedback packet indicating that a mode transition is desired (see chapter 5).



#### 4.4.2. Bidirectional Optimistic mode -- O-mode

The Bidirectional Optimistic mode is similar to the Unidirectional mode. The difference is that a feedback channel is used to send error recovery requests and (optionally) acknowledgments of significant context updates from decompressor to compressor (not, however, for pure sequence number updates). Periodic refreshes are not used in the Bidirectional Optimistic mode.

O-mode aims to maximize compression efficiency and sparse usage of the feedback channel. It reduces the number of damaged headers delivered to the upper layers due to residual errors or context invalidation. The frequency of context invalidation may be higher than for R-mode, in particular when long loss/error bursts occur. Refer to section 4.7 for more details.

#### 4.4.3. Bidirectional Reliable mode -- R-mode

The Bidirectional Reliable mode differs in many ways from the previous two. The most important differences are a more intensive usage of the feedback channel and a stricter logic at both the compressor and the decompressor that prevents loss of context synchronization between compressor and decompressor except for very high residual bit error rates. Feedback is sent to acknowledge all context updates, including updates of the sequence number field. However, not every packet updates the context in Reliable mode.

R-mode aims to maximize robustness against loss propagation and damage propagation, i.e., minimize the probability of context invalidation, even under header loss/error burst conditions. It may have a lower probability of context invalidation than O-mode, but a larger number of damaged headers may be delivered when the context actually is invalidated. Refer to section 4.7 for more details.

### 4.5. Encoding methods

This chapter describes the encoding methods used for header fields. How the methods are applied to each field (e.g., values of associated parameters) is specified in section 5.7.

#### 4.5.1. Least Significant Bits (LSB) encoding

Least Significant Bits (LSB) encoding is used for header fields whose values are usually subject to small changes. With LSB encoding, the  $k$  least significant bits of the field value are transmitted instead of the original field value, where  $k$  is a positive integer. After receiving  $k$  bits, the decompressor derives the original value using a previously received value as reference ( $v_{ref}$ ).

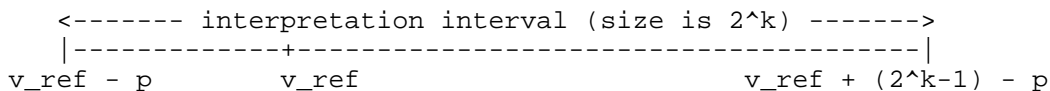
The scheme is guaranteed to be correct if the compressor and the decompressor each use interpretation intervals

- 1) in which the original value resides, and
- 2) in which the original value is the only value that has the exact same  $k$  least significant bits as those transmitted.

The interpretation interval can be described as a function  $f(v\_ref, k)$ . Let

$$f(v\_ref, k) = [v\_ref - p, v\_ref + (2^k - 1) - p]$$

where  $p$  is an integer.



The function  $f$  has the following property: for any value  $k$ , the  $k$  least significant bits will uniquely identify a value in  $f(v\_ref, k)$ .

The parameter  $p$  is introduced so that the interpretation interval can be shifted with respect to  $v\_ref$ . Choosing a good value for  $p$  will yield a more efficient encoding for fields with certain characteristics. Below are some examples:

- a) For field values that are expected always to increase,  $p$  can be set to  $-1$ . The interpretation interval becomes  $[v\_ref + 1, v\_ref + 2^k]$ .
- b) For field values that stay the same or increase,  $p$  can be set to  $0$ . The interpretation interval becomes  $[v\_ref, v\_ref + 2^k - 1]$ .
- c) For field values that are expected to deviate only slightly from a constant value,  $p$  can be set to  $2^{(k-1)} - 1$ . The interpretation interval becomes  $[v\_ref - 2^{(k-1)} + 1, v\_ref + 2^{(k-1)}]$ .
- d) For field values that are expected to undergo small negative changes and larger positive changes, such as the RTP TS for video, or RTP SN when there is misordering,  $p$  can be set to  $2^{(k-2)} - 1$ . The interval becomes  $[v\_ref - 2^{(k-2)} + 1, v\_ref + 3 * 2^{(k-2)}]$ , i.e.,  $3/4$  of the interval is used for positive changes.

The following is a simplified procedure for LSB compression and decompression; it is modified for robustness and damage propagation protection in the next subsection:

- 1) The compressor (decompressor) always uses  $v\_ref\_c$  ( $v\_ref\_d$ ), the last value that has been compressed (decompressed), as  $v\_ref$ ;
- 2) When compressing a value  $v$ , the compressor finds the minimum value of  $k$  such that  $v$  falls into the interval  $f(v\_ref\_c, k)$ . Call this function  $k = g(v\_ref\_c, v)$ . When only a few distinct values of  $k$  are possible, for example due to limitations imposed by packet formats (see section 5.7), the compressor will instead pick the smallest  $k$  that puts  $v$  in the interval  $f(v\_ref\_c, k)$ .
- 3) When receiving  $m$  LSBs, the decompressor uses the interpretation interval  $f(v\_ref\_d, m)$ , called  $interval\_d$ . It picks as the decompressed value the one in  $interval\_d$  whose LSBs match the received  $m$  bits.

Note that the values to be encoded have a finite range; for example, the RTP SN ranges from 0 to 0xFFFF. When the SN value is close to 0 or 0xFFFF, the interpretation interval can straddle the wraparound boundary between 0 and 0xFFFF.

The scheme is complicated by two factors: packet loss between the compressor and decompressor, and transmission errors undetected by the lower layer. In the former case, the compressor and decompressor will lose the synchronization of  $v\_ref$ , and thus also of the interpretation interval. If  $v$  is still covered by the intersection( $interval\_c$ ,  $interval\_d$ ), the decompression will be correct. Otherwise, incorrect decompression will result. The next section will address this issue further.

In the case of undetected transmission errors, the corrupted LSBs will give an incorrectly decompressed value that will later be used as  $v\_ref\_d$ , which in turn is likely to lead to damage propagation. This problem is addressed by using a secure reference, i.e., a reference value whose correctness is verified by a protecting CRC. Consequently, the procedure 1) above is modified as follows:

- 1) a) the compressor always uses as  $v\_ref\_c$  the last value that has been compressed and sent with a protecting CRC.
- b) the decompressor always uses as  $v\_ref\_d$  the last correct value, as verified by a successful CRC.

Note that in U/O-mode, 1) b) is modified so that if decompression of the SN fails using the last verified SN reference, another decompression attempt is made using the last but one verified SN reference. This procedure mitigates damage propagation when a small CRC fails to detect a damaged value. See section 5.3.2.2.3 for further details.

#### 4.5.2. Window-based LSB encoding (W-LSB encoding)

This section describes how to modify the simplified algorithm in 4.5.1 to achieve robustness.

The compressor may not be able to determine the exact value of `v_ref_d` that will be used by the decompressor for a particular value `v`, since some candidates for `v_ref_d` may have been lost or damaged. However, by using feedback or by making reasonable assumptions, the compressor can limit the candidate set. The compressor then calculates `k` such that no matter which `v_ref_d` in the candidate set the decompressor uses, `v` is covered by the resulting `interval_d`.

Since the decompressor always uses as the reference the last received value where the CRC succeeded, the compressor maintains a sliding window containing the candidates for `v_ref_d`. The sliding window is initially empty. The following operations are performed on the sliding window by the compressor:

- 1) After sending a value `v` (compressed or uncompressed) protected by a CRC, the compressor adds `v` to the sliding window.
- 2) For each value `v` being compressed, the compressor chooses  $k = \max(g(v_{\min}, v), g(v_{\max}, v))$ , where `v_min` and `v_max` are the minimum and maximum values in the sliding window, and `g` is the function defined in the previous section.
- 3) When the compressor is sufficiently confident that a certain value `v` and all values older than `v` will not be used as reference by the decompressor, the window is advanced by removing those values (including `v`). The confidence may be obtained by various means. In R-mode, an ACK from the decompressor implies that values older than the ACKed one can be removed from the sliding window. In U/O-mode there is always a CRC to verify correct decompression, and a sliding window with a limited maximum width is used. The window width is an implementation dependent optimization parameter.

Note that the decompressor follows the procedure described in the previous section, except that in R-mode it MUST ACK each header received with a succeeding CRC (see also section 5.5).

#### 4.5.3. Scaled RTP Timestamp encoding

The RTP Timestamp (TS) will usually not increase by an arbitrary number from packet to packet. Instead, the increase is normally an integral multiple of some unit (`TS_STRIDE`). For example, in the case of audio, the sample rate is normally 8 kHz and one voice frame may

cover 20 ms. Furthermore, each voice frame is often carried in one RTP packet. In this case, the RTP increment is always  $n * 160$  ( $= 8000 * 0.02$ ), for some integer  $n$ . Note that silence periods have no impact on this, as the sample clock at the source normally keeps running without changing either frame rate or frame boundaries.

In the case of video, there is usually a `TS_STRIDE` as well when the video frame level is considered. The sample rate for most video codecs is 90 kHz. If the video frame rate is fixed, say, to 30 frames/second, the TS will increase by  $n * 3000$  ( $= n * 90000 / 30$ ) between video frames. Note that a video frame is often divided into several RTP packets to increase robustness against packet loss. In this case several RTP packets will carry the same TS.

When using scaled RTP Timestamp encoding, the TS is downscaled by a factor of `TS_STRIDE` before compression. This saves

$$\text{floor}(\log_2(\text{TS\_STRIDE}))$$

bits for each compressed TS. TS and `TS_SCALED` satisfy the following equality:

$$\text{TS} = \text{TS\_SCALED} * \text{TS\_STRIDE} + \text{TS\_OFFSET}$$

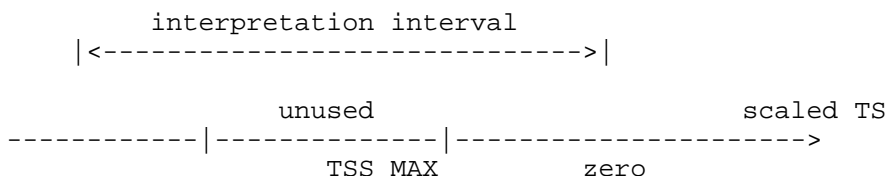
`TS_STRIDE` is explicitly, and `TS_OFFSET` implicitly, communicated to the decompressor. The following algorithm is used:

1. Initialization: The compressor sends to the decompressor the value of `TS_STRIDE` and the absolute value of one or several TS fields. The latter are used by the decompressor to initialize `TS_OFFSET` to (absolute value) modulo `TS_STRIDE`. Note that `TS_OFFSET` is the same regardless of which absolute value is used, as long as the unscaled TS value does not wrap around; see 4) below.
2. Compression: After initialization, the compressor no longer compresses the original TS values. Instead, it compresses the downscaled values:  $\text{TS\_SCALED} = \text{TS} / \text{TS\_STRIDE}$ . The compression method could be either W-LSB encoding or the timer-based encoding described in the next section.
3. Decompression: When receiving the compressed value of `TS_SCALED`, the decompressor first derives the value of the original `TS_SCALED`. The original RTP TS is then calculated as  $\text{TS} = \text{TS\_SCALED} * \text{TS\_STRIDE} + \text{TS\_OFFSET}$ .
4. Offset at wraparound: Wraparound of the unscaled 32-bit TS will invalidate the current value of `TS_OFFSET` used in the equation above. For example, let us assume  $\text{TS\_STRIDE} = 160 = 0xA0$  and the

current TS = 0xFFFFFFFF0. TS\_OFFSET is then 0x50 = 80. Then if the next RTP TS = 0x00000130 (i.e., the increment is 160 \* 2 = 320), the new TS\_OFFSET should be 0x00000130 modulo 0xA0 = 0x90 = 144. The compressor is not required to re-initialize TS\_OFFSET at wraparound. Instead, the decompressor MUST detect wraparound of the unscaled TS (which is trivial) and update TS\_OFFSET to

$$TS\_OFFSET = (\text{Wrapped around unscaled TS}) \text{ modulo } TS\_STRIDE$$

- 5. Interpretation interval at wraparound: Special rules are needed for the interpretation interval of the scaled TS at wraparound, since the maximum scaled TS, TSS\_MAX, (0xFFFFFFFF / TS\_STRIDE) may not have the form 2^m - 1. For example, when TS\_STRIDE is 160, the scaled TS is at most 26843545 which has LSBs 10011001. The wraparound boundary between the TSS\_MAX may thus not correspond to a natural boundary between LSBs.



When TSS\_MAX is part of the interpretation interval, a number of unused values are inserted into it after TSS\_MAX such that their LSBs follow naturally upon each other. For example, for TS\_STRIDE = 160 and k = 4, values corresponding to the LSBs 1010 through 1111 are inserted. The number of inserted values depends on k and the LSBs of the maximum scaled TS. The number of valid values in the interpretation interval should be high enough to maintain robustness. This can be ensured by the following rule:

Let a be the number of LSBs needed if there was no wraparound, and let b be the number of LSBs needed to disambiguate between TSS\_MAX and zero where the a LSBs of TSS\_MAX are set to zero. The number of LSB bits to send while TSS\_MAX or zero is part of the interpretation interval is b.

This scaling method can be applied to many frame-based codecs. However, the value of TS\_STRIDE might change during a session, for example as a result of adaptation strategies. If that happens, the unscaled TS is compressed until re-initialization of the new TS\_STRIDE and TS\_OFFSET is completed.

#### 4.5.4. Timer-based compression of RTP Timestamp

The RTP Timestamp [RFC 1889] is defined to identify the number of the first sample used to generate the payload. When 1) RTP packets carry payloads corresponding to a fixed sampling interval, 2) the sampling is done at a constant rate, and 3) packets are generated in lock-step with sampling, then the timestamp value will closely approximate a linear function of the time of day. This is the case for conversational media, such as interactive speech. The linear ratio is determined by the source sample rate. The linear pattern can be complicated by packetization (e.g., in the case of video where a video frame usually corresponds to several RTP packets) or frame rearrangement (e.g., B-frames are sent out-of-order by some video codecs).

With a fixed sample rate of 8 kHz, 20 ms in the time domain is equivalent to an increment of 160 in the unscaled TS domain, and to an increment of 1 in the scaled TS domain with `TS_STRIDE = 160`.

As a consequence, the (scaled) TS of headers arriving at the decompressor will be a linear function of time of day, with some deviation due to the delay jitter (and the clock inaccuracies) between the source and the decompressor. In normal operation, i.e., no crashes or failures, the delay jitter will be bounded to meet the requirements of conversational real-time traffic. Hence, by using a local clock the decompressor can obtain an approximation of the (scaled) TS in the header to be decompressed by considering its arrival time. The approximation can then be refined with the *k* LSBs of the (scaled) TS carried in the header. The value of *k* required to ensure correct decompression is a function of the jitter between the source and the decompressor.

If the compressor knows the potential jitter introduced between compressor and decompressor, it can determine *k* by using a local clock to estimate jitter in packet arrival times, or alternatively it can use a fixed *k* and discard packets arriving too much out of time.

The advantages of this scheme include:

- a) The size of the compressed TS is constant and small. In particular, it does NOT depend on the length of silence intervals. This is in contrast to other TS compression techniques, which at the beginning of a talkspurt require sending a number of bits dependent on the duration of the preceding silence interval.
- b) No synchronization is required between the clock local to the compressor and the clock local to the decompressor.

Note that although this scheme can be made to work using both scaled and unscaled TS, in practice it is always combined with scaled TS encoding because of the less demanding requirement on the clock resolution, e.g., 20 ms instead of 1/8 ms. Therefore, the algorithm described below assumes that the clock-based encoding scheme operates on the scaled TS. The case of unscaled TS would be similar, with changes to scale factors.

The major task of the compressor is to determine the value of  $k$ . Its sliding window now contains not only potential reference values for the TS but also their times of arrival at the compressor.

- 1) The compressor maintains a sliding window

$\{(T_j, a_j), \text{ for each header } j \text{ that can be used as a reference}\}$ ,

where  $T_j$  is the scaled TS for header  $j$ , and  $a_j$  is the arrival time of header  $j$ . The sliding window serves the same purpose as the W-LSB sliding window of section 4.5.2.

- 2) When a new header  $n$  arrives with  $T_n$  as the scaled TS, the compressor notes the arrival time  $a_n$ . It then calculates

$\text{Max\_Jitter\_BC} =$

$$\max \{ |(T_n - T_j) - ((a_n - a_j) / \text{TIME\_STRIDE})|, \\ \text{for all headers } j \text{ in the sliding window} \},$$

where  $\text{TIME\_STRIDE}$  is the time interval equivalent to one  $\text{TS\_STRIDE}$ , e.g., 20 ms.  $\text{Max\_Jitter\_BC}$  is the maximum observed jitter before the compressor, in units of  $\text{TS\_STRIDE}$ , for the headers in the sliding window.

- 3)  $k$  is calculated as

$$k = \text{ceiling}(\log_2(2 * J + 1)),$$

where  $J = \text{Max\_Jitter\_BC} + \text{Max\_Jitter\_CD} + 2$ .

$\text{Max\_Jitter\_CD}$  is the upper bound of jitter expected on the communication channel between compressor and decompressor (CD-CC). It depends only on the characteristics of CD-CC.



The constant 2 accounts for the quantization error introduced by the clocks at the compressor and decompressor, which can be +/-1.

Note that the calculation of  $k$  follows the compression algorithm described in section 4.5.1, with  $p = 2^{(k-1)} - 1$ .

- 4) The sliding window is subject to the same window operations as in section 4.5.2, 1) and 3), except that the values added and removed are paired with their arrival times.

Decompressor:

- 1) The decompressor uses as its reference header the last correctly (as verified by CRC) decompressed header. It maintains the pair  $(T_{ref}, a_{ref})$ , where  $T_{ref}$  is the scaled TS of the reference header, and  $a_{ref}$  is the arrival time of the reference header.
- 2) When receiving a compressed header  $n$  at time  $a_n$ , the approximation of the original scaled TS is calculated as:

$$T_{approx} = T_{ref} + (a_n - a_{ref}) / TIME\_STRIDE.$$

- 3) The approximation is then refined by the  $k$  least significant bits carried in header  $n$ , following the decompression algorithm of section 4.5.1, with  $p = 2^{(k-1)} - 1$ .

Note: The algorithm does not assume any particular pattern in the packets arriving at the compressor, i.e., it tolerates reordering before the compressor and nonincreasing RTP Timestamp behavior.

Note: Integer arithmetic is used in all equations above. If  $TIME\_STRIDE$  is not equal to an integral number of clock ticks, time must be normalized such that  $TIME\_STRIDE$  is an integral number of clock ticks. For example, if a clock tick is 20 ms and  $TIME\_STRIDE$  is 30 ms,  $(a_n - a_{ref})$  in 2) can be multiplied by 3 and  $TIME\_STRIDE$  can have the value 2.

Note: The clock resolution of the compressor or decompressor can be worse than  $TIME\_STRIDE$ , in which case the difference, i.e., actual resolution -  $TIME\_STRIDE$ , is treated as additional jitter in the calculation of  $k$ .

Note: The clock resolution of the decompressor may be communicated to the compressor using the `CLOCK` feedback option.

Note: The decompressor may observe the jitter and report this to the compressor using the `JITTER` feedback option. The compressor may use this information to refine its estimate of `Max_Jitter_CD`.

#### 4.5.5. Offset IP-ID encoding

As all IPv4 packets have an IP Identifier to allow for fragmentation, ROHC provides for transparent compression of this ID. There is no explicit support in ROHC for the IPv6 fragmentation header, so there is never a need to discuss IP IDs outside the context of IPv4.

This section assumes (initially) that the IPv4 stack at the source host assigns IP-ID according to the value of a 2-byte counter which is increased by one after each assignment to an outgoing packet. Therefore, the IP-ID field of a particular IPv4 packet flow will increment by 1 from packet to packet except when the source has emitted intermediate packets not belonging to that flow.

For such IPv4 stacks, the RTP SN will increase by 1 for each packet emitted and the IP-ID will increase by at least the same amount. Thus, it is more efficient to compress the offset, i.e., (IP-ID - RTP SN), instead of IP-ID itself.

The remainder of section 4.5.5 describes how to compress/decompress the sequence of offsets using W-LSB encoding/decoding, with  $p = 0$  (see section 4.5.1). All IP-ID arithmetic is done using unsigned 16-bit quantities, i.e., modulo  $2^{16}$ .

##### Compressor:

The compressor uses W-LSB encoding (section 4.5.2) to compress a sequence of offsets

$$\text{Offset}_i = \text{ID}_i - \text{SN}_i,$$

where  $\text{ID}_i$  and  $\text{SN}_i$  are the values of the IP-ID and RTP SN of header  $i$ . The sliding window contains such offsets and not the values of header fields, but the rules for adding and deleting offsets from the window otherwise follow section 4.5.2.

##### Decompressor:

The reference header is the last correctly (as verified by CRC) decompressed header.

When receiving a compressed packet  $m$ , the decompressor calculates  $\text{Offset}_{\text{ref}} = \text{ID}_{\text{ref}} - \text{SN}_{\text{ref}}$ , where  $\text{ID}_{\text{ref}}$  and  $\text{SN}_{\text{ref}}$  are the values of IP-ID and RTP SN in the reference header, respectively.

Then W-LSB decoding is used to decompress `Offset_m`, using the received LSBs in packet `m` and `Offset_ref`. Note that `m` may contain zero LSBs for `Offset_m`, in which case `Offset_m = Offset_ref`.

Finally, the IP-ID for packet `m` is regenerated as

IP-ID for `m` = decompressed SN of packet `m` + `Offset_m`

Network byte order:

Some IPv4 stacks do use a counter to generate IP ID values as described, but do not transmit the contents of this counter in network byte order, but instead send the two octets reversed. In this case, the compressor can compress the IP-ID field after swapping the bytes. Consequently, the decompressor also swaps the bytes of the IP-ID after decompression to regenerate the original IP-ID. This requires that the compressor and the decompressor synchronize on the byte order of the IP-ID field using the NBO or NBO2 flag (see section 5.7).

Random IP Identifier:

Some IPv4 stacks generate the IP Identifier values using a pseudo-random number generator. While this may provide some security benefits, it makes it pointless to attempt compressing the field. Therefore, the compressor should detect such random behavior of the field. After detection and synchronization with the decompressor using the RND or RND2 flag, the field is sent as-is in its entirety as additional octets after the compressed header.

#### 4.5.6. Self-describing variable-length values

The values of `TS_STRIDE` and a few other compression parameters can vary widely. `TS_STRIDE` can be 160 for voice and 90 000 for 1 f/s video. To optimize the transfer of such values, a variable number of octets is used to encode them. The number of octets used is determined by the first few bits of the first octet:

First bit is 0: 1 octet.

7 bits transferred.

Up to 127 decimal.

Encoded octets in hexadecimal: 00 to 7F

First bits are 10: 2 octets.

14 bits transferred.

Up to 16 383 decimal.

Encoded octets in hexadecimal: 80 00 to BF FF

First bits are 110: 3 octets.  
21 bits transferred.  
Up to 2 097 151 decimal.  
Encoded octets in hexadecimal: C0 00 00 to DF FF FF

First bits are 111: 4 octets.  
29 bits transferred.  
Up to 536 870 911 decimal.  
Encoded octets in hexadecimal: E0 00 00 00 to FF FF FF FF

#### 4.5.7. Encoded values across several fields in compressed headers

When a compressed header has an extension, pieces of an encoded value can be present in more than one field. When an encoded value is split over several fields in this manner, the more significant bits of the value are closer to the beginning of the header. If the number of bits available in compressed header fields exceeds the number of bits in the value, the most significant field is padded with zeroes in its most significant bits.

For example, an unscaled TS value can be transferred using an UOR-2 header (see section 5.7) with an extension of type 3. The Tsc bit of the extension is then unset (zero) and the variable length TS field of the extension is 4 octets, with 29 bits available for the TS (see section 4.5.6). The UOR-2 TS field will contain the three most significant bits of the unscaled TS, and the 4-octet TS field in the extension will contain the remaining 29 bits.

#### 4.6. Errors caused by residual errors

ROHC is designed under the assumption that packets can be damaged between the compressor and decompressor, and that such damaged packets can be delivered to the decompressor ("residual errors").

Residual errors may damage the SN in compressed headers. Such damage will cause generation of a header which upper layers may not be able to distinguish from a correct header. When the compressed header contains a CRC, the CRC will catch the bad header with a probability dependent on the size of the CRC. When ROHC does not detect the bad header, it will be delivered to upper layers.

Damage is not confined to the SN:

- a) Damage to packet type indication bits can cause a header to be interpreted as having a different packet type.

- b) Damage to CID information may cause a packet to be interpreted according to another context and possibly also according to another profile. Damage to CIDs will be more harmful when a large part of the CID space is being used, so that it is likely that the damaged CID corresponds to an active context.
- c) Feedback information can also be subject to residual errors, both when feedback is piggybacked and when it is sent in separate ROHC packets. ROHC uses sanity checks and adds CRCs to vital feedback information to allow detection of some damaged feedback.

Note that context damage can also result in generation of incorrect headers; section 4.7 elaborates further on this.

#### 4.7. Impairment considerations

Impairments to headers can be classified into the following types:

- (1) the lower layer was not able to decode the packet and did not deliver it to ROHC,
- (2) the lower layer was able to decode the packet, but discarded it because of a detected error,
- (3) ROHC detected an error in the generated header and discarded the packet, or
- (4) ROHC did not detect that the regenerated header was damaged and delivered it to upper layers.

Impairments cause loss or damage of individual headers. Some impairment scenarios also cause context invalidation, which in turn results in loss propagation and damage propagation. Damage propagation and undetected residual errors both contribute to the number of damaged headers delivered to upper layers. Loss propagation and impairments resulting in loss or discarding of single packets both contribute to the packet loss seen by upper layers.

Examples of context invalidating scenarios are:

- (a) Impairment of type (4) on the forward channel, causing the decompressor to update its context with incorrect information;

- (b) Loss/error burst of pattern update headers: Impairments of types (1),(2) and (3) on consecutive pattern update headers; a pattern update header is a header carrying a new pattern information, e.g., at the beginning of a new talk spurt; this causes the decompressor to lose the pattern update information;
- (c) Loss/error burst of headers: Impairments of types (1),(2) and (3) on a number of consecutive headers that is large enough to cause the decompressor to lose the SN synchronization;
- (d) Impairment of type (4) on the feedback channel which mimics a valid ACK and makes the compressor update its context;
- (e) a burst of damaged headers (3) erroneously triggers the "k-out-of-n" rule for detecting context invalidation, which results in a NACK/update sequence during which headers are discarded.

Scenario (a) is mitigated by the CRC carried in all context updating headers. The larger the CRC, the lower the chance of context invalidation caused by (a). In R-mode, the CRC of context updating headers is always 7 bits or more. In U/O-mode, it is usually 3 bits and sometimes 7 or 8 bits.

Scenario (b) is almost completely eliminated when the compressor ensures through ACKs that no context updating headers are lost, as in R-mode.

Scenario (c) is almost completely eliminated when the compressor ensures through ACKs that the decompressor will always detect the SN wraparound, as in R-mode. It is also mitigated by the SN repair mechanisms in U/O-mode.

Scenario (d) happens only when the compressor receives a damaged header that mimics an ACK of some header present in the W-LSB window, say ACK of header 2, while in reality header 2 was never received or accepted by the decompressor, i.e., header 2 was subject to impairment (1), (2) or (3). The damaged header must mimic the feedback packet type, the ACK feedback type, and the SN LSBs of some header in the W-LSB window.

Scenario (e) happens when a burst of residual errors causes the CRC check to fail in k out of the last n headers carrying CRCs. Large k and n reduces the probability of scenario (e), but also increases the number of headers lost or damaged as a consequence of any context invalidation.

ROHC detects damaged headers using CRCs over the original headers. The smallest headers in this document either include a 3-bit CRC (U/O-mode) or do not include a CRC (R-mode). For the smallest headers, damage is thus detected with a probability of roughly 7/8 for U/O-mode. For R-mode, damage to the smallest headers is not detected.

All other things (coding scheme at lower layers, etc.) being equal, the rate of headers damaged by residual errors will be lower when headers are compressed compared when they are not, since fewer bits are transmitted. Consequently, for a given ROHC CRC setup the rate of incorrect headers delivered to applications will also be reduced.

The above analysis suggests that U/O-mode may be more prone than R-mode to context invalidation. On the other hand, the CRC present in all U/O-mode headers continuously screens out residual errors coming from lower layers, reduces the number of damaged headers delivered to upper layers when context is invalidated, and permits quick detection of context invalidation.

R-mode always uses a stronger CRC on context updating headers, but no CRC in other headers. A residual error on a header which carries no CRC will result in a damaged header being delivered to upper layers (4). The number of damaged headers delivered to the upper layers depends on the ratio of headers with CRC vs. headers without CRC, which is a compressor parameter.

## 5. The protocol

### 5.1. Data structures

The ROHC protocol is based on a number of parameters that form part of the negotiated channel state and the per-context state. This section describes some of this state information in an abstract way. Implementations can use a different structure for and representation of this state. In particular, negotiation protocols that set up the per-channel state need to establish the information that constitutes the negotiated channel state, but it is not necessary to exchange it in the form described here.

#### 5.1.1. Per-channel parameters

**MAX\_CID:** Nonnegative integer; highest context ID number to be used by the compressor (note that this parameter is not coupled to, but in effect further constrained by, **LARGE\_CIDS**).

LARGE\_CIDS: Boolean; if false, the short CID representation (0 bytes or 1 prefix byte, covering CID 0 to 15) is used; if true, the embedded CID representation (1 or 2 embedded CID bytes covering CID 0 to 16383) is used.

PROFILES: Set of nonnegative integers, each integer indicating a profile supported by the decompressor. The compressor MUST NOT compress using a profile not in PROFILES.

FEEDBACK\_FOR: Optional reference to a channel in the reverse direction. If provided, this parameter indicates which channel any feedback sent on this channel refers to (see 5.7.6.1).

MRRU: Maximum reconstructed reception unit. This is the size of the largest reconstructed unit in octets that the decompressor is expected to reassemble from segments (see 5.2.5). Note that this size includes the CRC. If MRRU is negotiated to be 0, no segment headers are allowed on the channel.

#### 5.1.2. Per-context parameters, profiles

Per-context parameters are established with IR headers (see section 5.2.3). An IR header contains a profile identifier, which determines how the rest of the header is to be interpreted. Note that the profile parameter determines the syntax and semantics of the packet type identifiers and packet types used in conjunction with a specific context. This document describes profiles 0x0000, 0x0001, 0x0002, and 0x0003; further profiles may be defined when ROHC is extended in the future.

Profile 0x0000 is for sending uncompressed IP packets. See section 5.10.

Profile 0x0001 is for RTP/UDP/IP compression, see sections 5.3 through 5.9.

Profile 0x0002 is for UDP/IP compression, i.e., compression of the first 12 octets of the UDP payload is not attempted. See section 5.11.

Profile 0x0003 is for ESP/IP compression, i.e., compression of the header chain up to and including the first ESP header, but not subsequent subheaders. See section 5.12.

Initially, all contexts are in no context state, i.e., all packets referencing this context except IR packets are discarded. If defined by a "ROHC over X" document, per-channel negotiation can be used to pre-establish state information for a context (e.g., negotiating



profile 0x0000 for CID 15). Such state information can also be marked read-only in the negotiation, which would cause the decompressor to discard any IR packet attempting to modify it.

### 5.1.3. Contexts and context identifiers

Associated with each compressed flow is a context, which is the state compressor and decompressor maintain in order to correctly compress or decompress the headers of the packet stream. Contexts are identified by a context identifier, CID, which is sent along with compressed headers and feedback information.

The CID space is distinct for each channel, i.e., CID 3 over channel A and CID 3 over channel B do not refer to the same context, even if the endpoints of A and B are the same nodes. In particular, CIDs for any pairs of forward and reverse channels are not related (forward and reverse channels need not even have CID spaces of the same size).

Context information is conceptually kept in a table. The context table is indexed using the CID which is sent along with compressed headers and feedback information. The CID space can be negotiated to be either small, which means that CIDs can take the values 0 through 15, or large, which means that CIDs take values between 0 and  $2^{14} - 1 = 16383$ . Whether the CID space is large or small is negotiated no later than when a channel is established.

A small CID with the value 0 is represented using zero bits. A small CID with a value from 1 to 15 is represented by a four-bit field in place of a packet type field (Add-CID) plus four more bits. A large CID is represented using the encoding scheme of section 4.5.6, limited to two octets.

## 5.2. ROHC packets and packet types

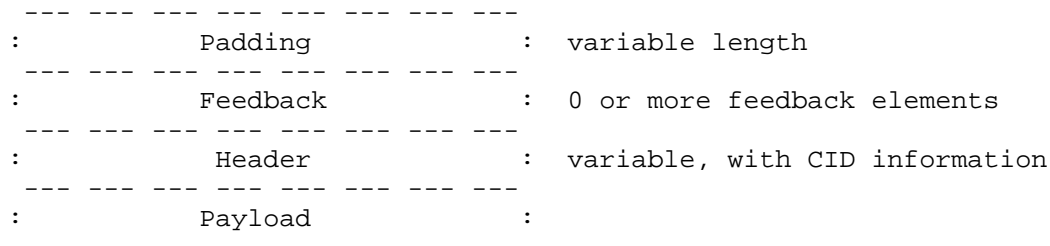
The packet type indication scheme for ROHC has been designed under the following constraints:

- a) it must be possible to use only a limited number of packet sizes;
- b) it must be possible to send feedback information in separate ROHC packets as well as piggybacked on forward packets;
- c) it is desirable to allow elimination of the CID for one packet stream when few packet streams share a channel;
- d) it is anticipated that some packets with large headers may be larger than the MTU of very constrained lower layers.

These constraints have led to a design which includes

- optional padding,
- a feedback packet type,
- an optional Add-CID octet which provides 4 bits of CID, and
- a simple segmentation and reassembly mechanism.

A ROHC packet has the following general format (in the diagram, colons ":" indicate that the part is optional):



Padding is any number (zero or more) of padding octets. Either of Feedback or Header must be present.

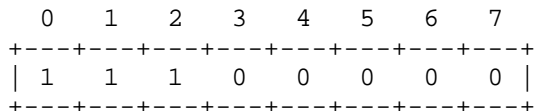
Feedback elements always start with a packet type indication. Feedback elements carry internal CID information. Feedback is described in section 5.2.2.

Header is either a profile-specific header or an IR or IR-DYN header (see sections 5.2.3 and 5.2.4). Header either

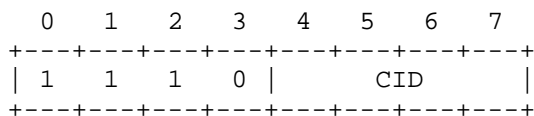
- 1) does not carry any CID information (indicating CID zero), or
- 2) includes one Add-CID Octet (see below), or
- 3) contains embedded CID information of length one or two octets.

Alternatives 1) and 2) apply only to compressed headers in channels where the CID space is small. Alternative 3) applies only to compressed headers in channels where the CID space is large.

Padding Octet



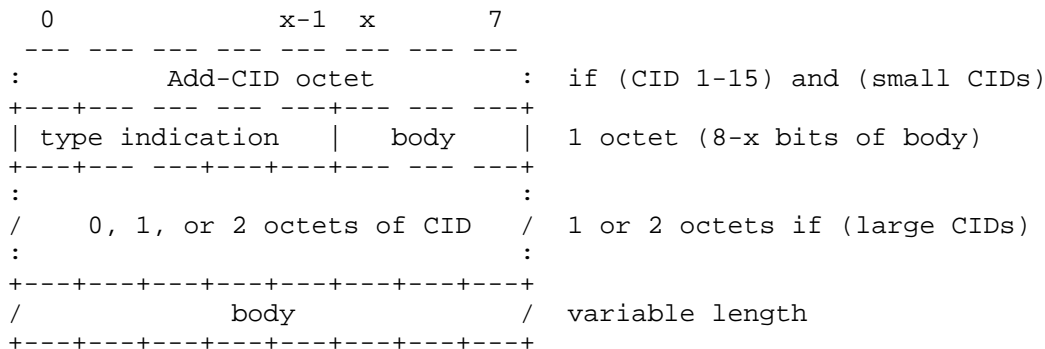
Add-CID Octet



CID: 0x1 through 0xF indicates CIDs 1 through 15.

Note: The Padding Octet looks like an Add-CID octet for CID 0.

Header either starts with a packet type indication or has a packet type indication immediately following an Add-CID Octet. All Header packet types have the following general format (in the diagram, slashes "/" indicate variable length):



The large CID, if present, is encoded according to section 4.5.6.

5.2.1. ROHC feedback

Feedback carries information from decompressor to compressor. The following principal kinds of feedback are supported. In addition to the kind of feedback, other information may be included in profile-specific feedback information.

- ACK : Acknowledges successful decompression of a packet, which means that the context is up-to-date with a high probability.
- NACK : Indicates that the dynamic context of the decompressor is out of sync. Generated when several successive packets have failed to be decompressed correctly.

STATIC-NACK : Indicates that the static context of the decompressor is not valid or has not been established.

It is anticipated that feedback to the compressor can be realized in many ways, depending on the properties of the particular lower layer. The exact details of how feedback is realized is to be specified in a "ROHC over X" document, for each lower layer X in question. For example, feedback might be realized using

- 1) lower-layer specific mechanisms
- 2) a dedicated feedback-only channel, realized for example by the lower layer providing a way to indicate that a packet is a feedback packet
- 3) a dedicated feedback-only channel, where the timing of the feedback provides information about which compressed packet caused the feedback
- 4) interspersing of feedback packets among normal compressed packets going in the same direction as the feedback (lower layers do not indicate feedback)
- 5) piggybacking of feedback information in compressed packets going in the same direction as the feedback (this technique may reduce the per-feedback overhead)
- 6) interspersing and piggybacking on the same channel, i.e., both 4) and 5).

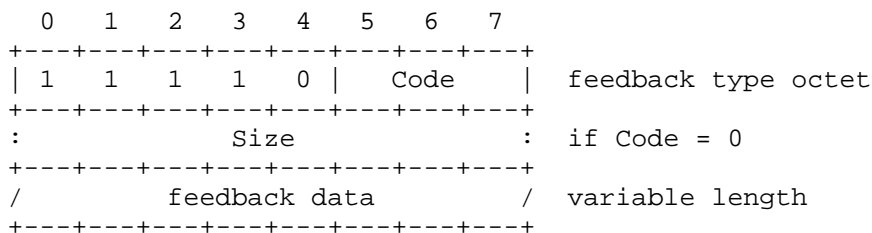
Alternatives 1-3 do not place any particular requirements on the ROHC packet type scheme. Alternatives 4-6 do, however. The ROHC packet type scheme has been designed to allow alternatives 4-6 (these may be used for example over PPP):

- a) The ROHC scheme provides a feedback packet type. The packet type is able to carry variable-length feedback information.
- b) The feedback information sent on a particular channel is passed to, and interpreted by, the compressor associated with feedback on that channel. Thus, the feedback information must contain CID information if the associated compressor can use more than one context. The ROHC feedback scheme requires that a channel carries feedback to at most one compressor. How a compressor is associated with feedback on a particular channel needs to be defined in a "ROHC over X" document.

- c) The ROHC feedback information format is octet-aligned, i.e., starts at an octet boundary, to allow using the format over a dedicated feedback channel, 2).
- d) To allow piggybacking, 5), it is possible to deduce the length of feedback information by examining the first few octets of the feedback. This allows the decompressor to pass piggybacked feedback information to the associated same-side compressor without understanding its format. The length information decouples the decompressor from the compressor in the sense that the decompressor can process the compressed header immediately without waiting for the compressor to hand it back after parsing the feedback information.

5.2.2. ROHC feedback format

Feedback sent on a ROHC channel consists of one or more concatenated feedback elements, where each feedback element has the following format:



Code: 0 indicates that a Size octet is present.  
 1-7 indicates the size of the feedback data field in octets.

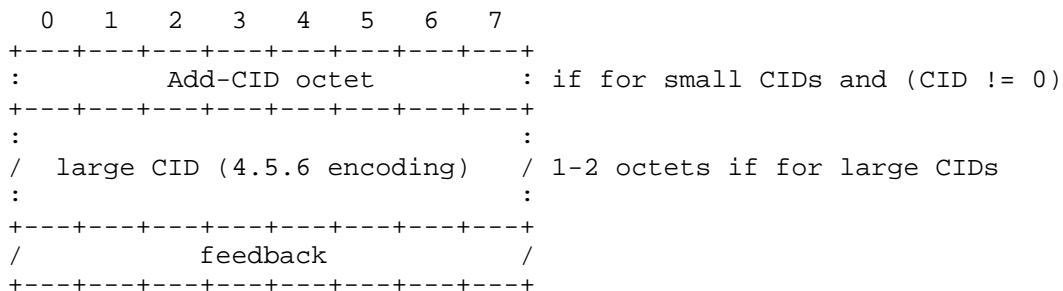
Size: Optional octet indicating the size of the feedback data field in octets.

feedback data: Profile-specific feedback information. Includes CID information.

The total size of the feedback data field is determinable upon reception by the decompressor, by inspection of the Code field and possibly the Size field. This explicit length information allows piggybacking and also sending more than one feedback element in a packet.

When the decompressor has determined the size of the feedback data field, it removes the feedback type octet and the Size field (if present) and hands the rest to the same-side associated compressor

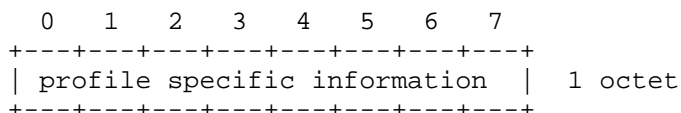
together with an indication of the size. The feedback data received by the compressor has the following structure (feedback sent on a dedicated feedback channel MAY also use this format):



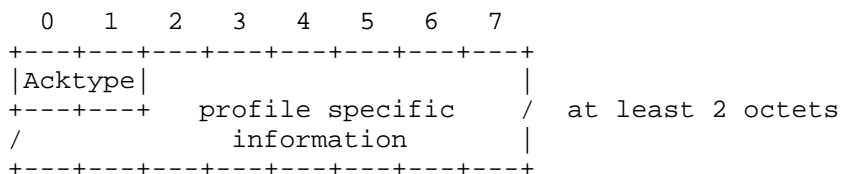
The large CID, if present, is encoded according to section 4.5.6. CID information in feedback data indicates the CID of the packet stream for which feedback is sent. Note that the LARGE\_CIDS parameter that controls whether a large CID is present is taken from the channel state of the receiving compressor's channel, NOT from that of the channel carrying the feedback.

It is REQUIRED that the feedback field have either of the following two formats:

FEEDBACK-1



FEEDBACK-2



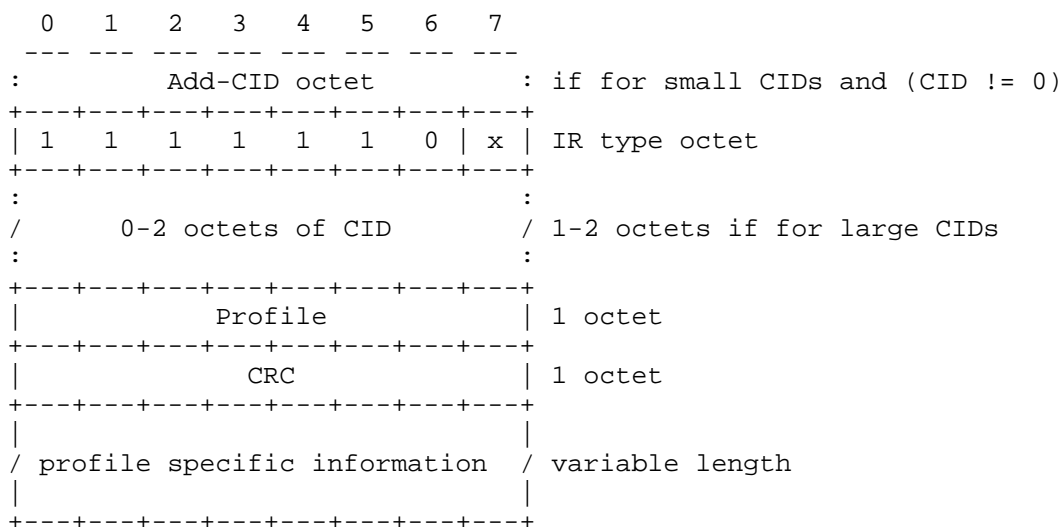
- Acktype: 0 = ACK
- 1 = NACK
- 2 = STATIC-NACK
- 3 is reserved (MUST NOT be used. Otherwise unparseable.)

The compressor can use the following logic to parse the feedback field.

- 1) If for large CIDs, the feedback will always start with a CID encoded according to section 4.5.6. If the first bit is 0, the CID uses one octet. If the first bit is 1, the CID uses two octets.
- 2) If for small CIDs, and the size is one octet, the feedback is a FEEDBACK-1.
- 3) If for small CIDs, and the size is larger than one octet, and the feedback starts with the two bits 11, the feedback starts with an Add-CID octet. If the size is 2, it is followed by FEEDBACK-1. If the size is larger than 2, the Add-CID is followed by FEEDBACK-2.
- 4) Otherwise, there is no Add-CID octet, and the feedback starts with a FEEDBACK-2.

5.2.3. ROHC IR packet type

The IR header associates a CID with a profile, and typically also initializes the context. It can typically also refresh (parts of) the context. It has the following general format.



x: Profile specific information. Interpreted according to the profile indicated in the Profile field.

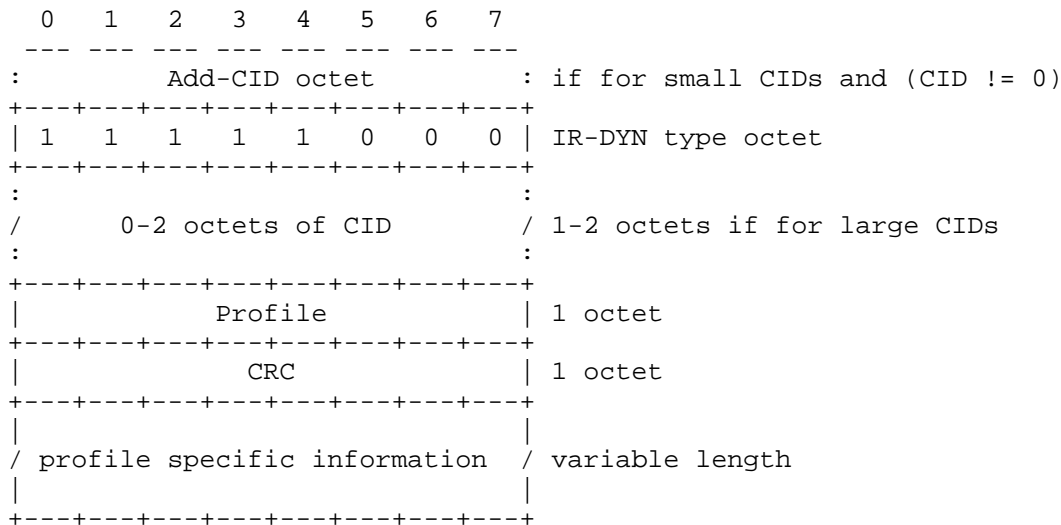
Profile: The profile to be associated with the CID. In the IR packet, the profile identifier is abbreviated to the 8 least significant bits. It selects the highest-number profile in the channel state parameter PROFILES that matches the 8 LSBs given.

CRC: 8-bit CRC computed using the polynomial of section 5.9.1. Its coverage is profile-dependent, but it MUST cover at least the initial part of the packet ending with the Profile field. Any information which initializes the context of the decompressor should be protected by the CRC.

Profile specific information: The contents of this part of the IR packet are defined by the individual profiles. Interpreted according to the profile indicated in the Profile field.

5.2.4. ROHC IR-DYN packet type

In contrast to the IR header, the IR-DYN header can never initialize an uninitialized context. However, it can redefine what profile is associated with a context, see for example 5.11 (ROHC UDP) and 5.12 (ROHC ESP). Thus the type needs to be reserved at the framework level. The IR-DYN header typically also initializes or refreshes parts of a context, typically the dynamic part. It has the following general format:



Profile: The profile to be associated with the CID. This is abbreviated in the same way as with IR packets.



CRC: 8-bit CRC computed using the polynomial of section 5.9.1. Its coverage is profile-dependent, but it MUST cover at least the initial part of the packet ending with the Profile field. Any information which initializes the context of the decompressor should be protected by the CRC.

Profile specific information: This part of the IR packet is defined by individual profiles. It is interpreted according to the profile indicated in the Profile field.

#### 5.2.5. ROHC segmentation

Some link layers may provide a much more efficient service if the set of different packet sizes to be transported is kept small. For such link layers, these sizes will normally be chosen to transport frequently occurring packets efficiently, with less frequently occurring packets possibly adapted to the next larger size by the addition of padding. The link layer may, however, be limited in the size of packets it can offer in this efficient mode, or it may be desirable to request only a limited largest size. To accommodate the occasional packet that is larger than that largest size negotiated, ROHC defines a simple segmentation protocol.

##### 5.2.5.1. Segmentation usage considerations

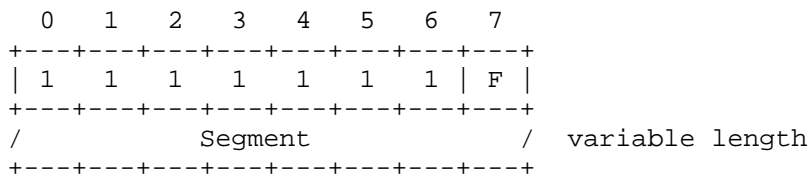
The segmentation protocol defined in ROHC is not particularly efficient. It is not intended to replace link layer segmentation functions; these SHOULD be used whenever available and efficient for the task at hand.

ROHC segmentation should only be used for occasional packets with sizes larger than what is efficient to accommodate, e.g., due to exceptionally large ROHC headers. The segmentation scheme was designed to reduce packet size variations that may occur due to outliers in the header size distribution. In other cases, segmentation should be done at lower layers. The segmentation scheme should only be used for packet sizes that are larger than the maximum size in the allowed set of sizes from the lower layers.

In summary, ROHC segmentation should be used with a relatively low frequency in the packet flow. If this cannot be ensured, segmentation should be performed at lower layers.

5.2.5.2. Segmentation protocol

Segment Packet

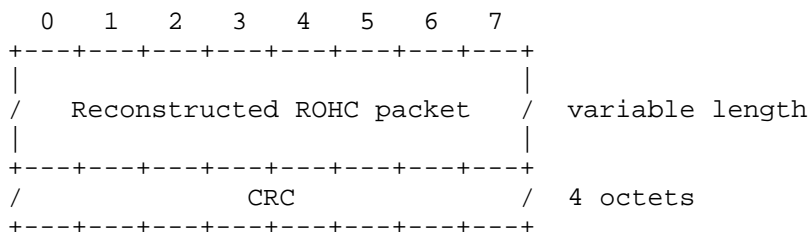


F: Final bit. If set, it indicates that this is the last segment of a reconstructed unit.

The segment header may be preceded by padding octets and/or feedback. It never carries a CID.

All segment header packets for one reconstructed unit have to be sent consecutively on a channel, i.e., any non-segment-header packet following a nonfinal segment header aborts the reassembly of the current reconstructed unit and causes the decompressor to discard the nonfinal segments received on this channel so far. When a final segment header is received, the decompressor reassembles the segment carried in this packet and any nonfinal segments that immediately preceded it into a single reconstructed unit, in the order they were received. The reconstructed unit has the format:

Reconstructed Unit



The CRC is used by the decompressor to validate the reconstructed unit. It uses the FCS-32 algorithm with the following generator polynomial:  $x^0 + x^1 + x^2 + x^4 + x^5 + x^7 + x^8 + x^{10} + x^{11} + x^{12} + x^{16} + x^{22} + x^{23} + x^{26} + x^{32}$  [HDLC]. If the reconstructed unit is 4 octets or less, or if the CRC fails, or if it is larger than the channel parameter MRRU (see 5.1.1), the reconstructed unit MUST be discarded by the decompressor.

If the CRC succeeds, the reconstructed ROHC packet is interpreted as a ROHC Header, optionally followed by a payload. Note that this means that there can be no padding and no feedback in the reconstructed unit, and that the CID is derived from the initial octets of the reconstructed unit.

(It should be noted that the ROHC segmentation protocol was inspired by SEAL by Steve Deering et al., which later became ATM AAL5. The same arguments for not having sequence numbers in the segments but instead providing a strong CRC in the reconstructed unit apply here as well. Note that, as a result of this protocol, there is no way in ROHC to make any use of a segment that has residual bit errors.)

#### 5.2.6. ROHC initial decompressor processing

The following packet types are reserved at the framework level in the ROHC scheme:

1110:       Padding or Add-CID octet  
11110:      Feedback  
11111000:   IR-DYN packet  
1111110:    IR packet  
1111111:    Segment

Other packet types can be used at will by individual profiles.

The following steps is an outline of initial decompressor processing which upon reception of a ROHC packet can determine its contents.

- 1) If the first octet is a Padding Octet (11100000), strip away all initial Padding Octets and goto next step.
- 2) If the first remaining octet starts with 1110, it is an Add-CID octet:  
  
    remember the Add-CID octet; remove the octet.
- 3) If the first remaining octet starts with 11110, and an Add-CID octet was found in step 2),  
  
    an error has occurred; the header MUST be discarded without further action.
- 4) If the first remaining octet starts with 11110, and an Add-CID octet was not found in step 2), this is feedback:  
  
    find the size of the feedback data, call it *s*;  
    remove the feedback type octet;

remove the Size octet if Code is 0;  
send feedback data of length *s* to the same-side associated  
compressor;  
if packet exhausted, stop; otherwise goto 2).

- 5) If the first remaining octet starts with 11111111, this is a segment:
  - attempt reconstruction using the segmentation protocol (5.2.5). If a reconstructed packet is not produced, this finishes the processing of the original packet. If a reconstructed packet is produced, it is fed into step 1) above. Padding, segments, and feedback are not allowed in reconstructed packets, so when processing them, steps 1), 4), and 5) are modified so that the packet is discarded without further action when their conditions match.
- 6) Here, it is known that the rest is forward information (unless the header is damaged).
- 7) If the forward traffic uses small CIDs, there is no large CID in the packet. If an Add-CID immediately preceded the packet type (step 2), it has the CID of the Add-CID; otherwise it has CID 0.
- 8) If the forward traffic uses large CIDs, the CID starts with the second remaining octet. If the first bit(s) of that octet are not 0 or 10, the packet MUST be discarded without further action. If an Add-CID octet immediately preceded the packet type (step 2), the packet MUST be discarded without further action.
- 9) Use the CID to find the context.
- 10) If the packet type is IR, the profile indicated in the IR packet determines how it is to be processed. If the CRC fails to verify the packet, it MUST be discarded. If a profile is indicated in the context, the logic of that profile determines what, if any, feedback is to be sent. If no profile is noted in the context, no further action is taken.
- 11) If the packet type is IR-DYN, the profile indicated in the IR-DYN packet determines how it is to be processed.
  - a) If the CRC fails to verify the packet, it MUST be discarded. If a profile is indicated in the context, the logic of that profile determines what, if any, feedback is to be sent. If no profile is noted in the context, no further action is taken.

- b) If the context has not been initialized by an IR packet, the packet MUST be discarded. The logic of the profile indicated in the IR-DYN header (if verified by the CRC), determines what, if any, feedback is to be sent.
- 12) Otherwise, the profile noted in the context determines how the rest of the packet is to be processed. If the context has not been initialized by an IR packet, the packet MUST be discarded without further action.

The procedure for finding the size of the feedback data is as follows:

Examine the three bits which immediately follow the feedback packet type. When these bits are

- 1-7, the size of the feedback data is given by the bits;
- 0, a Size octet, which explicitly gives the size of the feedback data, is present after the feedback type octet.

#### 5.2.7. ROHC RTP packet formats from compressor to decompressor

ROHC RTP uses three packet types to identify compressed headers, and two for initialization/refresh. The format of a compressed packet can depend on the mode. Therefore a naming scheme of the form

<modes format is used in>-<packet type number>-<some property>

is used to uniquely identify the format when necessary, e.g., UOR-2, R-1. For exact formats of the packet types, see section 5.7.

Packet type zero: R-0, R-0-CRC, UO-0.

This, the minimal, packet type is used when parameters of all SN-functions are known by the decompressor, and the header to be compressed adheres to these functions. Thus, only the W-LSB encoded RTP SN needs to be communicated.

R-mode: Only if a CRC is present (packet type R-0-CRC) may the header be used as a reference for subsequent decompression.

U-mode and O-mode: A small CRC is present in the UO-0 packet.

Packet type 1: R-1, R-1-ID, R-1-TS, UO-1, UO-1-ID, UO-1-TS.

This packet type is used when the number of bits needed for the SN exceeds those available in packet type zero, or when the parameters of the SN-functions for RTP TS or IP-ID change.

R-mode: R-1-\* packets are not used as references for subsequent decompression. Values for other fields than the RTP TS or IP-ID can be communicated using an extension, but they do not update the context.

U-mode and O-mode: Only the values of RTP SN, RTP TS and IP-ID can be used as references for future compression. Nonupdating values can be provided for other fields using an extension (UO-1-ID).

Packet type 2: UOR-2, UOR-2-ID, UOR-2-TS

This packet type can be used to change the parameters of any SN-function, except those for most static fields. Headers of packets transferred using packet type 2 can be used as references for subsequent decompression.

Packet type: IR

This packet type communicates the static part of the context, i.e., the value of the constant SN-functions. It can optionally also communicate the dynamic part of the context, i.e., the parameters of the nonconstant SN-functions.

Packet type: IR-DYN

This packet type communicates the dynamic part of the context, i.e., the parameters of nonconstant SN-functions.

#### 5.2.8. Parameters needed for mode transition in ROHC RTP

The packet types IR (with dynamic information), IR-DYN, and UOR-2 are common for all modes. They can carry a mode parameter which can take the values U = Unidirectional, O = Bidirectional Optimistic, and R = Bidirectional Reliable.

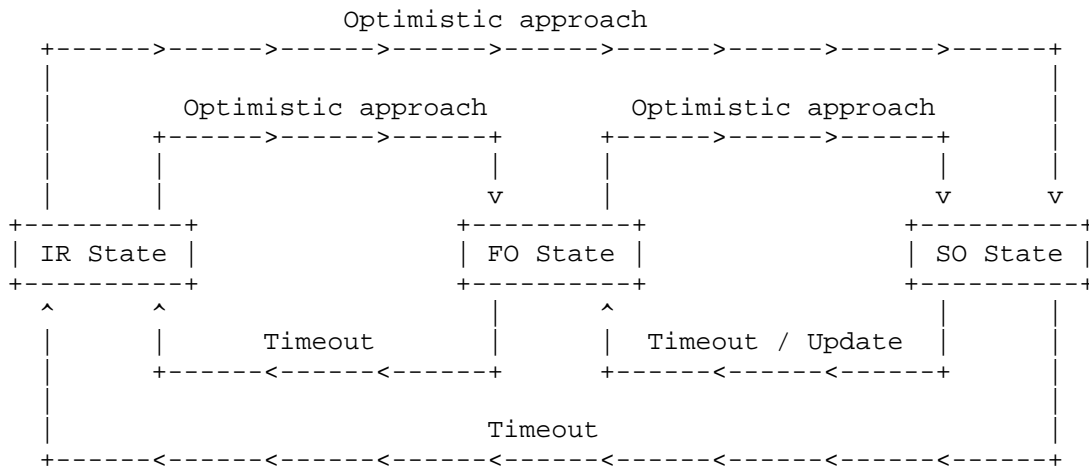
Feedback of types ACK, NACK, and STATIC-NACK carry sequence numbers, and feedback packets can also carry a mode parameter indicating the desired compression mode: U, O, or R.

As a shorthand, the notation PACKET(mode) is used to indicate which mode value a packet carries. For example, an ACK with mode parameter R is written ACK(R), and an UOR-2 with mode parameter O is written UOR-2(O).

5.3. Operation in Unidirectional mode

5.3.1. Compressor states and logic (U-mode)

Below is the state machine for the compressor in Unidirectional mode. Details of the transitions between states and compression logic are given subsequent to the figure.



5.3.1.1. State transition logic (U-mode)

The transition logic for compression states in Unidirectional mode is based on three principles: the optimistic approach principle, timeouts, and the need for updates.

5.3.1.1.1. Optimistic approach, upwards transition

Transition to a higher compression state in Unidirectional mode is carried out according to the optimistic approach principle. This means that the compressor transits to a higher compression state when it is fairly confident that the decompressor has received enough information to correctly decompress packets sent according to the higher compression state.

When the compressor is in the IR state, it will stay there until it assumes that the decompressor has correctly received the static context information. For transition from the FO to the SO state, the compressor should be confident that the decompressor has all parameters needed to decompress according to a fixed pattern.

The compressor normally obtains its confidence about decompressor status by sending several packets with the same information according to the lower compression state. If the decompressor receives any of these packets, it will be in sync with the compressor. The number of consecutive packets to send for confidence is not defined in this document.

#### 5.3.1.1.2. Timeouts, downward transition

When the optimistic approach is taken as described above, there will always be a possibility of failure since the decompressor may not have received sufficient information for correct decompression. Therefore, the compressor **MUST** periodically transit to lower compression states. Periodic transition to the IR state **SHOULD** be carried out less often than transition to the FO state. Two different timeouts **SHOULD** therefore be used for these transitions. For an example of how to implement periodic refreshes, see [IPHC] chapters 3.3.1-3.3.2.

#### 5.3.1.1.3. Need for updates, downward transition

In addition to the downward state transitions carried out due to periodic timeouts, the compressor must also immediately transit back to the FO state when the header to be compressed does not conform to the established pattern.

#### 5.3.1.2. Compression logic and packets used (U-mode)

The compressor chooses the smallest possible packet format that can communicate the desired changes, and has the required number of bits for W-LSB encoded values.

#### 5.3.1.3. Feedback in Unidirectional mode

The Unidirectional mode of operation is designed to operate over links where a feedback channel is not available. If a feedback channel is available, however, the decompressor **MAY** send an acknowledgment of successful decompression with the mode parameter set to U (send an ACK(U)). When the compressor receives such a message, it **MAY** disable (or increase the interval between) periodic IR refreshes.

#### 5.3.2. Decompressor states and logic (U-mode)

Below is the state machine for the decompressor in Unidirectional mode. Details of the transitions between states and decompression logic are given subsequent to the figure.





The sequence number is reconstructed by replacing the sequence number LSBs in the context with those received in the header. The resulting value is then verified to be within the interpretation interval by comparison with a previously reconstructed reference value `v_ref` (see section 4.5.1). If it is not within this interval, an adjustment is applied by adding  $N \times \text{interval\_size}$  to the reconstructed value so that the result is brought within the interpretation interval. Note that  $N$  can be negative.

If RTP Timestamp and IP Identification fields are not included in the received header, they are supposed to be calculated from the sequence number. The IP Identifier usually increases by the same delta as the sequence number and the timestamp by the same delta times a fixed value. See chapters 4.5.3 and 4.5.5 for details about how these fields are encoded in compressed headers.

When working in Unidirectional mode, all compressed headers carry a CRC which MUST be used to verify decompression.

#### 5.3.2.2.3. Actions upon CRC failure

This section is written so that it is applicable to all modes.

A mismatch in the CRC can be caused by one or more of:

1. residual bit errors in the current header
2. a damaged context due to residual bit errors in previous headers
3. many consecutive packets being lost between compressor and decompressor (this may cause the LSBs of the SN in compressed packets to be interpreted wrongly, because the decompressor has not moved the interpretation interval for lack of input -- in essence, a kind of context damage).

(Cases 2 and 3 do not apply to IR packets; case 3 does not apply to IR-DYN packets.) The 3-bit CRC present in some header formats will eventually detect context damage reliably, since the probability of undetected context damage decreases exponentially with each new header processed. However, residual bit errors in the current header are only detected with good probability, not reliably.

When a CRC mismatch is caused by residual bit errors in the current header (case 1 above), the decompressor should stay in its current state to avoid unnecessary loss of subsequent packets. On the other hand, when the mismatch is caused by a damaged context (case 2), the decompressor should attempt to repair the context locally. If the local repair attempt fails, it must move to a lower state to avoid

delivering incorrect headers. When the mismatch is caused by prolonged loss (case 3), the decompressor might attempt additional decompression attempts. Note that case 3 does not occur in R-mode.

The following actions MUST be taken when a CRC check fails:

First, attempt to determine whether SN LSB wraparound (case 3) is likely, and if so, attempt a correction. For this, the algorithm of section 5.3.2.2.4 MAY be used. If another algorithm is used, it MUST have at least as high a rate of correct repairs as the one in 5.3.2.2.4. (This step is not applicable to R-mode.)

Second, if the previous step did not attempt a correction, a repair should be attempted under the assumption that the reference SN has been incorrectly updated. For this, the algorithm of section 5.3.2.2.5 MAY be used. If another algorithm is used, it MUST have at least as high a rate of correct repairs as the one in 5.3.2.2.5. (This step is not applicable to R-mode.)

If both the above steps fail, additional decompression attempts SHOULD NOT be made. There are two possible reasons for the CRC failure: case 1 or unrecoverable context damage. It is impossible to know for certain which of these is the actual cause. The following rules are to be used:

- a. When CRC checks fail only occasionally, assume residual errors in the current header and simply discard the packet. NACKs SHOULD NOT be sent at this time.
- b. In the Full Context state: When the CRC check of  $k_1$  out of the last  $n_1$  decompressed packets have failed, context damage SHOULD be assumed and a NACK SHOULD be sent in O- and R-mode. The decompressor moves to the Static Context state and discards all packets until an update (IR, IR-DYN, UOR-2) which passes the CRC check is received.
- c. In the Static Context state: When the CRC check of  $k_2$  out of the last  $n_2$  updates (IR, IR-DYN, UOR-2) have failed, static context damage SHOULD be assumed and a STATIC-NACK is sent in O- and R-mode. The decompressor moves to the No Context state.
- d. In the No Context state: The decompressor discards all packets until a static update (IR) which passes the CRC check is received. (In O-mode and R-mode, feedback is sent according to sections 5.4.2.2 and 5.5.2.2, respectively.)

Note that appropriate values for  $k_1$ ,  $n_1$ ,  $k_2$ , and  $n_2$ , are related to the residual error rate of the link. When the residual error rate is close to zero,  $k_1 = n_1 = k_2 = n_2 = 1$  may be appropriate.

#### 5.3.2.2.4. Correction of SN LSB wraparound

When many consecutive packets are lost there will be a risk of sequence number LSB wraparound, i.e., the SN LSBs being interpreted wrongly because the interpretation interval has not moved for lack of input. The decompressor might be able to detect this situation and avoid context damage by using a local clock. The following algorithm MAY be used:

- a. The decompressor notes the arrival time,  $a(i)$ , of each incoming packet  $i$ . Arrival times of packets where decompression fails are discarded.
- b. When decompression fails, the decompressor computes  $INTERVAL = a(i) - a(i - 1)$ , i.e., the time elapsed between the arrival of the previous, correctly decompressed packet and the current packet.
- c. If wraparound has occurred,  $INTERVAL$  will correspond to at least  $2^k$  inter-packet times, where  $k$  is the number of SN bits in the current header. On the basis of an estimate of the packet inter-arrival time, obtained for example using a moving average of arrival times,  $TS\_STRIDE$ , or  $TS\_TIME$ , the decompressor judges if  $INTERVAL$  can correspond to  $2^k$  inter-packet times.
- d. If  $INTERVAL$  is judged to be at least  $2^k$  packet inter-arrival times, the decompressor adds  $2^k$  to the reference SN and attempts to decompress the packet using the new reference SN.
- e. If this decompression succeeds, the decompressor updates the context but SHOULD NOT deliver the packet to upper layers. The following packet is also decompressed and updates the context if its CRC succeeds, but SHOULD be discarded. If decompression of the third packet using the new context also succeeds, the context repair is deemed successful and this and subsequent decompressed packets are delivered to the upper layers.
- f. If any of the three decompression attempts in d. and e. fails, the decompressor discards the packets and acts according to rules a) through c) of section 5.3.2.2.3.

Using this mechanism, the decompressor may be able to repair the context after excessive loss, at the expense of discarding two packets.

## 5.3.2.2.5. Repair of incorrect SN updates

The CRC can fail to detect residual errors in the compressed header because of its limited length, i.e., the incorrectly decompressed packet can happen to have the same CRC as the original uncompressed packet. The incorrect decompressed header will then update the context. This can lead to an erroneous reference SN being used in W-LSB decoding, as the reference SN is updated for each successfully decompressed header of certain types.

In this situation, the decompressor will detect the incorrect decompression of the following packet with high probability, but it does not know the reason for the failure. The following mechanism allows the decompressor to judge if the context was updated incorrectly by an earlier packet and, if so, to attempt a repair.

- a. The decompressor maintains two decompressed sequence numbers: the last one (ref 0) and the one before that (ref -1).
- b. When receiving a compressed header the SN (SN curr1) is decompressed using ref 0 as the reference. The other header fields are decompressed using this decompressed SN curr1. (This is part of the normal decompression procedure prior to any CRC test failures.)
- c. If the decompressed header generated in b. passes the CRC test, the references are shifted as follows:  
  
    ref -1 = ref 0  
    ref 0 = SN curr1.
- d. If the header generated in b. does not pass the CRC test, and the SN (SN curr2) generated when using ref -1 as the reference is different from SN curr1, an additional decompression attempt is performed based on SN curr2 as the decompressed SN.
- e. If the decompressed header generated in b. does not pass the CRC test and SN curr2 is the same as SN curr1, an additional decompression attempt is not useful and is not attempted.
- f. If the decompressed header generated in d. passes the CRC test, ref -1 is not changed while ref 0 is set to SN curr2.
- g. If the decompressed header generated in d. does not pass the CRC test, the decompressor acts according to rules a) through c) of section 5.3.2.2.3.

The purpose of this algorithm is to repair the context. If the header generated in d. passes the CRC test, the references are updated according to f., but two more headers MUST also be successfully decompressed before the repair is deemed successful. Of the three successful headers, the first two SHOULD be discarded and only the third delivered to upper layers. If decompression of any of the three headers fails, the decompressor MUST discard that header and the previously generated headers, and act according to rules a) through c) of section 5.3.2.2.3.

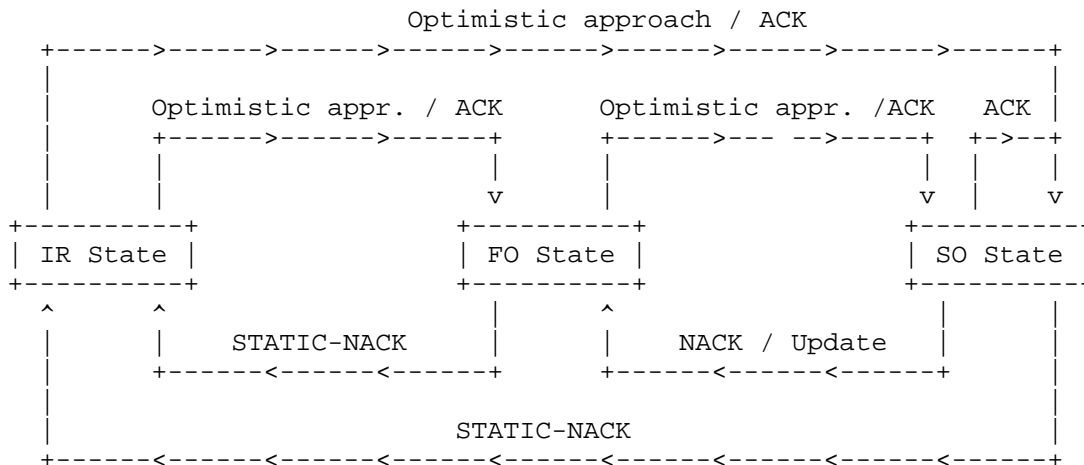
5.3.2.3. Feedback in Unidirectional mode

To improve performance for the Unidirectional mode over a link that does have a feedback channel, the decompressor MAY send an acknowledgment when decompression succeeds. Setting the mode parameter in the ACK packet to U indicates that the compressor is to stay in Unidirectional mode. When receiving an ACK(U), the compressor should reduce the frequency of IR packets since the static information has been correctly received, but it is not required to stop sending IR packets. If IR packets continue to arrive, the decompressor MAY repeat the ACK(U), but it SHOULD NOT repeat the ACK(U) continuously.

5.4. Operation in Bidirectional Optimistic mode

5.4.1. Compressor states and logic (O-mode)

Below is the state machine for the compressor in Bidirectional Optimistic mode. The details of each state, state transitions, and compression logic are given subsequent to the figure.



#### 5.4.1.1. State transition logic

The transition logic for compression states in Bidirectional Optimistic mode has much in common with the logic of the Unidirectional mode. The optimistic approach principle and transitions occasioned by the need for updates work in the same way as described in chapter 5.3.1. However, in Optimistic mode there are no timeouts. Instead, the Optimistic mode makes use of feedback from decompressor to compressor for transitions in the backward direction and for OPTIONAL improved forward transition.

##### 5.4.1.1.1. Negative acknowledgments (NACKs), downward transition

Negative acknowledgments (NACKs), also called context requests, obviate the periodic updates needed in Unidirectional mode. Upon reception of a NACK the compressor transits back to the FO state and sends updates (IR-DYN, UOR-2, or possibly IR) to the decompressor. NACKs carry the SN of the latest packet successfully decompressed, and this information MAY be used by the compressor to determine what fields need to be updated.

Similarly, reception of a STATIC-NACK packet makes the compressor transit back to the IR state.

##### 5.4.1.1.2. Optional acknowledgments, upwards transition

In addition to NACKs, positive feedback (ACKs) MAY also be used for UOR-2 packets in the Bidirectional Optimistic mode. Upon reception of an ACK for an updating packet, the compressor knows that the decompressor has received the acknowledged packet and the transition to a higher compression state can be carried out immediately. This functionality is optional, so a compressor MUST NOT expect to get such ACKs initially.

The compressor MAY use the following algorithm to determine when to expect ACKs for UOR-2 packets. Let an update event be when a sequence of UOR-2 headers are sent to communicate an irregularity in the packet stream. When ACKs have been received for  $k_3$  out of the last  $n_3$  update events, the compressor will expect ACKs. A compressor which expects ACKs will repeat updates (possibly not in every packet) until an ACK is received.

##### 5.4.1.2. Compression logic and packets used

The compression logic is the same for the Bidirectional Optimistic mode as for the Unidirectional mode (see section 5.3.1.2).

#### 5.4.2. Decompressor states and logic (O-mode)

The decompression states and the state transition logic are the same as for the Unidirectional case (see section 5.3.2). What differs is the decompression and feedback logic.

##### 5.4.2.1. Decompression logic, timer-based timestamp decompression

In Bidirectional mode (or if there is some other way for the compressor to obtain the decompressor's clock resolution and the link's jitter), timer-based timestamp decompression may be used to improve compression efficiency when RTP Timestamp values are proportional to wall-clock time. The mechanisms used are those described in 4.5.4.

##### 5.4.2.2. Feedback logic (O-mode)

The feedback logic defines what feedback to send due to different events when operating in the various states. As mentioned above, there are three principal kinds of feedback; ACK, NACK and STATIC-NACK. Further, the logic described below will refer to different kinds of packets that can be received by the decompressor; Initialization and Refresh (IR) packets, IR packets without static information (IR-DYN) and type 2 packets (UOR-2), or type 1 (UO-1) and type 0 packets (UO-0). A type 0 packet carries a packet header compressed according to a fixed pattern, while type 1, 2 and IR-DYN packets are used when this pattern is broken.

Below, rules are defined stating which feedback to use when. If the optional feedback is used once, the decompressor is REQUIRED to continue to send optional feedback for the lifetime of the packet stream.

#### State Actions

- NC:
- When an IR packet passes the CRC check, send an ACK(O).
  - When receiving a type 0, 1, 2 or IR-DYN packet, or an IR packet has failed the CRC check, send a STATIC-NACK(O), subject to the considerations at the beginning of section 5.7.6.
- SC:
- When an IR packet is correctly decompressed, send an ACK(O).
  - When a type 2 or an IR-DYN packet is correctly decompressed, optionally send an ACK(O).
  - When a type 0 or 1 packet is received, treat it as a mismatching CRC and use the logic of section 5.3.2.2.3 to decide if a NACK(O) should be sent.



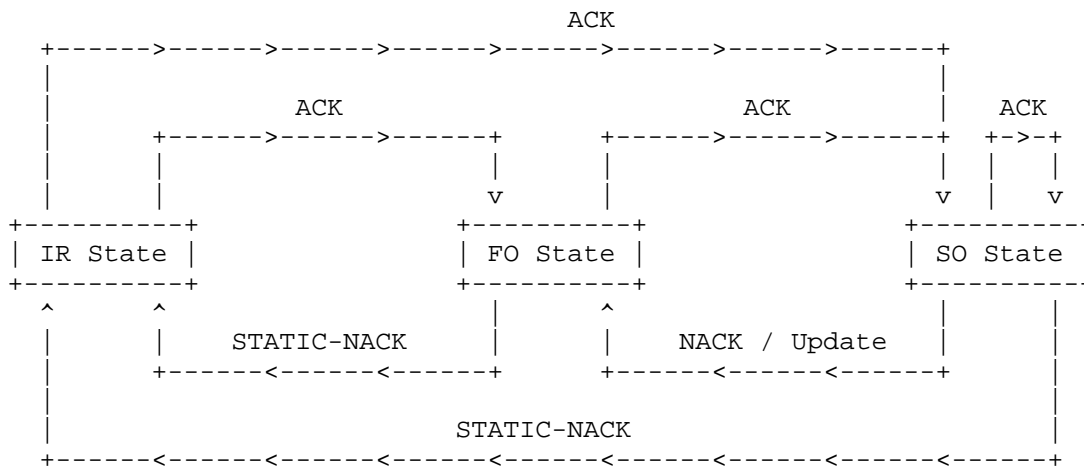
- When decompression of a type 2 packet, an IR-DYN packet or an IR packet has failed, use the logic of section 5.3.2.2.3 to decide if a STATIC-NACK(0) should be sent.

- FC:
- When an IR packet is correctly decompressed, send an ACK(0).
  - When a type 2 or an IR-DYN packet is correctly decompressed, optionally send an ACK(0).
  - When a type 0 or 1 packet is correctly decompressed, no feedback is sent.
  - When any packet fails the CRC check, use the logic of 5.3.2.2.3 to decide if a NACK(0) should be sent.

5.5. Operation in Bidirectional Reliable mode

5.5.1. Compressor states and logic (R-mode)

Below is the state machine for the compressor in Bidirectional Reliable mode. The details of each state, state transitions, and compression logic are given subsequent to the figure.



5.5.1.1. State transition logic (R-mode)

The transition logic for compression states in Reliable mode is based on three principles: the secure reference principle, the need for updates, and negative acknowledgments.

5.5.1.1.1. Upwards transition

The upwards transition is determined by the secure reference principle. The transition procedure is similar to the one described in section 5.3.1.1.1, with one important difference: the compressor

bases its confidence only on acknowledgments received from the decompressor. This ensures that the synchronization between the compression context and decompression context will never be lost due to packet losses.

#### 5.5.1.1.2. Downward transition

Downward transitions are triggered by the need for updates or by negative acknowledgment (NACKs and STATIC\_NACKs), as described in section 5.3.1.1.3 and 5.4.1.1.1, respectively. Note that NACKs should rarely occur in R-mode because of the secure reference used (see fourth paragraph of next section).

#### 5.5.1.2. Compression logic and packets used (R-mode)

The compressor starts in the IR state by sending IR packets. It transits to the FO state once it receives a valid ACK for an IR packet sent (an ACK can only be valid if it refers to an SN sent earlier). In the FO state, it sends the smallest packets that can communicate the changes, according to W-LSB or other encoding rules. Those packets could be of type R-1\*, UOR-2, or even IR-DYN.

The compressor will transit to the SO state after it has determined the presence of a string (see section 2), while also being confident that the decompressor has the string parameters. The confidence can be based on ACKs. For example, in a typical case where the string pattern has the form of non-SN-field = SN \* slope + offset, one ACK is enough if the slope has been previously established by the decompressor (i.e., only the new offset needs to be synchronized). Otherwise, two ACKs are required since the decompressor needs two headers to learn both the new slope and the new offset. In the SO state, R-0\* packets will be sent.

Note that a direct transition from the IR state to the SO state is possible.

The secure reference principle is enforced in both compression and decompression logic. The principle means that only a packet carrying a 7- or 8-bit CRC can update the decompression context and be used as a reference for subsequent decompression. Consequently, only field values of update packets need to be added to the encoding sliding windows (see 4.5) maintained by the compressor.

Reasons for the compressor to send update packets include:

- 1) The update may lead to a transition to higher compression efficiency (meaning either a higher compression state or smaller packets in the same state).

- 2) It is desirable to shrink sliding windows. Windows are only shrunk when an ACK is received.

The generation of a CRC is infrequent since it is only needed for an update packet.

One algorithm for sending update packets could be:

- \* Let pRTT be the number of packets that are sent during one round-trip time. In the SO state, when  $(64 - pRTT)$  headers have been sent since the last acked reference, the compressor will send  $m1$  consecutive R-0-CRC headers, then send  $(pRTT - m1)$  R-0 headers. After these headers have been sent, if the compressor has not received an ACK to at least one of the previously sent R0-CRC, it sends R-0-CRC headers continuously until it receives a corresponding ACK.  $m1$  is an implementation parameter, which can be as large as pRTT.
- \* In the FO state,  $m2$  UOR-2 headers are sent when there is a pattern change, after which the compressor sends  $(pRTT - m2)$  R-1-\* headers.  $m2$  is an implementation parameter, which can be as large as pRTT. At that time, if the compressor has not received enough ACKs to the previously sent UOR-2 packets in order to transit to SO state, it can repeat the cycle with the same  $m2$ , or repeat the cycle with a larger  $m2$ , or send UOR-2 headers continuously ( $m2 = pRTT$ ). The operation stops when the compressor has received enough ACKs to make the transition.

An algorithm for processing ACKs could be:

- \* Upon reception of an ACK, the compressor first derives the complete SN (see section 5.7.6.1). Then it searches the sliding window for an update packet that has the same SN. If found, that packet is the one being ACKed. Otherwise, the ACK is invalid and MUST be discarded.
- \* It is possible, although unlikely, that residual errors on the reverse channel could cause a packet to mimic a valid ACK feedback. The compressor may use a local clock to reduce the probability of processing such a mistaken ACK. After finding the update packet as described above, the compressor can check the time elapsed since the packet was sent. If the time is longer than RTT\_U, or shorter than RTT\_L, the compressor may choose to discard the ACK. RTT\_U and RTT\_L correspond to an upper bound and lower bound of the RTT, respectively. (These bounds should be chosen appropriately to allow some variation of RTT.) Note that the only side effect of discarding a good ACK is slightly reduced compression efficiency.

### 5.5.2. Decompressor states and logic (R-mode)

The decompression states and the state transition logic are the same as for the Unidirectional case (see section 5.3.2). What differs is the decompression and feedback logic.

#### 5.5.2.1. Decompression logic (R-mode)

The rules for when decompression is allowed are the same as for U-mode. Although the acking scheme in R-mode guarantees that non-decompressible packets are never sent by the compressor, residual errors can cause delivery of unexpected packets for which decompression should not be attempted.

Decompression MUST follow the secure reference principle as described in 5.5.1.2.

CRC verification is infrequent since only update packets carry CRCs. A CRC mismatch can only occur due to 1) residual bit errors in the current header, and/or 2) a damaged context due to residual bit errors in previous headers or feedback. Although it is impossible to determine which is the actual cause, case 1 is more likely, as a previous header reconstructed according to a damaged packet is unlikely to pass the 7- or 8-bit CRC, and damaged packets are unlikely to result in feedback that damages the context. The decompressor SHOULD act according to section 5.3.2.2.3 when CRCs fail, except that no local repair is performed. Note that all the parameter numbers, *k*<sub>1</sub>, *n*<sub>1</sub>, *k*<sub>2</sub>, and *n*<sub>2</sub>, are applied to the update packets only (i.e., exclude R-0, R-1\*).

#### 5.5.2.2. Feedback logic (R-mode)

The feedback logic for the Bidirectional Reliable mode is as follows:

- When an updating packet (i.e., a packet carrying a 7- or 8-bit CRC) is correctly decompressed, send an ACK(R), subject to the sparse ACK mechanism described below.
- When context damage is detected, send a NACK(R) if in Full Context state, or a STATIC-NACK(R) if in Static Context state.
- In No Context state, send a STATIC-NACK(R) when receiving a non-IR packet, subject to the considerations at the beginning of section 5.7.6. The decompressor SHOULD NOT send STATIC-NACK(R) when receiving an IR packet that fails the CRC check, as the compressor will stay in IR state and thus continue sending IR packets until a valid ACK is received (see section 5.5.1.2).

- Feedback is never sent for packets not updating the context (i.e., packets that do not carry a CRC)

A mechanism called "Sparse ACK" can be applied to reduce the feedback overhead caused by a large RTT. For a sequence of ACK-triggering events, a minimal set of ACKs MUST be sent:

- 1) For a sequence of R-0-CRC packets, the first one MUST be ACKed.
- 2) For a sequence of UOR-2, IR, or IR-DYN packets, the first N of them MUST be ACKed, where N is the number of ACKs needed to give the compressor confidence that the decompressor has acquired the new string parameters (see second paragraph of 5.5.1.2). In case the decompressor cannot determine the value of N, the default value 2 SHOULD be used. If the subsequently received packets continue the same change pattern of header fields, sparse ACK can be applied. Otherwise, each new pattern MUST be treated as a new sequence, i.e., the first N packets that exhibit a new pattern MUST be ACKed.

After sending these minimal ACKs, the decompressor MAY choose to ACK only k subsequent packets per RTT ("Sparse ACKs"), where k is an implementation parameter. To achieve robustness against loss of ACKs, k SHOULD be at least 1.

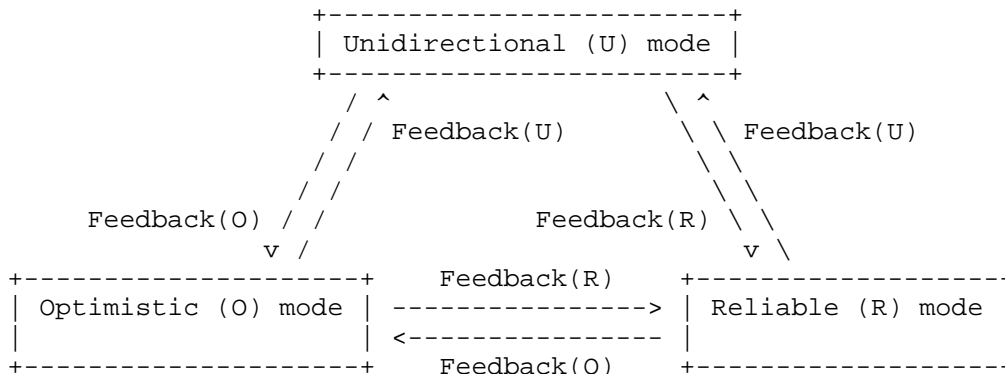
To avoid ambiguity at the compressor, the decompressor MUST use the feedback format whose SN field length is equal to or larger than the one in the compressed packet that triggered the feedback.

Context damage is detected according to the principles in 5.3.2.2.3.

When the decompressor is capable of timer-based compression of the RTP Timestamp (e.g., it has access to a clock with sufficient resolution, and the jitter introduced internally in the receiving node is sufficiently small) it SHOULD signal that it is ready to do timer-based compression of the RTP Timestamp. The compressor will then make a decision based on its knowledge of the channel and the observed properties of the packet stream.

## 5.6. Mode transitions

The decision to move from one compression mode to another is taken by the decompressor and the possible mode transitions are shown in the figure below. Subsequent chapters describe how the transitions are performed together with exceptions for the compression and decompression functionality during transitions.



5.6.1. Compression and decompression during mode transitions

The following sections assume that, for each context, the compressor and decompressor maintain a variable whose value is the current compression mode for that context. The value of the variable controls, for the context in question, which packet types to use, which actions to be taken, etc.

As a safeguard against residual errors, all feedback sent during a mode transition MUST be protected by a CRC, i.e., the CRC option MUST be used. A mode transition MUST NOT be initiated by feedback which is not protected by a CRC.

The subsequent subsections define exactly when to change the value of the MODE variable. When ROHC transits between compression modes, there are several cases where the behavior of compressor or decompressor must be restricted during the transition phase. These restrictions are defined by exception parameters that specify which restrictions to apply. The transition descriptions in subsequent chapters refer to these exception parameters and defines when they are set and to what values. All mode related parameters are listed below together with their possible values, with explanations and restrictions:

Parameters for the compressor side:

- C\_MODE:
  - Possible values for the C\_MODE parameter are (U)NIDIRECTIONAL, (O)PTIMISTIC and (R)ELIABLE. C\_MODE MUST be initialized to U.
- C\_TRANS:
  - Possible values for the C\_TRANS parameter are (P)ENDING and (D)ONE. C\_TRANS MUST be initialized to D. When C\_TRANS is P, it is REQUIRED











The following fields occur in several headers and extensions:

SN: The compressed RTP Sequence Number.

Compressed with W-LSB. The interpretation intervals, see section 4.5.1, are defined as follows:

$$\begin{aligned} p &= 1 && \text{if bits(SN) } \leq 4 \\ p &= 2^{(\text{bits(SN)}-5)} - 1 && \text{if bits(SN) } > 4 \end{aligned}$$

IP-ID: A compressed IP-ID field.

IP-ID fields in compressed base headers carry the compressed IP-ID of the innermost IPv4 header whose corresponding RND flag is not 1. The rules below assume that the IP-ID is for the innermost IP header. If it is for an outer IP header, the RND2 and NBO2 flags are used instead of RND and NBO.

If  $\text{value(RND)} = 0$ ,  $\text{hdr(IP-ID)}$  is compressed using Offset IP-ID encoding (see section 4.5.5) using  $p = 0$  and  $\text{default-slope(IP-ID offset)} = 0$ .

If  $\text{value(RND)} = 1$ , IP-ID is the uncompressed  $\text{hdr(IP-ID)}$ . IP-ID is then passed as additional octets at the end of the compressed header, after any extensions.

If  $\text{value(NBO)} = 0$ , the octets of  $\text{hdr(IP-ID)}$  are swapped before compression and after decompression. The value of NBO is ignored when  $\text{value(RND)} = 1$ .

TS: The compressed RTP Timestamp value.

If  $\text{value(TIME_STRIDE)} > 0$ , timer-based compression of the RTP Timestamp is used (see section 4.5.4).

If  $\text{value(Tsc)} = 1$ , Scaled RTP Timestamp encoding is used before compression (see section 4.5.3), and  $\text{default-slope(TS)} = 1$ .

If  $\text{value(Tsc)} = 0$ , the Timestamp value is compressed as-is, and  $\text{default-slope(TS)} = \text{value(TS\_STRIDE)}$ .

The interpretation intervals, see section 4.5.1, are defined as follows:

$$p = 2^{(\text{bits(TS)}-2)} - 1$$

CRC: The CRC over the original, uncompressed, header.

For 3-bit CRCs, the polynomial of section 5.9.2 is used.

For 7-bit CRCs, the polynomial of section 5.9.2 is used.

For 8-bit CRCs, the polynomial of section 5.9.1 is used.

M: RTP Marker bit.

Context(M) is initially zero and is never updated. value(M) = 1 only when field(M) = 1.

The general format for a compressed RTP header is as follows:

```

  0   1   2   3   4   5   6   7
  ---
  :           Add-CID octet           :   if for small CIDs and CID 1-15
+-----+-----+-----+-----+-----+
| first octet of base header | (with type indication)
+-----+-----+-----+-----+-----+
:                                     :
/   0, 1, or 2 octets of CID   /   1-2 octets if large CIDs
:                               :
+-----+-----+-----+-----+-----+
/ remainder of base header / variable number of bits
+-----+-----+-----+-----+-----+
:                               :
/ Extension (see 5.7.5) / extension, if X = 1 in base header
:                               :
:                               :
+ IP-ID of outer IPv4 header + 2 octets, if value(RND2) = 1
:                               :
+-----+-----+-----+-----+-----+
/ AH data for outer list / variable (see 5.8.4.2)
+-----+-----+-----+-----+-----+
:                               :
+ GRE checksum (see 5.8.4.4) + 2 octets, if GRE flag C = 1
:                               :
:                               :
+ IP-ID of inner IPv4 header + 2 octets, if value(RND) = 1
:                               :
+-----+-----+-----+-----+-----+
/ AH data for inner list / variable (see 5.8.4.2)
+-----+-----+-----+-----+-----+
:                               :
+ GRE checksum (see 5.8.4.4) + 2 octets, if GRE flag C = 1
:                               :
+-----+-----+-----+-----+-----+
:                               :
+           UDP Checksum           + 2 octets,
:                               :   if context(UDP Checksum) != 0
+-----+-----+-----+-----+-----+

```

Note that the order of the fields following the optional extension is the same as the order between the fields in an uncompressed header.

In subsequent sections, the position of the large CID in the diagrams is indicated using this notation:

+====+====+====+====+====+====+====+====+

Whether the UDP Checksum field is present or not is controlled by the value of the UDP Checksum in the context. If nonzero, the UDP Checksum is enabled and sent along with each packet. If zero, the UDP Checksum is disabled and not sent. Should `hdr(UDP Checksum)` be nonzero when `context(UDP Checksum)` is zero, the header cannot be compressed. It must be sent uncompressed or the context reinitialized using an IR packet. `Context(UDP Checksum)` is updated only by IR or IR-DYN headers, never by UDP checksums sent in headers of type 2, 1, or 0.

When an IPv4 header is present in the static context, for which the corresponding RND flag has not been established to be 1, the packet types R-1 and UO-1 MUST NOT be used.

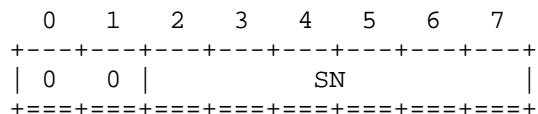
When no IPv4 header is present in the static context, or the RND flags for all IPv4 headers in the context have been established to be 1, the packet types R-1-ID, R-1-TS, UO-1-ID, and UO-1-TS MUST NOT be used.

While in the transient state in which an RND flag is being established, the packet types R-1-ID, R-1-TS, UO-1-ID, and UO-1-TS MUST NOT be used. This implies that the RND flag(s) of the Extension 3 may have to be inspected before the format of a base header carrying an Extension 3 can be determined.

5.7.1. Packet type 0: UO-0, R-0, R-0-CRC

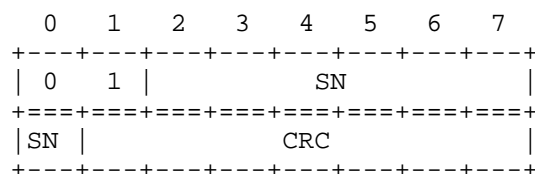
Packet type 0 is indicated by the first bit being 0:

R-0



Updating properties: R-0 packets do not update any part of the context.

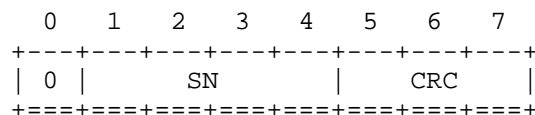
## R-0-CRC



Note: The SN field straddles the CID field.

Updating properties: R-0-CRC packets update context(RTP Sequence Number).

## UO-0

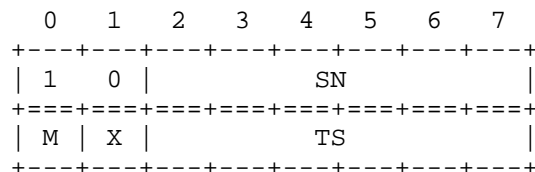


Updating properties: UO-0 packets update the current value of context(RTP Sequence Number).

## 5.7.2. Packet type 1 (R-mode): R-1, R-1-TS, R-1-ID

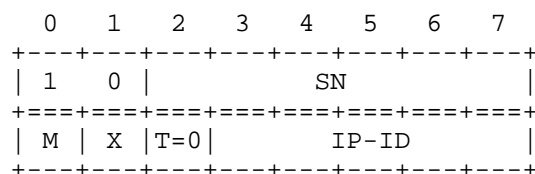
Packet type 1 is indicated by the first bits being 10:

## R-1



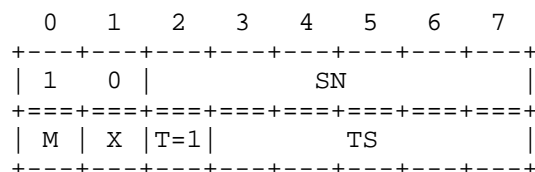
Note: R-1 cannot be used if the context contains at least one IPv4 header with value(RND) = 0. This disambiguates it from R-1-ID and R-1-TS.

## R-1-ID



Note: R-1-ID cannot be used if there is no IPv4 header in the context or if value(RND) and value(RND2) are both 1.

## R-1-TS



Note: R-1-TS cannot be used if there is no IPv4 header in the context or if value(RND) and value(RND2) are both 1.

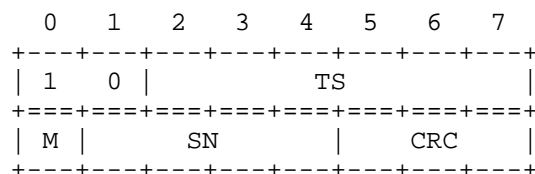
X: X = 0 indicates that no extension is present;  
X = 1 indicates that an extension is present.

T: T = 0 indicates format R-1-ID;  
T = 1 indicates format R-1-TS.

Updating properties: R-1\* headers do not update any part of the context.

## 5.7.3. Packet type 1 (U/O-mode): UO-1, UO-1-ID, UO-1-TS

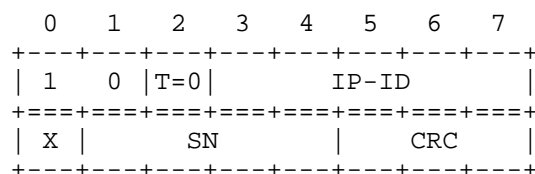
## UO-1



Note: UO-1 cannot be used if the context contains at least one IPv4 header with value(RND) = 0. This disambiguates it from UO-1-ID and UO-1-TS.

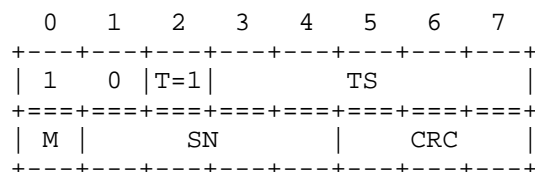


## UO-1-ID



Note: UO-1-ID cannot be used if there is no IPv4 header in the context or if value(RND) and value(RND2) are both 1.

## UO-1-TS



Note: UO-1-TS cannot be used if there is no IPv4 header in the context or if value(RND) and value(RND2) are both 1.

X: X = 0 indicates that no extension is present;  
X = 1 indicates that an extension is present.

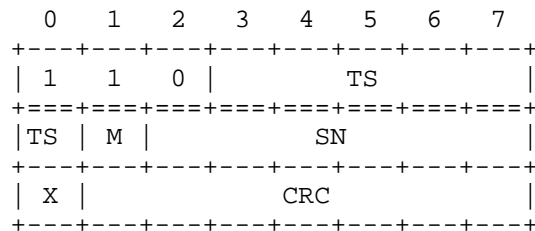
T: T = 0 indicates format UO-1-ID;  
T = 1 indicates format UO-1-TS.

Updating properties: UO-1\* packets update context(RTP Sequence Number). UO-1 and UO-1-TS packets update context(RTP Timestamp). UO-1-ID packets update context(IP-ID). Values provided in extensions, except those in other SN, TS, or IP-ID fields, do not update the context.

## 5.7.4. Packet type 2: UOR-2

Packet type 2 is indicated by the first bits being 110:

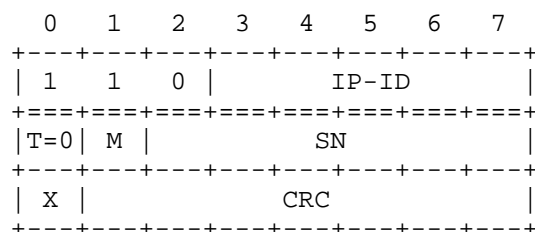
UOR-2



Note: UOR-2 cannot be used if the context contains at least one IPv4 header with value(RND) = 0. This disambiguates it from UOR-2-ID and UOR-2-TS.

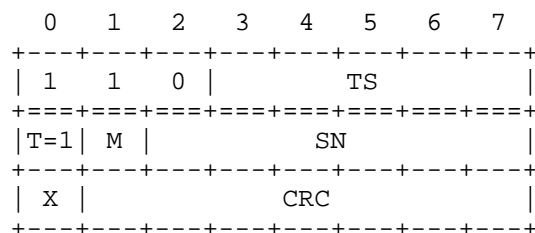
Note: The TS field straddles the CID field.

UOR-2-ID



Note: UOR-2-ID cannot be used if there is no IPv4 header in the context or if value(RND) and value(RND2) are both 1.

UOR-2-TS



Note: UOR-2-TS cannot be used if there is no IPv4 header in the context or if value(RND) and value(RND2) are both 1.

X: X = 0 indicates that no extension is present;  
X = 1 indicates that an extension is present.

T: T = 0 indicates format UOR-2-ID;  
T = 1 indicates format UOR-2-TS.

Updating properties: All values provided in UOR-2\* packets update the context, unless explicitly stated otherwise.

5.7.5. Extension formats

(Note: the term extension as used for additional information contained in the ROHC headers does not bear any relationship to the term extension header used in IP.)

Fields in extensions are concatenated with the corresponding field in the base compressed header, if there is one. Bits in an extension are less significant than bits in the base compressed header (see section 4.5.7).

The TS field is scaled in all extensions, as it is in the base header, except optionally when using Extension 3 where the Tsc flag can indicate that the TS field is not scaled. Value(TS\_STRIDE) is used as the scale factor when scaling the TS field.

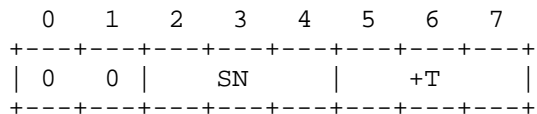
In the following three extensions, the interpretation of the fields depends on whether there is a T-bit in the base compressed header, and if so, on the value of that field. When there is no T-bit, +T and -T both mean TS. This is the case when there are no IPv4 headers in the static context, and when all IPv4 headers in the static context have their corresponding RND flag set (i.e., RND = 1).

If there is a T-bit,

T = 1 indicates that +T is TS, and  
-T is IP-ID;

T = 0 indicates that +T is IP-ID, and  
-T is TS.

Extension 0:





Extension 3:

0	1	2	3	4	5	6	7	
1	1	S	R-TS	Tsc	I	ip	rtp	(FLAGS)
Inner IP header flags							ip2	if ip = 1
Outer IP header flags								if ip2 = 1
SN								if S = 1
TS (encoded as in section 4.5.6)								1-4 octets, if R-TS = 1
Inner IP header fields								variable, if ip = 1
IP-ID								2 octets, if I = 1
Outer IP header fields								variable, if ip2 = 1
RTP header flags and fields								variable, if rtp = 1

S, R-TS, I, ip, rtp, ip2: Indicate presence of fields as shown to the right of each field above.

Tsc: Tsc = 0 indicates that TS is not scaled;  
 Tsc = 1 indicates that TS is scaled according to section 4.5.3, using value(TS\_STRIDE).  
 Context(Tsc) is always 1. If scaling is not desired, the compressor will establish TS\_STRIDE = 1.

SN: See the beginning of section 5.7.

TS: Variable number of bits of TS, encoded according to section 4.5.6. See the beginning of section 5.7.

IP-ID: See the beginning of section 5.7.

Inner IP header flags

These correspond to the inner IP header if there are two, and the single IP header otherwise.

0	1	2	3	4	5	6	7	
.....	.....	.....	.....	.....	.....	.....	.....	.....
TOS	TTL	DF	PR	IPX	NBO	RND	ip2	if ip = 1
.....	.....	.....	.....	.....	.....	.....	.....	.....

TOS, TTL, PR, IPX: Indicates presence of fields as shown to the right of the field in question below.

DF: Don't Fragment bit of IP header.

NBO: Indicates whether the octets of hdr(IP identifier) of this IP header are swapped before compression and after decompression.

NBO = 1 indicates that the octets need not be swapped. NBO = 0 indicates that the octets are to be swapped. See section 4.5.5.

RND: Indicates whether hdr(IP identifier) is not to be compressed but instead sent as-is in compressed headers.

IP2: Indicates presence of Outer IP header fields. Unless the static context contains two IP headers, IP2 is always zero.

Inner IP header fields

.....	.....	.....	.....	.....	.....	.....	.....	.....
	Type of Service/Traffic Class		if TOS = 1					
.....	.....	.....	.....	.....	.....	.....	.....	.....
	Time to Live/Hop Limit		if TTL = 1					
.....	.....	.....	.....	.....	.....	.....	.....	.....
	Protocol/Next Header		if PR = 1					
.....	.....	.....	.....	.....	.....	.....	.....	.....
/	IP extension headers	/	variable,					
.....	.....	.....	if IPX = 1					

Type of Service/Traffic Class: That field in the uncompressed IP header (absolute value).

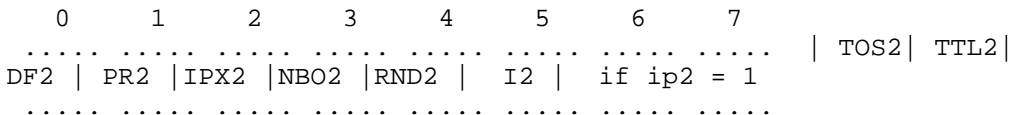
Time to Live/Hop Limit: That field in the uncompressed IP header.

Protocol/Next Header: That field in the uncompressed IP header.

IP extension header(s): According to section 5.8.5.

Outer IP header flags

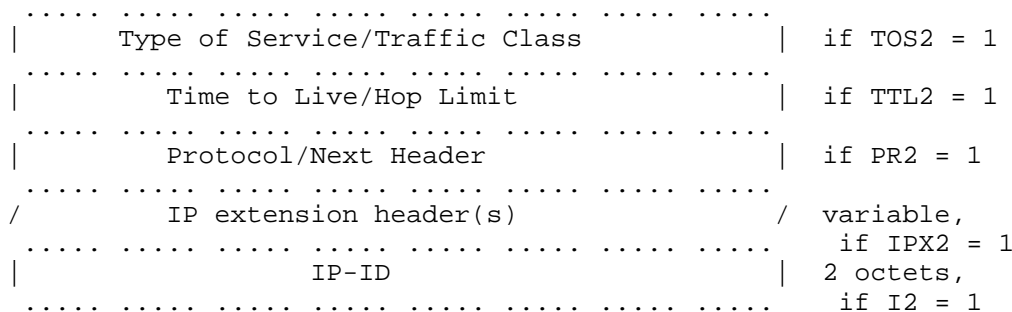
The fields in this part of the Extension 3 header refer to the outermost IP header:



These flags are the same as the Inner IP header flags, but refer to the outer IP header instead of the inner IP header. The following flag, however, has no counterpart in the Inner IP header flags:

I2: Indicates presence of the IP-ID field.

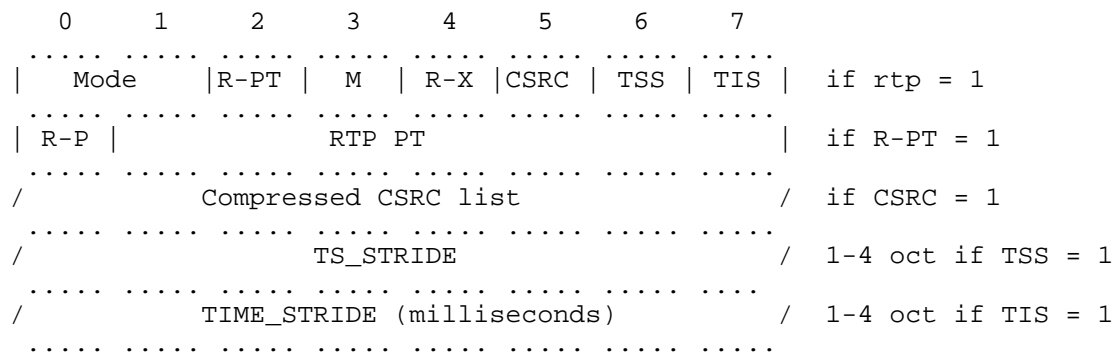
Outer IP header fields



The fields in this part of Extension 3 are as for the Inner IP header fields, but they refer to the outer IP header instead of the inner IP header. The following field, however, has no counterpart among the Inner IP header fields:

IP-ID: The IP Identifier field of the outer IP header, unless the inner header is an IPv6 header, in which case I2 is always zero.

RTP header flags and fields



Mode: Compression mode. 0 = Reserved,  
 1 = Unidirectional,  
 2 = Bidirectional Optimistic,  
 3 = Bidirectional Reliable.

R-PT, CSRC, TSS, TIS: Indicate presence of fields as shown to the right of each field above.

R-P: RTP Padding bit, absolute value (presumed zero if absent).

R-X: RTP eXtension bit, absolute value.

M: See the beginning of section 5.7.

RTP PT: Absolute value of RTP Payload type field.

Compressed CSRC list: See section 5.8.1.

TS\_STRIDE: Predicted increment/decrement of the RTP Timestamp field when it changes. Encoded as in section 4.5.6.

TIME\_STRIDE: Predicted time interval in milliseconds between changes in the RTP Timestamp. Also an indication that the compressor desires to perform timer-based compression of the RTP Timestamp field: see section 4.5.4. Encoded as in section 4.5.6.

5.7.5.1. RND flags and packet types

The values of the RND and RND2 flags are changed by sending UOR-2 headers with Extension 3, or IR-DYN headers, where the flag(s) have their new values. The establishment procedure of the flags is the normal one for the current mode, i.e., in U-mode and O-mode the values are repeated several times to ensure that the decompressor



receives at least one. In R-mode, the flags are sent until an acknowledgment for a packet with the new RND flag values is received.

The decompressor updates the values of its RND and RND2 flags whenever it receives an UOR-2 with Extension 3 carrying values for RND or RND2, and the UOR-2 CRC verifies successful decompression.

When an IPv4 header for which the corresponding RND flag has not been established to be 1 is present in the static context, the packet types R-1 and UO-1 MUST NOT be used.

When no IPv4 header is present in the static context, or the RND flags for all IPv4 headers in the context have been established to be 1, the packet types R-1-ID, R-1-TS, UO-1-ID, and UO-1-TS MUST NOT be used.

While in the transient state in which an RND flag is being established, the packet types R-1-ID, R-1-TS, UO-1-ID, and UO-1-TS MUST NOT be used. This implies that the RND flag(s) of Extension 3 may have to be inspected before the exact format of a base header carrying an Extension 3 can be determined, i.e., whether a T-bit is present or not.

#### 5.7.5.2. Flags/Fields in context

Some flags and fields in Extension 3 need to be maintained in the context of the decompressor. Their values are established using the mechanism appropriate to the compression mode, unless otherwise indicated in the table below and in referred sections.

Flag/Field	Initial value	Comment
Mode	Unidirectional	See section 5.6
NBO	1	See section 4.5.5
RND	0	See sections 4.5.5, 5.7.5.1
NBO2	1	As NBO, but for outer header
RND2	0	As RND, but for outer header
TS_STRIDE	1	See section 4.5.3
TIME_STRIDE	0	See section 4.5.4
Tsc	1	Tsc is always 1 in context; can be 0 only when an Extension 3 is present. See the discussion of the TS field in the beginning of section 5.7.

#### 5.7.6. Feedback packets and formats

When the round-trip time between compressor and decompressor is large, several packets can be in flight concurrently. Therefore, several packets may be received by the decompressor after feedback has been sent and before the compressor has reacted to feedback. Moreover, decompression may fail due to residual errors in the compressed header.

Therefore,

- a) in 0-mode, the decompressor SHOULD limit the rate at which feedback on successful decompression is sent (if it is sent at all);
- b) when decompression fails, feedback SHOULD be sent only when decompression of several consecutive packets has failed, and when this occurs, the feedback rate SHOULD be limited;
- c) when packets are received which belong to a rejected packet stream, the feedback rate SHOULD be limited.

A decompressor MAY limit the feedback rate by sending feedback only for one out of every  $k$  packets provoking the same (kind of) feedback. The appropriate value of  $k$  is implementation dependent;  $k$  might be chosen such that feedback is sent 1-3 times per link round-trip time.

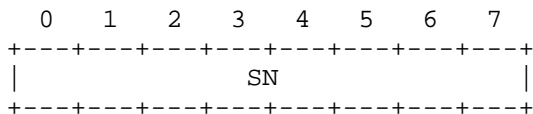
See section 5.2.2 for a discussion concerning ways to provide feedback information to the compressor.

##### 5.7.6.1. Feedback formats for ROHC RTP

This section describes the format for feedback information in ROHC RTP. See also 5.2.2.

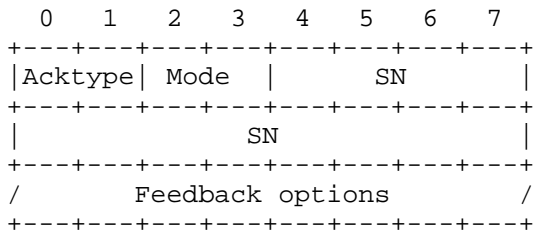
Several feedback formats carry a field labeled SN. The SN field contains LSBs of an RTP Sequence Number. The sequence number to use is the sequence number of the header which caused the feedback information to be sent. If that sequence number cannot be determined, for example when decompression fails, the sequence number to use is that of the last successfully decompressed header. If no sequence number is available, the feedback MUST carry a SN-NOT-VALID option. Upon reception, the compressor matches valid SN LSBs with the most recent header sent with a SN with matching LSBs. The decompressor must ensure that it sends enough SN LSBs in its feedback that this correlation does not become ambiguous; e.g., if an 8-bit SN LSB field could wrap around within a round-trip time, the FEEDBACK-1 format cannot be used.

FEEDBACK-1



A FEEDBACK-1 is an ACK. In order to send a NACK or a STATIC-NACK, FEEDBACK-2 must be used. FEEDBACK-1 does not contain any mode information; FEEDBACK-2 must be used when mode information is required.

FEEDBACK-2



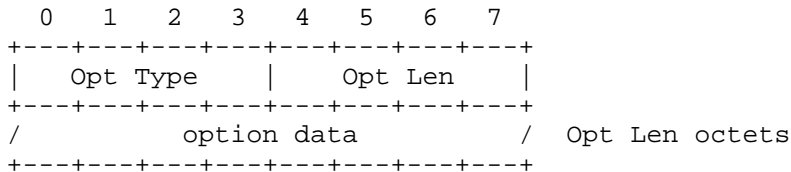
Acktype: 0 = ACK  
 1 = NACK  
 2 = STATIC-NACK  
 3 is reserved (MUST NOT be used for parseability)

Mode: 0 is reserved  
 1 = Unidirectional mode  
 2 = Bidirectional Optimistic mode  
 3 = Bidirectional Reliable mode

Feedback options: A variable number of feedback options, see section 5.7.6.2. Options may appear in any order.

5.7.6.2. ROHC RTP Feedback options

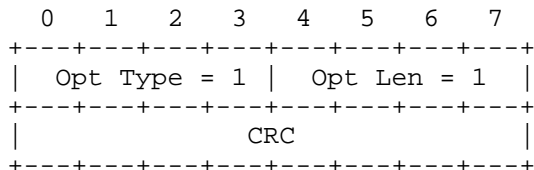
A ROHC RTP Feedback option has variable length and the following general format:



Sections 5.7.6.3-9 describe the currently defined ROHC RTP feedback options.

5.7.6.3. The CRC option

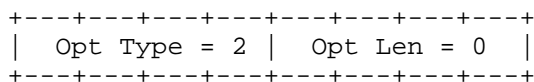
The CRC option contains an 8-bit CRC computed over the entire feedback payload, without the packet type and code octet, but including any CID fields, using the polynomial of section 5.9.1. If the CID is given with an Add-CID octet, the Add-CID octet immediately precedes the FEEDBACK-1 or FEEDBACK-2 format. For purposes of computing the CRC, the CRC fields of all CRC options are zero.



When receiving feedback information with a CRC option, the compressor MUST verify the information by computing the CRC and comparing the result with the CRC carried in the CRC option. If the two are not identical, the feedback information MUST be ignored.

5.7.6.4. The REJECT option

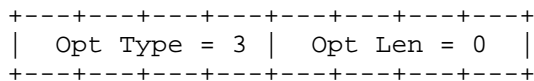
The REJECT option informs the compressor that the decompressor does not have sufficient resources to handle the flow.



When receiving a REJECT option, the compressor stops compressing the packet stream, and should refrain from attempting to increase the number of compressed packet streams for some time. Any FEEDBACK packet carrying a REJECT option MUST also carry a CRC option.

5.7.6.5. The SN-NOT-VALID option

The SN-NOT-VALID option indicates that the SN of the feedback is not valid. A compressor MUST NOT use the SN of the feedback to find the corresponding sent header when this option is present.



## 5.7.6.6. The SN option

The SN option provides 8 additional bits of SN.

```

+---+---+---+---+---+---+---+
| Opt Type = 4 | Opt Len = 1 |
+---+---+---+---+---+---+---+
|                SN                |
+---+---+---+---+---+---+---+

```

## 5.7.6.7. The CLOCK option

The CLOCK option informs the compressor of the clock resolution of the decompressor. This is needed to allow the compressor to estimate the jitter introduced by the clock of the decompressor when doing timer-based compression of the RTP Timestamp.

```

+---+---+---+---+---+---+---+
| Opt Type = 5 | Opt Len = 1 |
+---+---+---+---+---+---+---+
|   clock resolution (ms)   |
+---+---+---+---+---+---+---+

```

The smallest clock resolution which can be indicated is 1 millisecond. The value zero has a special meaning: it indicates that the decompressor cannot do timer-based compression of the RTP Timestamp. Any FEEDBACK packet carrying a CLOCK option SHOULD also carry a CRC option.

## 5.7.6.8. The JITTER option

The JITTER option allows the decompressor to report the maximum jitter it has observed lately, using the following formula which is very similar to the formula for Max\_Jitter\_BC in section 4.5.4.

Let observation window *i* contain the decompressor's best approximation of the sliding window of the compressor (see section 4.5.4) when header *i* is received.

$$\text{Max\_Jitter}_i =$$

$$\max \{ |(T_i - T_j) - ((a_i - a_j) / \text{TIME\_STRIDE})|, \\ \text{for all headers } j \text{ in observation window } i \}$$

$$\text{Max\_Jitter} =$$

$$\max \{ \text{Max\_Jitter}_i, \text{ for a large number of recent headers } i \}$$

This information may be used by the compressor to refine the formula for determining k when doing timer-based compression of the RTP Timestamp.

```

+-----+-----+-----+-----+-----+
| Opt Type = 6 | Opt Len = 1 |
+-----+-----+-----+-----+
|               Max_Jitter |
+-----+-----+-----+-----+

```

The decompressor MAY ignore the oldest observed values of Max\_Jitter\_i. Thus, the reported Max\_Jitter may decrease. Robustness will be reduced if the compressor uses a jitter estimate which is too small. Therefore, a FEEDBACK packet carrying a JITTER option SHOULD also carry a CRC option. Moreover, the compressor MAY ignore decreasing Max\_Jitter values.

5.7.6.9. The LOSS option

The LOSS option allows the decompressor to report the largest observed number of packets lost in sequence. This information MAY be used by the compressor to adjust the size of the reference window used in U- and O-mode.

```

+-----+-----+-----+-----+-----+
| Opt Type = 7 | Opt Len = 1 |
+-----+-----+-----+-----+
| longest loss event (packets) |
+-----+-----+-----+-----+

```

The decompressor MAY choose to ignore the oldest loss events. Thus, the value reported may decrease. Since setting the reference window too small can reduce robustness, a FEEDBACK packet carrying a LOSS option SHOULD also carry a CRC option. The compressor MAY choose to ignore decreasing loss values.

5.7.6.10. Unknown option types

If an option type unknown to the compressor is encountered, it must continue parsing the rest of the FEEDBACK packet, which is possible since the length of the option is explicit, but MUST otherwise ignore the unknown option.

5.7.6.11. RTP feedback example

Feedback for CID 8 indicating an ACK for SN 17 and Bidirectional Reliable mode can have the following formats.

Assuming small CIDs:

0	1	2	3	4	5	6	7	
1	1	1	1	0	0	1	1	feedback packet type, Code = 3
1	1	1	0	1	0	0	0	Add-CID octet with CID = 8
0	0	1	1	SN MSB = 0			AckType = ACK, Mode = Reliable	
SN LSB = 17								

The second, third, and fourth octet are handed to the compressor.

The FEEDBACK-1 format may also be used. Assuming large CIDs:

0	1	2	3	4	5	6	7	
1	1	1	1	0	0	1	0	feedback packet type, Code = 2
0	0	0	0	1	0	0	0	large CID with value 8
SN LSB = 17								

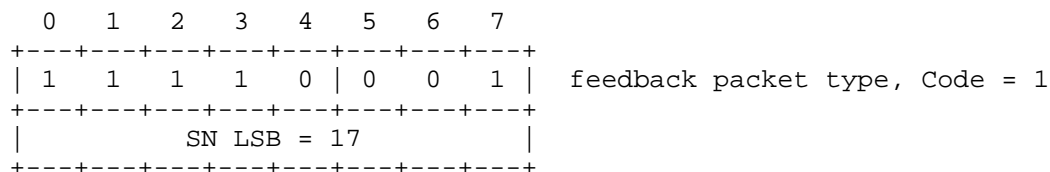
The second and third octet are handed to the compressor.

Assuming small CIDs:

0	1	2	3	4	5	6	7	
1	1	1	1	0	0	1	0	feedback packet type, Code = 2
1	1	1	0	1	0	0	0	Add-CID octet with CID = 8
SN LSB = 17								

The second and third octet are handed to the compressor.

Assuming small CIDs and CID 0 instead of CID 8:



The second octet is handed to the compressor.

5.7.7. RTP IR and IR-DYN packets

The subheaders which are compressible are split into a STATIC part and a DYNAMIC part. These parts are defined in sections 5.7.7.3 through 5.7.7.7.

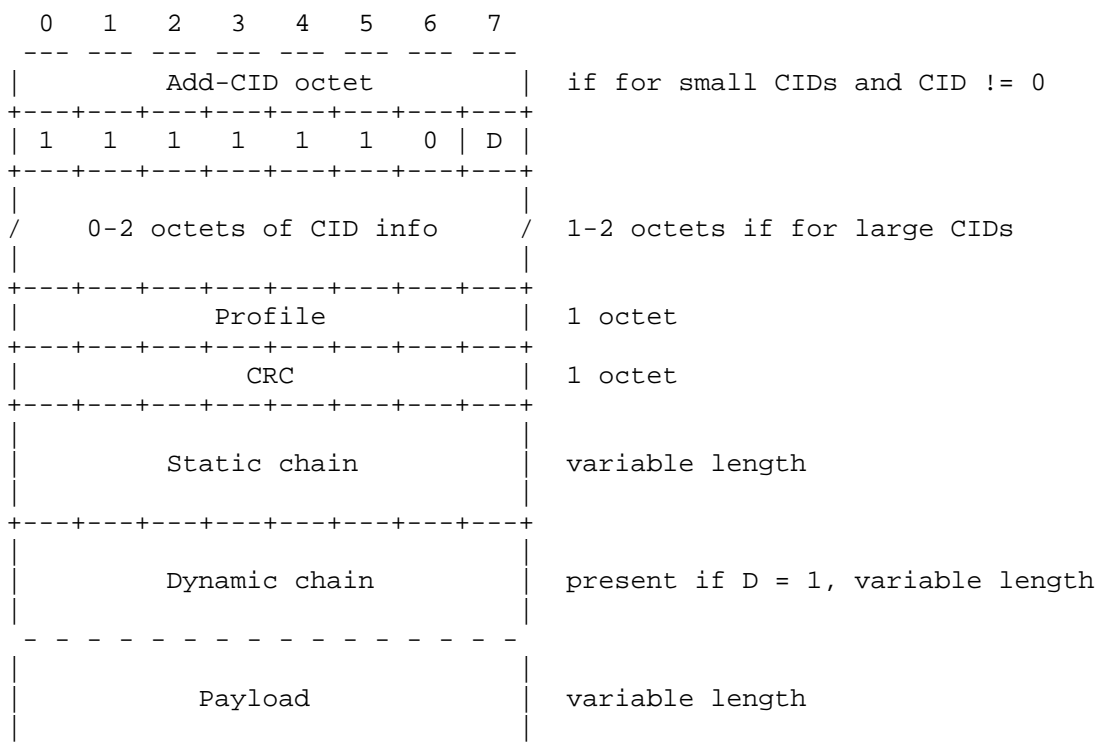
The structure of a chain of subheaders is determined by each header having a Next Header, or Protocol, field. This field identifies the type of the following header. Each Static part below that is followed by another Static part contains the Next Header/Protocol field and allows parsing of the Static chain; the Dynamic chain, if present, is structured analogously.

IR and IR-DYN packets will cause a packet to be delivered to upper layers if and only if the payload is non-empty. This means that an IP/UDP/RTP packet where the UDP length indicates a UDP payload of size 12 octets cannot be represented by an IR or IR-DYN packet. Such packets can instead be represented using the UNCOMPRESSED profile (section 5.10).

5.7.7.1. Basic structure of the IR packet

This packet type communicates the static part of the context, i.e., the values of the constant SN functions. It can optionally also communicate the dynamic part of the context, i.e., the parameters of nonconstant SN functions. It can also optionally communicate the payload of an original packet, if any.





D: D = 1 indicates that the dynamic chain is present.

Profile: Profile identifier, abbreviated as defined in section 5.2.3.

CRC: 8-bit CRC, computed according to section 5.9.1.

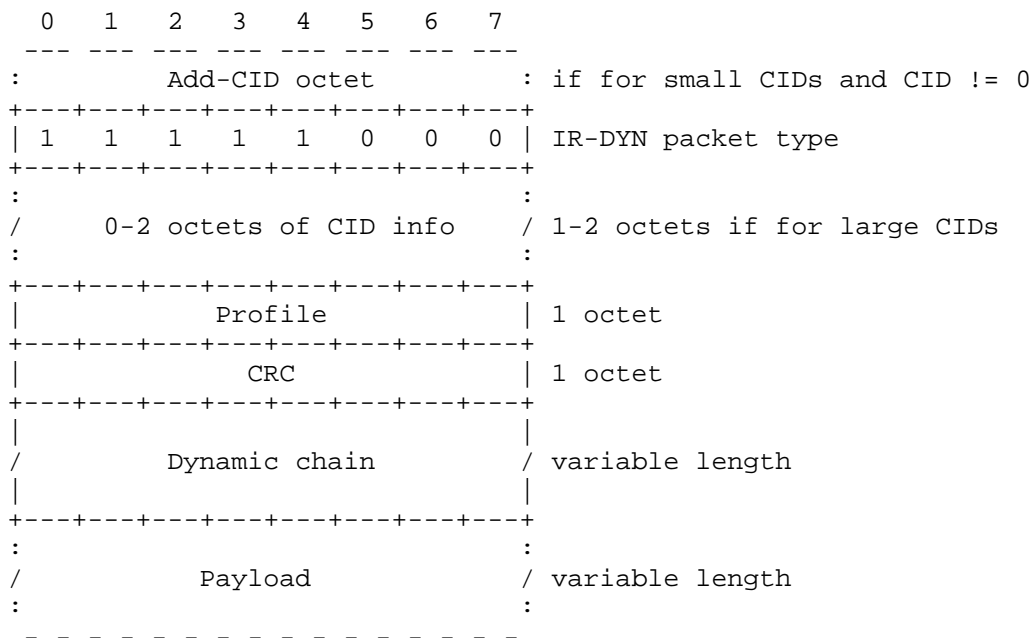
Static chain: A chain of static subheader information.

Dynamic chain: A chain of dynamic subheader information. What dynamic information is present is inferred from the Static chain.

Payload: The payload of the corresponding original packet, if any. The presence of a payload is inferred from the packet length.

5.7.7.2. Basic structure of the IR-DYN packet

This packet type communicates the dynamic part of the context, i.e., the parameters of nonconstant SN functions.



Profile: Profile identifier, abbreviated as defined in section 5.2.3.

CRC: 8-bit CRC, computed according to section 5.9.1.

NOTE: As the CRC checks only the integrity of the header itself, an acknowledgment of this header does not signify that previous changes to the static chain in the context are also acknowledged. In particular, care should be taken when IR packets that update an existing context are followed by IR-DYN packets.

Dynamic chain: A chain of dynamic subheader information. What dynamic information is present is inferred from the Static chain of the context.

Payload: The payload of the corresponding original packet, if any. The presence of a payload is inferred from the packet length.

Note: The static and dynamic chains of IR or IR-DYN packets for profile 0x0001 (ROHC RTP) MUST end with the static and dynamic parts of an RTP header. If not, the packet MUST be discarded and the context MUST NOT be updated.

Note: The static or dynamic chains of IR or IR-DYN packets for profile 0x0002 (ROHC UDP) MUST end with the static and dynamic parts of a UDP header. If not, the packet MUST be discarded and the context MUST NOT be updated.

Note: The static or dynamic chains of IR or IR-DYN packets for profile 0x0003 (ROHC ESP) MUST end with the static and dynamic parts of an ESP header. If not, the packet MUST be discarded and the context MUST NOT be updated.

5.7.7.3. Initialization of IPv6 Header [IPv6]

Static part:

```

+-----+-----+-----+-----+-----+-----+
| Version = 6 |Flow Label(msb)| 1 octet
+-----+-----+-----+-----+-----+-----+
/          Flow Label (lsb)          / 2 octets
+-----+-----+-----+-----+-----+-----+
|                Next Header                | 1 octet
+-----+-----+-----+-----+-----+-----+
/          Source Address          / 16 octets
+-----+-----+-----+-----+-----+-----+
/          Destination Address      / 16 octets
+-----+-----+-----+-----+-----+-----+

```

Dynamic part:

```

+-----+-----+-----+-----+-----+-----+
|                Traffic Class                | 1 octet
+-----+-----+-----+-----+-----+-----+
|                Hop Limit                    | 1 octet
+-----+-----+-----+-----+-----+-----+
/ Generic extension header list / variable length
+-----+-----+-----+-----+-----+-----+

```

Eliminated:

Payload Length

Extras:

Generic extension header list: Encoded according to section 5.8.6.1, with all header items present in uncompressed form.

CRC-DYNAMIC: Payload Length field (octets 5-6).

CRC-STATIC: All other fields (octets 1-4, 7-40).

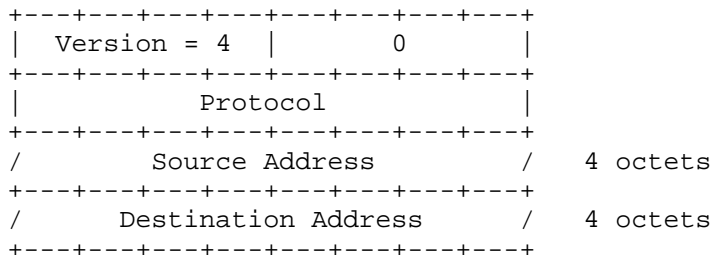
CRC coverage for extension headers is defined in section 5.8.7.

Note: The Next Header field indicates the type of the following header in the static chain, rather than being a copy of the Next Header field of the original IPv6 header. See also section 5.7.7.8.

5.7.7.4. Initialization of IPv4 Header [IPv4, section 3.1].

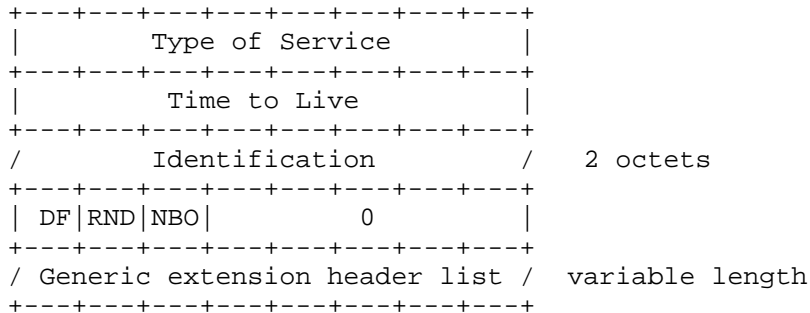
Static part:

Version, Protocol, Source Address, Destination Address.



Dynamic part:

Type of Service, Time to Live, Identification, DF, RND, NBO, extension header list.



Eliminated:

IHL	(IP Header Length, must be 5)
Total Length	(inferred in decompressed packets)
MF flag	(More Fragments flag, must be 0)
Fragment Offset	(must be 0)
Header Checksum	(inferred in decompressed packets)
Options, Padding	(must not be present)

Extras:

RND, NBO                    See section 5.7.

Generic extension header list: Encoded according to section 5.8.6.1, with all header items present in uncompressed form.

CRC-DYNAMIC: Total Length, Identification, Header Checksum (octets 3-4, 5-6, 11-12).

CRC-STATIC: All other fields (octets 1-2, 7-10, 13-20)

CRC coverage for extension headers is defined in section 5.8.7.

Note: The Protocol field indicates the type of the following header in the static chain, rather than being a copy of the Protocol field of the original IPv4 header. See also section 5.7.7.8.

5.7.7.5. Initialization of UDP Header [RFC-768].

Static part:

```

+-----+-----+-----+-----+-----+
/           Source Port           /   2 octets
+-----+-----+-----+-----+-----+
/           Destination Port       /   2 octets
+-----+-----+-----+-----+-----+

```

Dynamic part:

```

+-----+-----+-----+-----+-----+
/           Checksum               /   2 octets
+-----+-----+-----+-----+-----+

```

Eliminated:

Length

The Length field of the UDP header MUST match the Length field(s) of the preceding subheaders, i.e., there must not be any padding after the UDP payload that is covered by the IP Length.

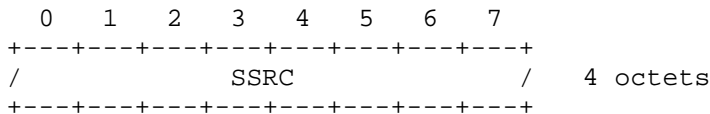
CRC-DYNAMIC: Length field, Checksum (octets 5-8).

CRC-STATIC: All other fields (octets 1-4).

5.7.7.6. Initialization of RTP Header [RTP].

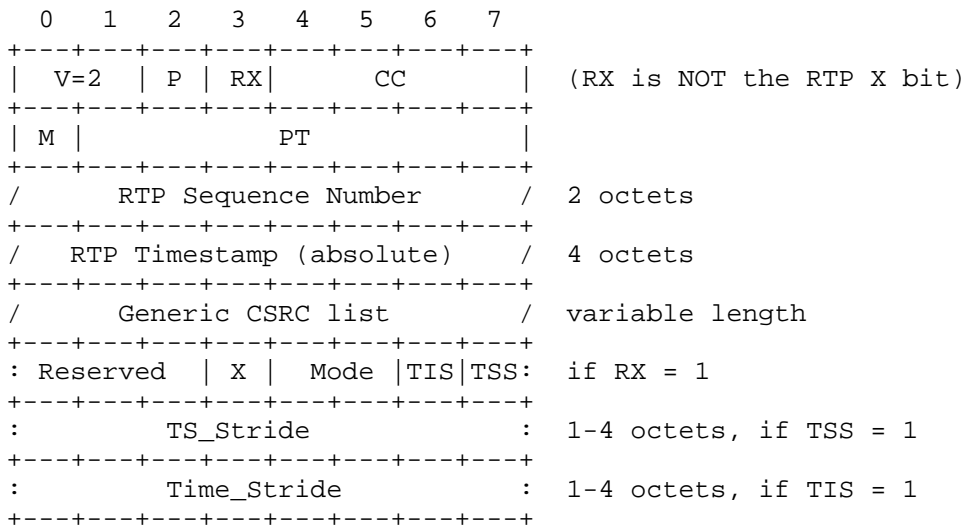
Static part:

SSRC.



Dynamic part:

P, X, CC, PT, M, sequence number, timestamp, timestamp stride, CSRC identifiers.



Eliminated:

Nothing.

Extras:

RX: Controls presence of extension.

Mode: Compression mode. 0 = Reserved,
1 = Unidirectional,
2 = Bidirectional Optimistic,
3 = Bidirectional Reliable.

X: Copy of X bit from RTP header (presumed 0 if RX = 0)

Reserved: Set to zero when sending, ignored when received.

Generic CSRC list: CSRC list encoded according to section
5.8.6.1, with all CSRC items present.

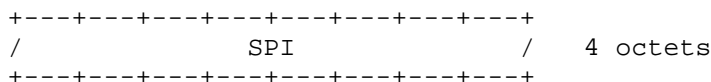
CRC-DYNAMIC: Octets containing M-bit, sequence number field,
and timestamp (octets 2-8).

CRC-STATIC: All other fields (octets 1, 9-12, original CSRC list).

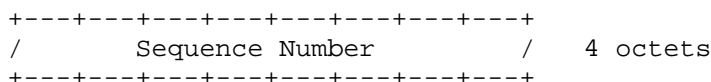
5.7.7.7. Initialization of ESP Header [ESP, section 2]

This is for the case when the NULL encryption algorithm [NULL] is NOT
being used with ESP, so that subheaders after the ESP header are
encrypted (see 5.12). See 5.8.4.3 for compression of the ESP header
when NULL encryption is being used.

Static part:



Dynamic part:



Eliminated:

Other fields are encrypted, and can neither be located nor
compressed.

CRC-DYNAMIC: Sequence number (octets 5-8)

CRC-STATIC: All other octets.

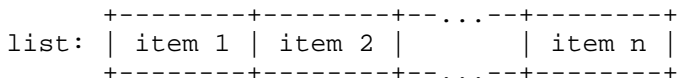
Note: No encrypted data is considered to be part of the header for purposes of computing the CRC, i.e., octets after the eight octet are not considered part of the header.

5.7.7.8. Initialization of Other Headers

Headers not explicitly listed in previous subsections can be compressed only by making them part of an extension header chain following an IPv4 or IPv6 header, see section 5.8.

5.8. List compression

Header information from the packet stream to be compressed can be structured as an ordered list, which is largely constant between packets. The generic structure of such a list is as follows.



This section describes the compression scheme for such information. The basic principles of list-based compression are the following:

- 1) While the list is constant, no information about the list is sent in compressed headers.
- 2) Small changes in the list are represented as additions (Insertion scheme), or deletions (Removal scheme), or both (Remove Then Insert scheme).
- 3) The list can also be sent in its entirety (Generic scheme).

There are two kinds of lists: CSRC lists in RTP packets, and extension header chains in IP packets (both IPv4 and IPv6).

IPv6 base headers and IPv4 headers cannot be part of an extension header chain. Headers which can be part of extension header chains include

- a) the AH header
- b) the null ESP header
- c) the minimal encapsulation header [RFC2004, section 3.1]
- d) the GRE header [GRE1, GRE2]
- e) IPv6 extension headers.



The table-based item compression scheme (5.8.1), which reduces the size of each item, is described first. Then it is defined which reference list to use in the insertion and removal schemes (5.8.2). List encoding schemes are described in section 5.8.3, and a few special cases in section 5.8.4. Finally, exact formats are described in sections 5.8.5-5.8.6.

5.8.1. Table-based item compression

The Table-based item compression scheme is a way to compress individual items sent in compressed lists. The compressor assigns each item in a list a unique identifier Index. The compressor conceptually maintains a table with all items, indexed by Index. The (Index, item) pair is sent together in compressed lists until the compressor gains enough confidence that the decompressor has observed the mapping between the item and its Index. Such confidence is obtained by receiving an acknowledgment from the decompressor in R-mode, and in U/O-mode by sending L (Index, item) pairs (not necessarily consecutively). After that, the Index alone is sent in compressed lists to indicate the corresponding item. The compressor may reassign an existing Index to a new item, and then needs to re-establish the mapping in the same manner as above.

The decompressor conceptually maintains a table that contains all (Index, item) pairs it knows about. The table is updated whenever an (Index, item) pair is received (and decompression is verified by a CRC). The decompressor retrieves the item from the table whenever an Index without an accompanying item is received.

5.8.1.1. Translation table in R-mode

At the compressor side, an entry in the Translation Table has the following structure.

Index i	Known	item	SN1, SN2, ...
---------	-------	------	---------------

The Known flag indicates whether the mapping between Index i and item has been established, i.e., if Index i alone can be sent in compressed lists. Known is initially zero. It is also set to zero whenever Index i is assigned to a new item. Known is set to one when the corresponding (Index, item) pair is acknowledged. Acknowledgments are based on the RTP Sequence Number, so a list of RTP Sequence Numbers of all packets which contain the (Index, item) pair is included in the translation table. When a packet with a sequence number in the sequence number list is acknowledged, the Known flag is set, and the sequence number list can be discarded.

Each entry in the Translation Table at the decompressor side has the following structure:

```

      +-----+-----+
Index i | Known | item |
      +-----+-----+

```

All Known fields are initialized to zero. Whenever the decompressor receives an (Index, item) pair, it inserts item into the table at position Index and sets the Known flag in that entry to one. If an index without an accompanying item is received for which the Known flag is zero, the header **MUST** be discarded and a **NACK SHOULD** be sent.

#### 5.8.1.2. Translation table in U/O-modes

At the compressor side, each entry in the Translation Table has the following structure:

```

      +-----+-----+-----+
Index | Known | item | Counter |
      +-----+-----+-----+

```

The Index, Known, and item fields have the same meaning as in section 5.8.1.1.

Known is set when the (Index, item) pair has been sent in L compressed lists (not necessarily consecutively). The Counter field keeps track of how many times the pair has been sent. Counter is set to 0 for each new entry added to the table, and whenever Index is assigned to a new item. Counter is incremented by 1 whenever an (Index, item) pair is sent. When the counter reaches L, the Known field is set and after that only the Index needs to be sent in compressed lists.

At the decompressor side, the Translation Table is the same as the Translation Table defined in R-mode.

#### 5.8.2. Reference list determination

In reference based compression schemes (i.e., addition or deletion based schemes), compression and decompression of a list (curr\_list) are based on a reference list (ref\_list) which is assumed to be present in the context of both compressor and decompressor. The compressed list is an encoding of the differences between curr\_list and ref\_list. Upon reception of a compressed list, the decompressor applies the differences to its reference list in order to obtain the original list.

To identify the reference list (to be) used, each compressed list carries an identifier (`ref_id`). The reference list is established by different methods in R-mode and U/O-mode.

#### 5.8.2.1. Reference list in R-mode and U/O-mode

In R-mode, the choice of reference list is based on acknowledgments, i.e., the compressor uses as `ref_list` the latest list which has been acknowledged by the decompressor. The `ref_list` is updated only upon receiving an acknowledgment. The least significant bits of the RTP Sequence Number of the acknowledged packet are used as the `ref_id`.

In U/O-mode, a sequence of identical lists are considered as belonging to the same generation and are all assigned the same generation identifier (`gen_id`). `Gen_id` increases by 1 each time the list changes and is carried in compressed and uncompressed lists that are candidates for being used as reference lists. Normally, `Gen_id` must have been repeated in at least `L` headers before the list can be used as a `ref_list`. However, some acknowledgments may be sent in O-mode (and also in U-mode), and whenever an acknowledgment for a header is received, the list of that header is considered known and need not be repeated further. The least significant bits of the `Gen_id` is used as the `ref_id` in U/O-mode.

The logic of the compressor and decompressor for reference based list compression is similar to that for SN and TS. The principal difference is that the decompressor maintains a sliding window with candidates for `ref_list`, and retrieves `ref_list` from the sliding window using the `ref_id` of the compressed list.

Logic of compressor:

- a) In the IR state, the compressor sends Generic lists (see 5.8.5) containing all items of the current list in order to establish or refresh the context of the decompressor.

In R-mode, such Generic lists are sent until a header is acknowledged. The list of that header can be used as a reference list to compress subsequent lists.

In U/O-mode, the compressor sends generation identifiers with the Generic lists until

- 1) a generation identifier has been repeated `L` times, or
- 2) an acknowledgment for a header carrying a generation identifier has been received.

The repeated (1) or acknowledged (2) list can be used as a reference list to compress subsequent lists and is kept together with its generation identifier.

- b) When not in the IR state, the compressor moves to the FO state when it observes a difference between `curr_list` and the previous list. It sends compressed lists based on `ref_list` to update the context of the decompressor. (However, see d).)

In R-mode, the compressor keeps sending compressed lists using the same reference until it receives an acknowledgment for a packet containing the newest list. The compressor may then move to the SO state with regard to the list.

In U/O-mode, the compressor keeps sending compressed lists with generation identifiers until

- 1) a generation identifier has been repeated L times, or
- 2) an acknowledgment for a header carrying the latest generation identifier has been received.

The repeated or acknowledged list is used as the future reference list. The compressor may move to the SO state with regard to the list.

- c) In R-mode, the compressor maintains a sliding window containing the lists which have been sent to update the context of the decompressor and have not yet been acknowledged. The sliding window shrinks when an acknowledgment arrives: all lists sent before the acknowledged list are removed. The compressor may use the Index to represent items of lists in the sliding window.

In U/O-mode, the compressor needs to store

- 1) the reference list and its generation identifier, and
  - 2) if the current generation identifier is different from the reference generation, the current list and the sequence numbers with which the current list has been sent.
- (2) is needed to determine if an acknowledgment concerns the latest generation. It is not needed in U-mode.

- d) In U/O-mode, the compressor may choose to not send a generation identifier with a compressed list. Such lists without generation identifiers are not assigned a new generation identifier and must

not be used as future reference lists. They do not update the context. This feature is useful when a new list is repeated few times and the list then reverts back to its old value.

Logic of decompressor:

- e) In R-mode, the decompressor acknowledges all received uncompressed or compressed lists which establish or update the context. (Such compressed headers contain a CRC.)

In O-mode, the decompressor MAY acknowledge a list with a new generation identifier, see section 5.4.2.2.

In U-mode, the decompressor MAY acknowledge a list sent in an IR packet, see section 5.3.2.3.

- f) The decompressor maintains a sliding window which contains the lists that may be used as reference lists.

In R-mode, the sliding window contains lists which have been acknowledged but not yet used as reference lists.

In U/O-mode, the sliding window contains at most one list per generation. It contains all generations seen by the decompressor newer than the last generation used as a reference.

- g) When the decompressor receives a compressed list, it retrieves the proper `ref_list` from the sliding window based on the `ref_id`, and decompresses the compressed list obtaining `curr_list`.

In R-mode, `curr_list` is inserted into the sliding window if an acknowledgment is sent for it. The sliding window is shrunk by removing all lists received before `ref_list`.

In U/O-mode, `curr_list` is inserted into the sliding window together with its generation identifier if the compressed list had a generation identifier and the sliding window does not contain a list with that generation identifier. All lists with generations older than `ref_id` are removed from the sliding window.

### 5.8.3. Encoding schemes for the compressed list

Four encoding schemes for the compressed list are described here. The exact formats of the compressed CSRC list and compressed IP extension header list using these encoding schemes are described in sections 5.8.5-5.8.6.

### Generic scheme

In contrast to subsequent schemes, this scheme does not rely on a reference list having been established. The entire list is sent, using table based compression for each individual item. The generic scheme is always used when establishing the context of the decompressor and may also be used at other times, as the compressor sees fit.

### Insertion Only scheme

When the new list can be constructed from `ref_list` by adding items, a list of the added items is sent (using table based compression), along with the positions in `ref_list` where the new items will be inserted. An insertion bit mask indicates the insertion positions in `ref_list`.

Upon reception of a list compressed according to the Insertion Only scheme, `curr_list` is obtained by scanning the insertion bit mask from left to right. When a '0' is observed, an item is copied from the `ref_list`. When a '1' is observed, an item is copied from the list of added items. If a '1' is observed when the list of added items has been exhausted, an error has occurred and decompression fails: The header MUST NOT be delivered to upper layers; it should be discarded, and MUST NOT be acknowledged nor used as a reference.

To construct the insertion bit mask and the list of added items, the compressor MAY use the following algorithm:

- 1) An empty bit list and an empty Inserted Item list are generated as the starting point.
  - 2) Start by considering the first item of `curr_list` and `ref_list`.
  - 3) If `curr_list` has a different item than `ref_list`,
    - a set bit (1) is appended to the bit list;
    - the first item in `curr_list` (represented using table-based item compression) is appended to the Inserted Item list;
    - advance to the next item of `curr_list`;
- otherwise,
- a zero bit (0) is appended to the bit list;
  - advance to the next item of `curr_list`;
  - advance to the next item of `ref_list`.

- 4) Repeat 3) until curr\_list has been exhausted.
- 5) If the length of the bit list is less than the required bit mask length, append additional zeroes.

#### Removal Only scheme

This scheme can be used when curr\_list can be obtained by removing some items in ref\_list. The positions of the items which are in ref\_list, but not in curr\_list, are sent as a removal bit mask.

Upon reception of the compressed list, the decompressor obtains curr\_list by scanning the removal bit mask from left to right. When a '0' is observed, the next item of ref\_list is copied into curr\_list. When a '1' is observed, the next item of ref\_list is skipped over without being copied. If a '0' is observed when ref\_list has been exhausted, an error has occurred and decompression fails: The header MUST NOT be delivered to upper layers; it should be discarded, and MUST NOT be acknowledged nor used as a reference.

To construct the removal bit mask and the list of added items, the compressor MAY use the following algorithm:

- 1) An empty bit list is generated as the starting point.
- 2) Start by considering the first item of curr\_list and ref\_list.
- 3) If curr\_list has a different item than ref\_list,  
    a set bit (1) is appended to the bit list;  
    advance to the next item of ref\_list;
- otherwise,  
    a zero bit (0) is appended to the bit list;  
    advance to the next item of curr\_list;  
    advance to the next item of ref\_list.
- 4) Repeat 3) until curr\_list has been exhausted.
- 5) If the length of the bit list is less than the required bit mask length, append additional ones.

#### Remove Then Insert scheme

In this scheme, `curr_list` is obtained by first removing items from `ref_list`, and then inserting items into the resulting list. A removal bit mask, an insertion bit mask, and a list of added items are sent.

Upon reception of the compressed list, the decompressor processes the removal bit mask as in the Removal Only scheme. The resulting list is then used as the reference list when the insertion bit mask and the list of added items are processed, as in the Insertion Only scheme.

#### 5.8.4. Special handling of IP extension headers

In CSRC list compression, each CSRC is assigned an index. In contrast, in IP extension header list compression an index is usually associated with a type of extension header. When there is more than one IP header, there is more than one list of extension headers. An index per type per list is then used.

The association with a type means that a new index need not always be used each time a field in an IP extension header changes. However, when a field in an extension header changes, the mapping between the index and the new value of the extension header needs to be established, except in the special handling cases defined in the following subsections.

##### 5.8.4.1. Next Header field

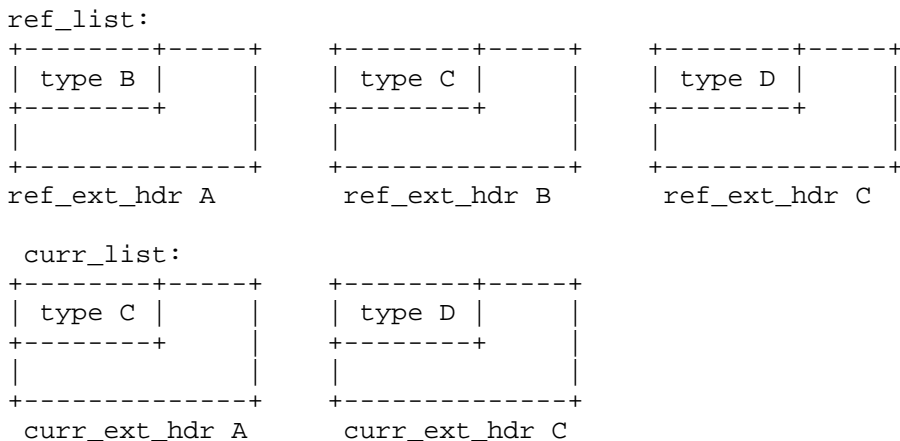
The next header field in an IP header or extension header changes whenever the type of the immediately following header changes, e.g., when a new extension header is inserted after it, when the immediate subsequent extension header is removed from the list, or when the order of extension headers is changed. Thus it may not be uncommon that, for a given header, the next header field changes while the remaining fields do not change.

Therefore, in the case that only the next header field changes, the extension header is considered to be unchanged and rules for special treatment of the change in the next header field are defined below.

All communicated uncompressed extension header items indicate their own type in their Next Header field. Note that the rules below explain how to treat the Next Header fields while showing the conceptual reference list as an exact recreation of the original uncompressed extension header list.

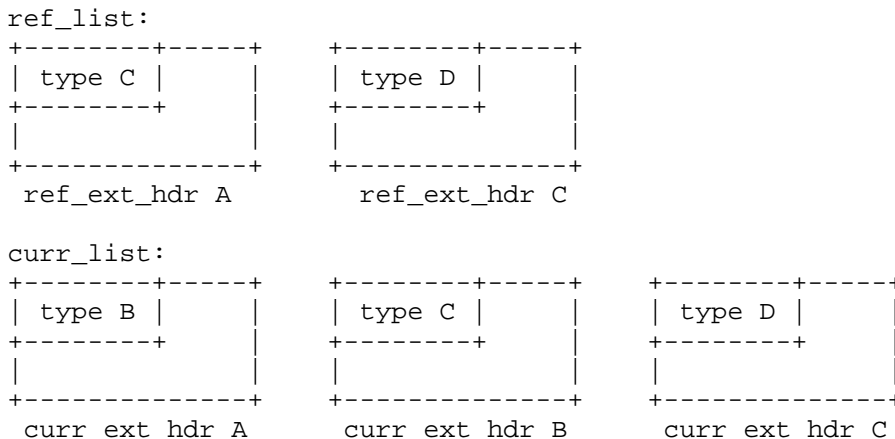


- a) When a subsequent extension header is removed from the list, the new value of the next header field is obtained from the reference extension header list. For example, assume that the reference header list (ref\_list) consists of headers A, B and C (ref\_ext\_hdr A, B, C), and the current extension header list (curr\_list) only consists of extension headers A and C (curr\_ext\_hdr A, C). The order and value of the next header fields of these extension headers are as follows.



Comparing the curr\_ext\_hdr A in curr\_list and the ref\_ext\_hdr A in ref\_list, the value of next header field is changed from "type B" to "type C" because of the removal of extension header B. The new value of the next header field in curr\_ext\_hdr A, i.e., "type C", does not need to be sent to the decompressor. Instead, it is retrieved from the next header field of the removed ref\_ext\_hdr B.

- b) When a new extension header is inserted after an existing extension header, the next header field in the communicated item will carry the type of itself, rather than the type of the header that follows. For example, assume that the reference header list (ref\_list) consists of headers A and C (ref\_ext\_hdr A, C), and the current header list (curr\_list) consists of headers A, B and C (curr\_ext\_hdr A, B, C). The order and the value of the next header fields of these extension headers are as follows.



Comparing the curr\_list and the ref\_list, the value of the next header field in extension header A is changed from "type C" to "type B".

The uncompressed curr\_ext\_hdr B is carried in the compressed header list. However, it carries "type B" instead of "type C" in its next header field. When the decompressor inserts a new header after curr\_ext\_hdr A, the next header field of A is taken from the new header, and the next header field of the new header is taken from ref\_ext\_hdr A.

- c) Some headers whose compression is defined in this document do not contain Next Header fields or do not have their Next Header field in the standard position (first octet of the header). The GRE and ESP headers are such headers. When sent as uncompressed items in lists, these headers are modified so that they do have a Next Header field as their first octet (see 5.8.4.3 and 5.8.4.4). This is necessary to enable the decompressor to decode the item.

5.8.4.2. Authentication Header (AH)

The sequence number field in the AH [AH] contains a monotonically increasing counter value for a security association. Therefore, when comparing curr\_list with ref\_list, if the sequence number in AH changes and SPI field does not change, the AH is not considered as changed.

If the sequence number in the AH linearly increases as the RTP Sequence Number increases, and the compressor is confident that the decompressor has obtained the pattern, the sequence number in AH need not be sent. The decompressor applies linear extrapolation to reconstruct the sequence number in the AH.

Otherwise, a compressed sequence number is included in the IPX compression field in an Extension 3 of an UOR-2 header.

The authentication data field in AH changes from packet to packet and is sent as-is. If the uncompressed AH is sent, the authentication data field is sent inside the uncompressed AH; otherwise, it is sent after the compressed IP/UDP/RTP and IPv6 extension headers and before the payload. See beginning of section 5.7.

Note: The payload length field of the AH uses a different notion of length than other IPv6 extension headers.

#### 5.8.4.3. Encapsulating Security Payload Header (ESP)

When the Encapsulating Security Payload Header (ESP) [ESP] is present and an encryption algorithm other than NULL is being used, the UDP and RTP headers are both encrypted and cannot be compressed. The ESP header thus ends the compressible header chain. The ROHC ESP profile defined in section 5.12 MAY be used for the stream in this case.

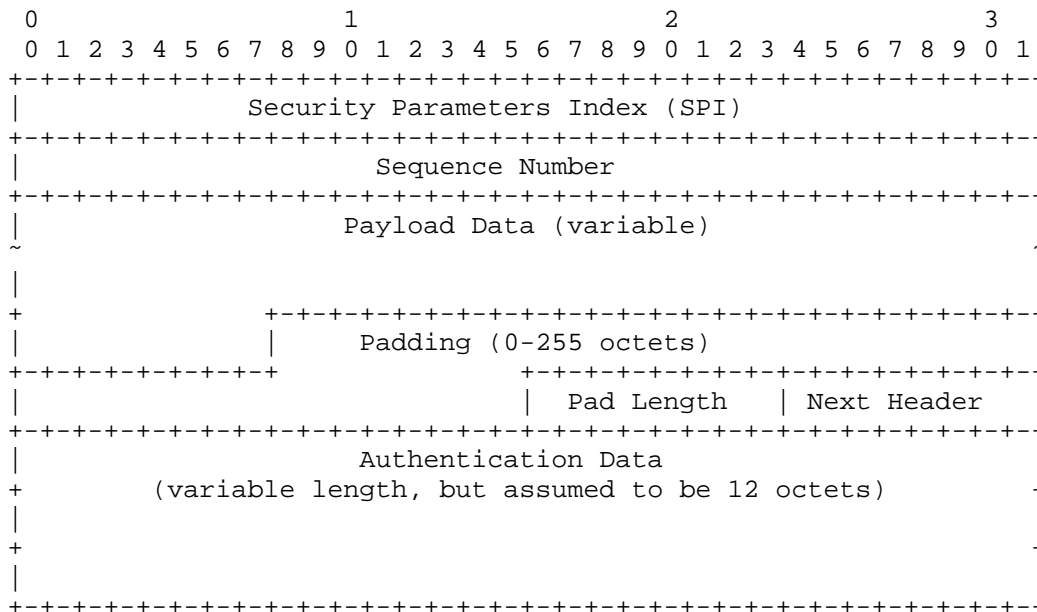
A special case is when the NULL encryption algorithm is used. This is the case when the ESP header is used for authentication only, and not for encryption. The payload is not encrypted by the NULL encryption algorithm, so compression of the rest of the header chain is possible. The rest of this section describes compression of the ESP header when the NULL encryption algorithm is used with ESP.

It is not possible to determine whether NULL encryption is used by inspecting a header in the stream, this information is present only at the encryption endpoints. However, a compressor may attempt compression under the assumption that the NULL encryption algorithm is being used, and later abort compression when the assumption proves to be false.

The compressor may, for example, inspect the Next Header fields and the header fields supposed to be static in subsequent headers in order to determine if NULL encryption is being used. If these change unpredictably, an encryption algorithm other than NULL is probably being used and compression of subsequent headers SHOULD be aborted. Compression of the stream is then either discontinued, or a profile that compresses only up to the ESP header may be used (see 5.12). While attempting to compress the header, the compressor should use the SPI of the ESP header together with the destination IP address as the defining fields for determining which packets belong to the stream.

In the ESP header [ESP, section 2], the fields that can be compressed are the SPI, the sequence number, the Next Header, and the padding bytes if they are in the standard format defined in [ESP]. (As always, the decompressor reinserts these fields based on the information in the context. Care must be taken to correctly reinsert all the information as the Authentication Data must be verified over the exact same information it was computed over.)

ESP header [ESP, section 2]:



SPI: Static. If it changes, it needs to be reestablished.

Sequence Number: Not sent when the offset from the sequence number of the compressed header is constant. When the offset is not constant, the sequence number may be compressed by sending LSBs. See 5.8.4.

Payload Data: This is where subsequent headers are to be found. Parsed according to the Next Header field.

Padding: The padding octets are assumed to be as defined in [ESP], i.e., to take the values 1, 2, ..., k, where k = Pad Length. If the padding in the static context has this pattern, padding in compressed headers is assumed to have this pattern as well and is removed. If padding in the static context does not have this pattern, the padding is not removed.

Pad Length: Dynamic. Always sent. 14th octet from end of packet.

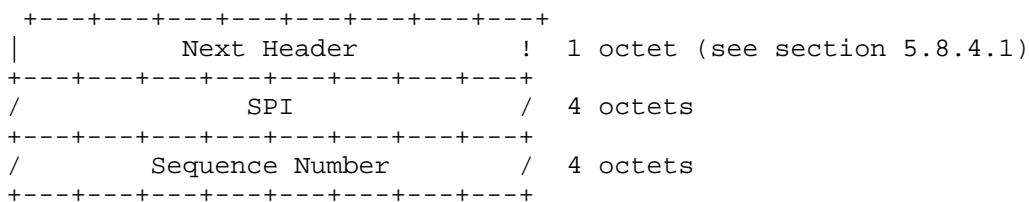
Next Header: Static. 13th octet from end of packet.

Authentication Data: Can have variable length, but when compression of NULL-encryption ESP header is attempted, it is assumed to have length 12 octets.

The sequence number in ESP has the same behavior as the sequence number field in AH. When it increases linearly, it can be compressed to zero bits. When it does not increase linearly, a compressed sequence number is included in the IPX compression field in an Extension 3 of an UOR-2 header.

The information which is part of an uncompressed item of a compressed list is the Next Header field, followed by the SPI and the Sequence Number. Padding, Pad Length, Next Header, and Authentication Data are sent as-is at the end of the packet. This means that the Next Header occurs in two places.

Uncompressed ESP list item:



When sending Uncompressed ESP list items, all ESP fields near the the end of the packet are left untouched (Padding, Pad Length, Next Header, Authentication Data).

A compressed item consists of a compressed sequence number. When an item is compressed, Padding (if it follows the 1, 2, ..., k pattern) and Next Header are removed near the end of the packet. Authentication Data and Pad Length remain as-is near the end of the packet.

5.8.4.4. GRE Header [RFC 2784, RFC 2890]

The GRE header is a set of flags, followed by a mandatory Protocol Type and optional parts as indicated by the flags.

The sequence number field in the GRE header contains a counter value for a GRE tunnel. Therefore, when comparing curr\_list with ref\_list, if the sequence number in GRE changes, the GRE is not considered as changed.

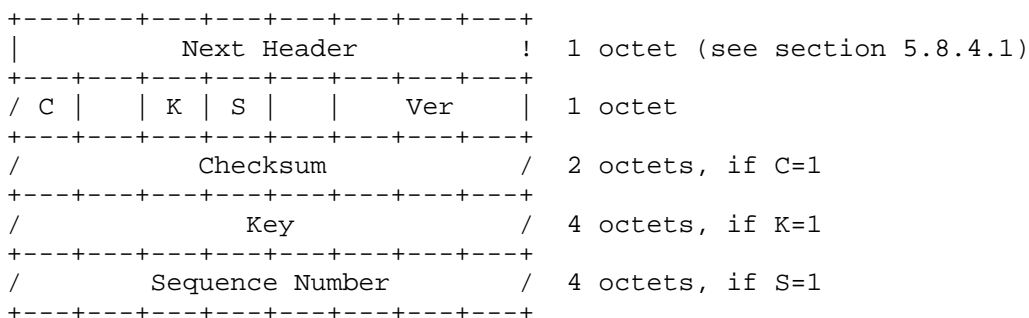
If the sequence number in the GRE header linearly increases as the RTP Sequence Number increases and the compressor is confident that the decompressor has received the pattern, the sequence number in GRE need not be sent. The decompressor applies linear extrapolation to reconstruct the sequence number in the GRE header.

Otherwise, a compressed sequence number is included in the IPX compression field in an Extension 3 of an UOR-2 header.

The checksum data field in GRE, if present, changes from packet to packet and is sent as-is. If the uncompressed GRE header is sent, the checksum data field is sent inside the uncompressed GRE header; otherwise, if present, it is sent after the compressed IP/UDP/RTP and IPv6 extension headers and before the payload. See beginning of section 5.7.

In order to allow simple parsing of lists of items, an uncompressed GRE header sent as an item in a list is modified from the original GRE header in the following manner: 1) the 16-bit Protocol Type field that encodes the type of the subsequent header using Ether types (see Ether types section in [ASSIGNED]) is removed. 2) A one-octet Next Header field is inserted as the first octet of the header. The value of the Next Header field corresponds to GRE (this value is 47 according to the Assigned Internet Protocol Number section of [ASSIGNED]) when the uncompressed item is to be inserted in a list, and to the type of the subsequent header when the uncompressed item is in a Generic list. Note that this implies that only GRE headers with Ether types that correspond to an IP protocol number can be compressed.

Uncompressed GRE list item:



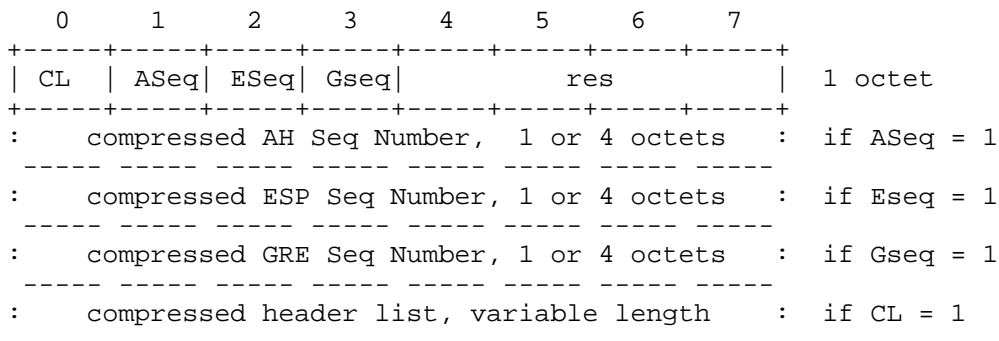
The bits left blank in the second octet are set to zero when sending and ignored when received.

The fields Reserved0 and Reserved1 of the GRE header [GRE2] must be all zeroes; otherwise, the packet cannot be compressed by this profile.

5.8.5. Format of compressed lists in Extension 3

5.8.5.1. Format of IP Extension Header(s) field

In Extension 3 (section 5.7.5), there is a field called IP extension header(s). This section describes the format of that field.



- ASeq: indicates presence of compressed AH Seq Number
- ESeq: indicates presence of compressed ESP Seq Number
- GSeq: indicates presence of compressed GRE Seq Number
- CL: indicates presence of compressed header list
- res: reserved; set to zero when sending, ignored when received

When Aseq, Eseq, or Gseq is set, the corresponding header item (AH, ESP, or GRE header) is compressed. When not set, the corresponding header item is sent uncompressed or is not present.

The format of compressed AH, ESP and GRE Sequence Numbers can each be either of the following:





GP: Indicates presence of `gen_id` field.

CC: CSRC counter from original RTP header.

`gen_id`: Identifier for a sequence of identical lists. It is present in U/O-mode when the compressor decides that it may use this list as a future reference list.

XI 1, ..., XI m: m XI items. The format of an XI item is as follows:

```

PS = 0:  +---+---+---+---+
          | X |   Index   |
          +---+---+---+---+

          0  1  2  3  4  5  6  7
PS = 1:  +---+---+---+---+---+---+---+
          | X |                   Index                   |
          +---+---+---+---+---+---+---+

```

X = 1 indicates that the item corresponding to the Index is sent in the item 0, ..., item n list.

X = 0 indicates that the item corresponding to the Index is not sent.

When 4-bit XI items are used and  $m > 1$ , the XI items are placed in octets in the following manner:

```

          0  1  2  3  4  5  6  7
          +---+---+---+---+---+---+---+
          |   XI k   |   XI k + 1   |
          +---+---+---+---+---+---+---+

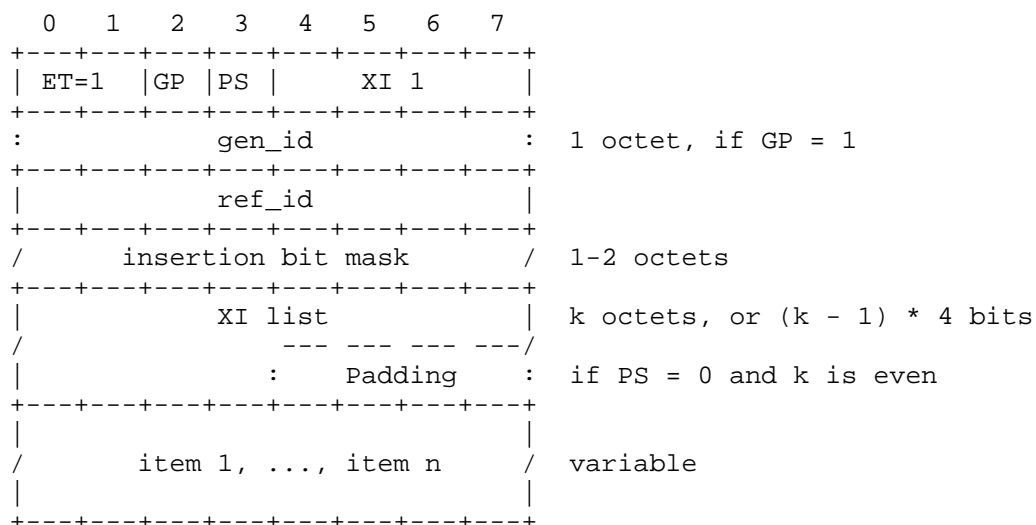
```

Padding: A 4-bit padding field is present when PS = 0 and m is odd. The Padding field is set to zero when sending and ignored when receiving.

Item 1, ..., item n:

Each item corresponds to an XI with X = 1 in XI 1, ..., XI m.

## 5.8.6.2. Encoding Type 1 (insertion only scheme)



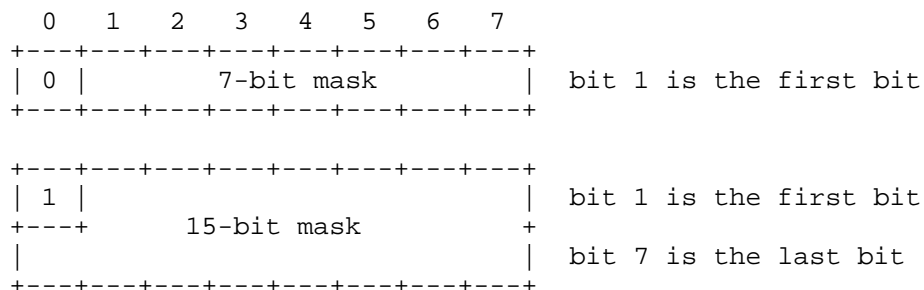
Unless explicitly stated otherwise, fields have the same meaning and values as for encoding type 0.

ET: Encoding type is one (1).

XI 1: When PS = 0, the first 4-bit XI item is placed here.  
When PS = 1, the field is set to zero when sending, and ignored when receiving.

ref\_id: The identifier of the reference CSRC list used when the list was compressed. It is the 8 least significant bits of the RTP Sequence Number in R-mode and gen\_id (see section 5.8.2) in U/O-mode.

insertion bit mask: Bit mask indicating the positions where new items are to be inserted. See Insertion Only scheme in section 5.8.3. The bit mask can have either of the following two formats:

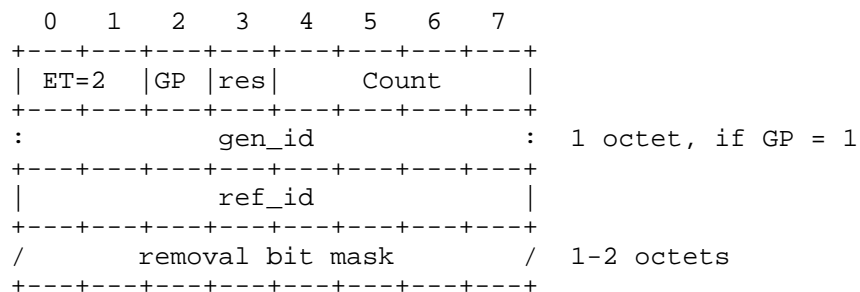


XI list: XI fields for items to be inserted. When the insertion bit mask has  $k$  ones, the total number of XI fields is  $k$ . When  $PS = 1$ , all XI fields are in the XI list. When  $PS = 0$ , the first XI field is in the XI 1 field, and the remaining  $k - 1$  XI fields are in the XI list.

Padding: Present when  $PS = 0$  and  $k$  is even.

item 1, ..., item  $n$ : One item for each XI field with the X bit set.

#### 5.8.6.3. Encoding Type 2 (removal only scheme)



Unless explicitly stated otherwise, fields have the same meaning and values as in section 5.8.5.2.

ET: Encoding type is 2.

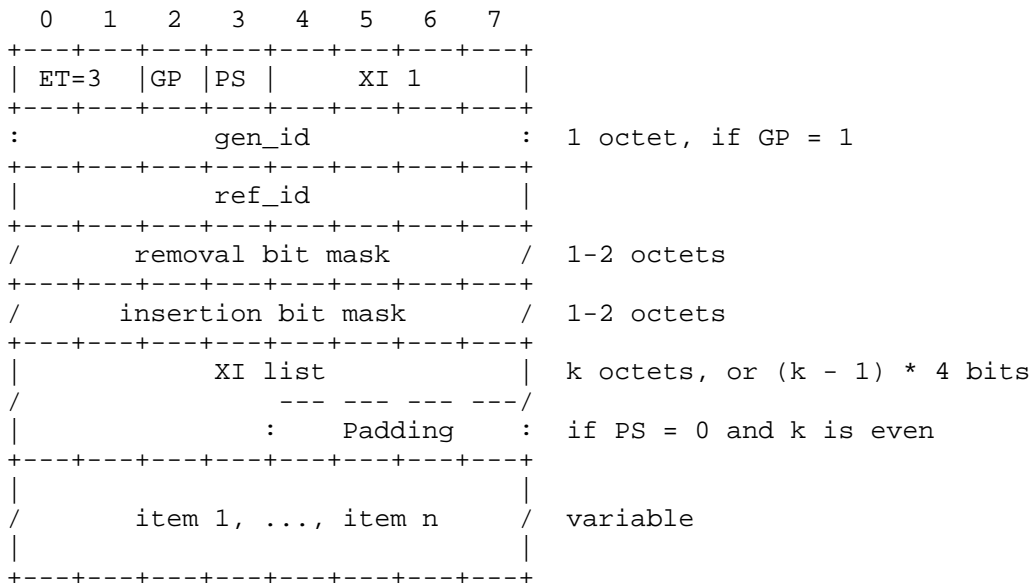
res: Reserved. Set to zero when sending, ignored when received.

Count: Number of elements in `ref_list`.

removal bit mask: Indicates the elements in ref\_list to be removed in order to obtain the current list. See section 5.8.3. The removal bit mask has the same format as the insertion bit mask of section 5.8.6.3.

5.8.6.4. Encoding Type 3 (remove then insert scheme)

See section 5.8.3 for a description of the Remove then insert scheme.



The fields in this header have the same meaning and formats as in section 5.8.5.2, except when explicitly stated otherwise below.

ET: Encoding type is 3.

removal bit mask: See section 5.8.6.3.

5.8.7. CRC coverage for extension headers

All fields of extension headers are CRC-STATIC, with the following exceptions which are CRC-DYNAMIC.

- 1) Entire AH header.
- 2) Entire ESP header.
- 3) Sequence number in GRE, Checksum in GRE

## 5.9. Header compression CRCs, coverage and polynomials

This chapter describes how to calculate the CRCs used in packet headers defined in this document. (Note that another type of CRC is defined for reconstructed units in section 5.2.5.)

### 5.9.1. IR and IR-DYN packet CRCs

The CRC in the IR and IR-DYN packet is calculated over the entire IR or IR-DYN packet, excluding Payload and including CID or any Add-CID octet. For purposes of computing the CRC, the CRC field in the header is set to zero.

The initial content of the CRC register is to be preset to all 1's.

The CRC polynomial to be used for the 8-bit CRC is:

$$C(x) = 1 + x + x^2 + x^8$$

### 5.9.2. CRCs in compressed headers

The CRC in compressed headers is calculated over all octets of the entire original header, before compression, in the following manner.

The octets of the header are classified as either CRC-STATIC or CRC-DYNAMIC, and the CRC is calculated over:

- 1) the concatenated CRC-STATIC octets of the original header, placed in the same order as they appear in the original header, followed by
- 2) the concatenated CRC-DYNAMIC octets of the original header, placed in the same order as they appear in the original header.

The intention is that the state of the CRC computation after 1) will be saved. As long as the CRC-STATIC octets do not change, the CRC calculation will then only need to process the CRC-DYNAMIC octets.

In a typical RTP/UDP/IPv4 header, 25 octets are CRC-STATIC and 15 are CRC-DYNAMIC. In a typical RTP/UDP/IPv6 header, 49 octets are CRC-STATIC and 11 are CRC-DYNAMIC. This technique will thus reduce the computational complexity of the CRC calculation by roughly 60% for RTP/UDP/IPv4 and by roughly 80% for RTP/UDP/IPv6.

Note: Whenever the CRC-STATIC fields change, the new saved CRC state after 1) is compared with the old state. If the states are identical, the CRC cannot catch the error consisting in the decompressor not having updated the static context. In U/O-mode the

compressor SHOULD then for a while use packet types with another CRC length, for which there is a difference in CRC state, to ensure error detection.

The initial content of the CRC register is preset to all 1's.

The polynomial to be used for the 3 bit CRC is:

$$C(x) = 1 + x + x^3$$

The polynomial to be used for the 7 bit CRC is:

$$C(x) = 1 + x + x^2 + x^3 + x^6 + x^7$$

The CRC in compressed headers is calculated over the entire original header, before compression.

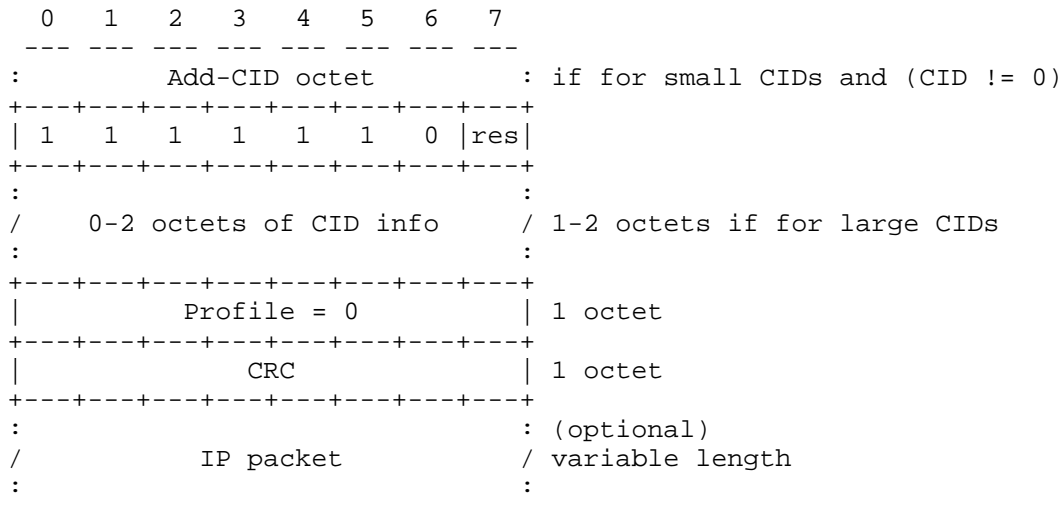
#### 5.10. ROHC UNCOMPRESSED -- no compression (Profile 0x0000)

In ROHC, compression has not been defined for all kinds of IP headers. Profile 0x0000 provides a way to send IP packets without compressing them. This can be used for IP fragments, RTCP packets, and in general for any packet for which compression of the header has not been defined, is not possible due to resource constraints, or is not desirable for some other reason.

After initialization, the only overhead for sending packets using Profile 0x0000 is the size of the CID. When uncompressed packets are frequent, Profile 0x0000 should be associated with a CID with size zero or one octet. There is no need to associate Profile 0x0000 with more than one CID.

##### 5.10.1. IR packet

The initialization packet (IR packet) for Profile 0x0000 has the following format:



res: Always zero.

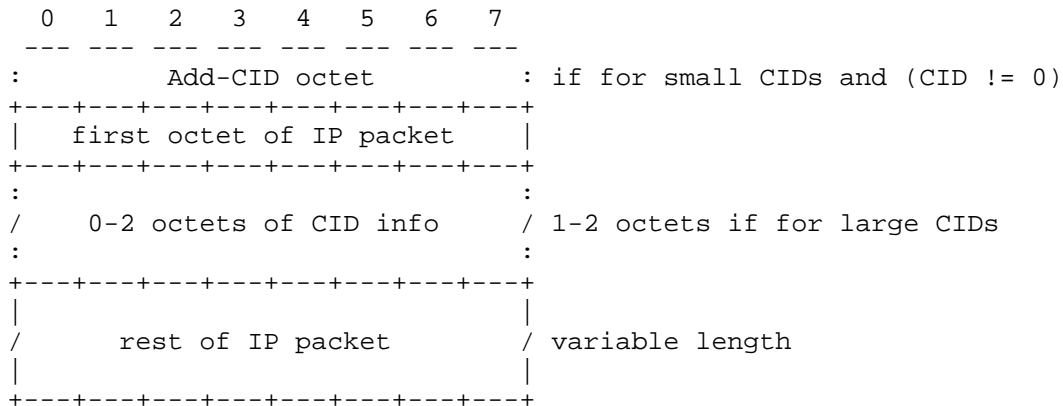
Profile: 0.

CRC: 8-bit CRC, computed using the polynomial of section 5.9.1. The CRC covers the first octet of the IR packet through the Profile octet of the IR packet, i.e., it does not cover the CRC itself or the IP packet.

IP packet: An uncompressed IP packet may be included in the IR packet. The decompressor determines if the IP packet is present by considering the length of the IR packet.

5.10.2. Normal packet

A Normal packet is a normal IP packet plus CID information. When the channel uses small CIDs, and profile 0x0000 is associated with a CID > 0, an Add-CID octet is prepended to the IP packet. When the channel uses large CIDs, the CID is placed so that it starts at the second octet of the Normal packet.



Note that the first octet of the IP packet starts with the bit pattern 0100 (IPv4) or 0110 (IPv6). This does not conflict with any reserved packet types. Hence, no bits in addition to the CID are needed. The profile is reasonably future-proof since problems do not occur until IP version 14.

### 5.10.3. States and modes

There are two modes in Profile 0x0000: Unidirectional mode and Bidirectional mode. In Unidirectional mode, the compressor repeats the IR packet periodically. In Bidirectional mode, the compressor never repeats the IR packet. The compressor and decompressor always start in Unidirectional mode. Whenever feedback is received, the compressor switches to Bidirectional mode.

The compressor can be in either of two states: the IR state or the Normal state. It starts in the IR state.

- a) IR state: Only IR packets can be sent. After sending a small number of IR packets (only one when refreshing), the compressor switches to the Normal state.
- b) Normal state: Only Normal packets can be sent. When in Unidirectional mode, the compressor periodically transits back to the IR state. The length of the period is implementation dependent, but should be fairly long. Exponential backoff may be used.
- c) When feedback is received in any state, the compressor switches to Bidirectional mode.



The decompressor can be in either of two states: NO\_CONTEXT or FULL\_CONTEXT. It starts in NO\_CONTEXT.

- d) When an IR packet is received in the NO\_CONTEXT state, the decompressor first verifies the packet using the CRC. If the packet is OK, the decompressor 1) moves to the FULL\_CONTEXT state, 2) delivers the IP packet to upper layers if present, 3) MAY send an ACK. If the packet is not OK, it is discarded without further action.
- e) When any other packet is received in the NO\_CONTEXT state, it is discarded without further action.
- f) When an IR packet is received in the FULL\_CONTEXT state, the packet is first verified using the CRC. If OK, the decompressor 1) delivers the IP packet to upper layers if present, 2) MAY send an ACK. If the packet is not OK, no action is taken.
- g) When a Normal packet is received in the FULL\_CONTEXT state, the CID information is removed and the IP packet is delivered to upper layers.

#### 5.10.4. Feedback

The only kind of feedback in Profile 0x0000 is ACKs. Profile 0x0000 MUST NOT be rejected. Profile 0x0000 SHOULD be associated with at most one CID. ACKs use the FEEDBACK-1 format of section 5.2. The value of the profile-specific octet in the FEEDBACK-1 ACK is 0 (zero).

#### 5.11. ROHC UDP -- non-RTP UDP/IP compression (Profile 0x0002)

UDP/IP headers do not have a sequence number which is as well-behaved as the RTP Sequence Number. For UDP/IPv4, there is an IP-ID field which may be echoed in feedback information, but when no IPv4 header is present such feedback identification becomes problematic.

Therefore, in the ROHC UDP profile, the compressor generates a 16-bit sequence number SN which increases by one for each packet received in the packet stream. This sequence number is thus relatively well-behaved and can serve as the basis for most mechanisms described for ROHC RTP. It is called SN or UDP SN below. Unless stated otherwise, the mechanisms of ROHC RTP are used also for ROHC UDP, with the UDP SN taking the role of the RTP Sequence Number.

The ROHC UDP profile always uses  $p = -1$  when interpreting the SN, since there will be no repetitions or reordering of the compressor-generated SN. The interpretation interval thus always starts with  $(\text{ref\_SN} + 1)$ .

#### 5.11.1. Initialization

The static context for ROHC UDP streams can be initialized in either of two ways:

- 1) By using an IR packet as in section 5.7.7.1, where the profile is two (2) and the static chain ends with the static part of an UDP packet. At the compressor, UDP SN is initialized to a random value when the IR packet is sent.
- 2) By reusing an existing context where the existing static chain contains the static part of a UDP packet, e.g., the context of a stream compressed using ROHC RTP (profile 0x0001). This is done with an IR-DYN packet (section 5.7.7.2) identifying profile 0x0002, where the dynamic chain corresponds to the prefix of the existing static chain that ends with the UDP header. UDP SN is initialized to the RTP Sequence Number if the earlier profile was profile 0x0001, and to a random number otherwise.

For ROHC UDP, the dynamic part of a UDP packet is different from section 5.7.7.5: a two-octet field containing the UDP SN is added after the Checksum field. This affects the format of dynamic chains in IR and IR-DYN packets.

Note: 2) can be used for packet streams which were initially assumed to be RTP streams, so that compression started with profile 0x0001, but were later found evidently not to be RTP streams.

#### 5.11.2. States and modes

ROHC UDP uses the same states and modes as ROHC RTP. Mode transitions and state logic are the same except when explicitly stated otherwise. Mechanisms dealing with fields in the RTP header (except the RTP SN) are not used. The decompressed UDP SN is never included in any header delivered to upper layers. The UDP SN is used in place of the RTP SN in feedback.

5.11.3. Packet types

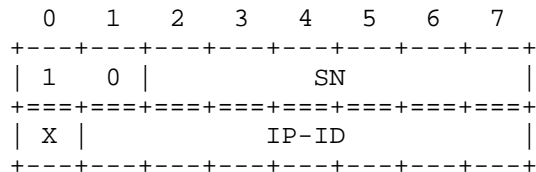
The general format of a ROHC UDP packet is the same as for ROHC RTP (see beginning of section 5.7). Padding and CIDs are the same, as is the feedback packet type (5.7.6.1) and the feedback. IR and IR-DYN packets (5.7.7) are changed as described in 5.11.1.

The general format of compressed packets is also the same, but there are differences in specific formats and extensions as detailed below. The differences are caused by removal of all RTP specific information except the RTP SN, which is replaced by the UDP SN.

Unless explicitly stated below, the packet formats are as in sections 5.7.1-6.

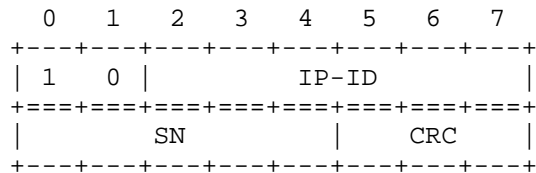
R-1

The TS field is replaced by an IP-ID field. The M flag has become part of IP-ID. The X bit has moved. The formats R-1-ID and R-1-TS are not used.



UO-1

The TS field is replaced by an IP-ID field. The M flag has become part of SN. Formats UO-1-ID and UO-1-TS are not used.



UOR-2

New format:

0	1	2	3	4	5	6	7	
1	1	0	SN					
X	CRC							

#### 5.11.4. Extensions

Extensions are as in 5.7.5, with the following exceptions:

Extension 0:

0	0	SN			IP-ID		
---	---	----	--	--	-------	--	--

Extension 1:

0	1	SN			IP-ID		
				IP-ID			

Extension 2:

1	0	SN			IP-ID2		
				IP-ID2			
				IP-ID			

IP-ID2: For outer IP-ID field.

Extension 3 is the same as Extension 3 in section 5.7.5, with the following exceptions.

1) The initial flag octet has the following format:

0	1	2	3	4	5	6	7
1	1	S	Mode		I	ip	ip2

Mode: Replaces R-TS and Tsc of 5.7.5. Provides mode information as was earlier done in RTP header flags and fields.

ip2: Replaces rtp bit of 5.7.5. Moved here from the Inner IP header flags octet.

- 2) The bit which was the ip2 flag in the Inner IP header flags in 5.7.5 is reserved. It is set to zero when sending and ignored when receiving.

#### 5.11.5. IP-ID

Treated as in ROHC RTP, but the offset is from UDP SN.

#### 5.11.6. Feedback

Feedback is as for ROHC RTP with the following exceptions:

- 1) UDP SN replaces RTP SN in feedback.
- 2) The CLOCK option (5.7.6.6) is not used.
- 3) The JITTER option (5.7.6.7) is not used.

#### 5.12. ROHC ESP -- ESP/IP compression (Profile 0x0003)

When the ESP header is being used with an encryption algorithm other than NULL, subheaders after the ESP header are encrypted and cannot be compressed. Profile 0x0003 is for compression of the chain of headers up to and including the ESP header in this case. When the NULL encryption algorithm is being used, other profiles can be used and could give higher compression rates. See section 5.8.4.3.

This profile is very similar to the ROHC UDP profile. It uses the ESP sequence number as the basis for compression instead of a generated number, but is otherwise very similar to ROHC UDP. The interpretation interval (value of p) for the ESP-based SN is as with ROHC RTP (profile 0x0001). Apart from this, unless stated explicitly below, mechanisms and formats are as for ROHC UDP.

##### 5.12.1. Initialization

The static context for ROHC ESP streams can be initialized in either of two ways:

- 1) by using an IR packet as in section 5.7.7.1, where the profile is three (3) and the static chain ends with the static part of an ESP header.

- 2) by reusing an existing context, where the existing static chain contains the static part of an ESP header. This is done with an IR-DYN packet (section 5.7.7.2) identifying profile 0x0003, where the dynamic chain corresponds to the prefix of the existing static chain that ends with the ESP header.

In contrast to ROHC UDP, no extra sequence number is added to the dynamic part of the ESP header: the ESP sequence number is the only element.

Note: 2) can be used for streams where compression has been initiated under the assumption that NULL encryption was being used with ESP. When it becomes obvious that an encryption algorithm other than NULL is being used, the compressor may send an IR-DYN according to 2) to switch to profile 0x0003 without having to send an IR packet.

#### 5.12.2. Packet types

The packet types for ROHC ESP are the same as for ROHC UDP, except that the ESP sequence number is used instead of the generated sequence number of ROHC UDP. The ESP header is not part of any compressed list in ROHC ESP.

### 6. Implementation issues

This document specifies mechanisms for the protocol and leaves many details on the use of these mechanisms to the implementers. This chapter is aimed to give guidelines, ideas and suggestions for implementing the scheme.

#### 6.1. Reverse decompression

This section describes an OPTIONAL decompressor operation to reduce the number of packets discarded due to an invalid context.

Once a context becomes invalid (e.g., when more consecutive packet losses than expected have occurred), subsequent compressed packets cannot immediately be decompressed correctly. Reverse decompression aims at decompressing such packets later instead of discarding them, by storing them until the context has been updated and validated and then attempting decompression.

Let the sequence of stored packets be  $i, i + 1, \dots, i + k$ , where  $i$  is the first packet and  $i + k$  is the last packet before the context was updated. The decompressor will attempt to recover the stored packets in reverse order, i.e., starting with  $i + k$ , and working back toward  $i$ . When a stored packet has been reconstructed, its correctness is verified using its CRC. Packets not carrying a CRC

must not be delivered to upper layers. Packets where the CRC succeeds are delivered to upper layers in their original order, i.e.,  $i, i + 1, \dots, i + k$ .

Note that this reverse decompression introduces buffering while waiting for the context to be validated and thereby introduces additional delay. Thus, it should be used only when some amount of delay is acceptable. For example, for video packets belonging to the same video frame, the delay in packet arrivals does not cause presentation time delay. Delay-insensitive streaming applications can also be tolerant of such delay. If the decompressor cannot determine whether the application can tolerate delay, it should not perform reverse decompression.

The following illustrates the decompression procedure in some detail:

1. The decompressor stores compressed packets that cannot be decompressed correctly due to an invalid context.
2. When the decompressor has received a context updating packet and the context has been validated, it proceeds to recover the last packet stored. After decompression, the decompressor checks the correctness of the reconstructed header using the CRC.
3. If the CRC indicates successful decompression, the decompressor stores the complete packet and attempts to decompress the preceding packet. In this way, the stored packets are recovered in reverse order until no compressed packets are left. For each packet, the decompressor checks the correctness of the decompressed headers using the header compression CRC.
4. If the CRC indicates an incorrectly decompressed packet, the reverse decompression attempt MUST be terminated and all remaining uncompressed packets MUST be discarded.
5. Finally, the decompressor forwards all the correctly decompressed packets to upper layers in their original order.

## 6.2. RTCP

RTCP is the RTP Control Protocol [RTP]. RTCP is based on periodic transmission of control packets to all participants in a session, using the same distribution mechanism as for data packets. Its primary function is to provide feedback from the data receivers on the quality of the data distribution. The feedback information may be used for issues related to congestion control functions, and directly useful for control of adaptive encodings.

In an RTP session there will be two types of packet streams: one with the RTP header and application data, and one with the RTCP control information. The difference between the streams at the transport level is in the UDP port numbers: the RTP port number is always even, the RTCP port number is that number plus one and therefore always odd [RTP, section 10]. The ROHC header compressor implementation has several ways at hand to handle the RTCP stream:

1. One compressor/decompressor entity carrying both types of streams on the same channel, using CIDs to distinguish between them. For sending a single RTP stream together with its RTCP packets on one channel, it is most efficient to set `LARGE_CIDS` to false, send the RTP packets with the implied CID 0 and use the Add-CID mechanism to send the RTCP packets.
2. Two compressor/decompressor entities, one for RTP and another one for RTCP, carrying the two types of streams on separate channels. This means that they will not share the same CID number space.

RTCP headers may simply be sent uncompressed using profile 0x0000. More efficiently, ROHC UDP compression (profile 0x0002) can be used.

### 6.3. Implementation parameters and signals

A ROHC implementation may have two kinds of parameters: configuration parameters that are mandatory and must be negotiated between compressor and decompressor peers, and implementation parameters that are optional and, when used, stipulate how a ROHC implementation is to operate.

Configuration parameters are mandatory and must be negotiated between compressor and decompressor, so that they have the same values at both compressor and decompressor, see section 5.1.1.

Implementation parameters make it possible for an external entity to stipulate how an implementation of a ROHC compressor or decompressor should operate. Implementation parameters have local significance, are optional to use and are thus not necessary to negotiate between compressor and decompressor. Note that this does not preclude signaling or negotiating implementation parameters using lower layer functionality in order to set the way a ROHC implementation should operate. Some implementation parameters are valid only at either of compressor or decompressor. Implementation parameters may further be divided into parameters that allow an external entity to describe the way the implementation should operate and parameters that allow an external entity to trigger a specific event, i.e., signals.



## 6.3.1. ROHC implementation parameters at compressor

**CONTEXT\_REINITIALIZATION** -- signal

This parameter triggers a reinitialization of the entire context at the decompressor, both the static and the dynamic part. The compressor **MUST**, when **CONTEXT\_REINITIALIZATION** is triggered, back off to the IR state and fully reinitialize the context by sending IR packets with both the static and dynamic chains covering the entire uncompressed headers until it is reasonably confident that the decompressor contexts are reinitialized. The context reinitialization **MUST** be done for all contexts at the compressor. This parameter may for instance be used to do context relocation at, e.g., a cellular handover that results in a change of compression point in the radio access network.

**NO\_OF\_PACKET\_SIZES\_ALLOWED** -- value: positive integer

This parameter may be set by an external entity to specify the number of packet sizes a ROHC implementation may use. However, the parameter may be used only if **PACKET\_SIZES** is not used by an external entity. With this parameter set, the ROHC implementation at the compressor **MUST NOT** use more different packet sizes than the value this parameter stipulates. The ROHC implementation must itself be able to determine which packet sizes will be used and describe these to an external entity using **PACKET\_SIZES\_USED**. It should be noted that one packet size might be used for several header formats, and that the number of packet sizes can be reduced by employing padding and segmentation.

**NO\_OF\_PACKET\_SIZES\_USED** -- value: positive integer

This parameter is set by the ROHC implementation to indicate how many packet sizes it will actually use. It can be set to a large value to indicate that no particular attempt is made to minimize that number.

**PACKET\_SIZES\_ALLOWED** -- value: list of positive integers (bytes)

This parameter, if set, governs which packet sizes in bytes may be used by the ROHC implementation. Thus, packet sizes not in the set of values for this parameter **MUST NOT** be used. Hence, an external entity can mandate a ROHC implementation to produce packet sizes that fit pre-configured lower layers better. If this parameter is used to stipulate which packet sizes a ROHC implementation can use, the following rules apply:

- A packet large enough to hold the entire IR header (both static and dynamic chain) **MUST** be part of the set of sizes, unless **MRRU** is set to a large enough value to allow segmentation.
- The packet size likely to be used most frequently in the **SO** state **SHOULD** be part of the set.

- The packet size likely to be used most frequently in the FO state SHOULD be part of the set.

PACKET\_SIZES\_USED -- values: set of positive integers (bytes)  
This parameter describes which packet sizes a ROHC implementation uses if NO\_OF\_PACKET\_SIZES\_ALLOWED or PACKET\_SIZES\_ALLOWED is used by an external entity to stipulate how many packet sizes a ROHC implementation should use. The information about used packet sizes (bytes) in this parameter, may then be used to configure lower layers.

PAYLOAD\_SIZES -- values: set of positive integer values (bytes)  
This parameter is set by an external entity that wants to make use of the PACKET\_SIZES\_USED parameter to indicate which payload sizes can be expected.

When a ROHC implementation has a limited set of allowed packet sizes, and the most preferable header format has a size that is not part of the set, it has the following options:

- Choose the next larger header format from the allowed set. This is probably the most efficient choice.
- Use the most preferable header format as if there were no restrictions on size, and then add padding octets to complete a packet of the next larger size in the allowed set.
- Use segmentation to fragment the packet into pieces that would make up packets of sizes that are permissible (possibly after the addition of padding to the last segment).

It should be noted that even if the two last parameters introduce the possibility of restricting the number of packet sizes used, such restrictions will have a negative impact on compression performance.

#### 6.3.2. ROHC implementation parameters at decompressor

MODE -- values: [U-mode, O-mode, R-mode]  
This parameter triggers a mode transition using the mechanism described in chapter 5 when the parameter changes value, i.e., to U-mode (Unidirectional mode), O-mode (Bidirectional Optimistic mode) or R-mode (Bidirectional Reliable mode). The mode transition is made from the current mode to the new mode as signaled by the implementation parameter. For example, if the current mode is Bidirectional Optimistic mode, MODE should have the value O-mode. If the MODE is changed to R-mode, a mode transition MUST be made from Bidirectional Optimistic mode to Bidirectional Reliable mode. MODE should not only serve as a trigger for mode transitions, but also make it visible which mode ROHC operates in.

CLOCK\_RESOLUTION -- value: nonnegative integer

This parameter indicates the system clock resolution in units of milliseconds. A zero (0) value means that there is no clock available. If nonzero, this parameter allows the decompressor to use timer-based TS compression (section 4.5.4) and SN wraparound detection (section 5.3.2.2.4). In this case, its specific value is also significant for correctness of the algorithms.

REVERSE\_DECOMPRESSION\_DEPTH -- value: nonnegative integer

This parameter determines whether reverse decompression as described in section 6.1 should be used or not, and if used, to what extent. The value indicates the maximum number of packets that can be buffered, and thus possibly be reverse decompressed by the decompressor. A zero (0) value means that reverse decompression MUST NOT be used.

#### 6.4. Handling of resource limitations at the decompressor

In a point-to-point link, the two nodes can agree on the number of compressed sessions they are prepared to support for this link. It may, however, not be possible for the decompressor to accurately predict when it will run out of resources. ROHC allows the negotiated number of contexts to be larger than could be accommodated in the worst case. Then, as context resources are consumed, an attempt to set up a new context may be rejected by the decompressor, using the REJECT option of the feedback payload.

Upon reception of a REJECT option, the compressor SHOULD wait for a while before attempting to compress additional streams destined for the rejecting node.

#### 6.5. Implementation structures

This section provides some explanatory material on data structures that a ROHC implementation will have to maintain in one form or another. It is not intended to constrain the implementations.

##### 6.5.1. Compressor context

The compressor context consists of a static part and a dynamic part. The content of the static part is the same as the static chain defined in section 5.7.7. The dynamic part consists of multiple elements which can be categorized into four types.

- a) Sliding Window (SW)
- b) Translation Table (TT)
- c) Flag
- d) Field

These elements may be common to all modes or mode specific. The following table summarizes all these elements.

	Common to all modes	Specific to R-mode	Specific to U/O-mode
SWs	GSW	R_CSW R_IESW	UO_CSW UO_IESW
TTs		R_CTT R_IETT	UO_CTT UO_IETT
Flags	UDP Chksum TSS, TIS RND, RND2 NBO, NBO2		ACKED
Fields	Profile C_MODE C_STATE C_TRANS TS_STRIDE (if TSS = 1) TS_OFFSET (if TSS = 1) TIME_STRIDE (if TIS = 1) CURR_TIME (if TIS = 1) MAX_JITTER_CD (if TIS = 1) LONGEST_LOSS_EVENT(O) CLOCK_RESOLUTION(O) MAX_JITTER(O)		CSRC_REF_ID CSRC_GEN_ID CSRC_GEN_COUNT IPEH_REF_ID IPEH_GEN_ID IPEH_GEN_COUNT

#### 1) GSW: Generic W\_LSB Sliding Window

Each element in GSW consists of all the dynamic fields in the dynamic chain (defined in section 5.7.7) plus the fields specified in a) but excluding the fields specified in b).

- a) Packet Arrival Time (if TIS = 1)  
Scaled RTP Time Stamp (if TSS = 1) (optional)  
Offset\_i (if RND = 0) (optional)

- b) UDP Checksum, TS Stride, CSRC list, IPv6 Extension Headers

#### 2) R\_CSW: CSRC Sliding Window in R-mode

R\_IESW: IPv6 Extension Header Sliding Window in R-mode

UO\_CSW: CSRC Sliding Window in U/O-mode

UO\_IESW: IPv6 Extension Header Sliding Window in U/O-mode

Each element in R\_CSW, R\_IESW, UO\_CSW and UO\_IESW is defined in section 6.5.3.

3) R\_CTT: CSRC Translation Table in R-mode

R\_IETT: IPv6 Extension Header Translation Table in U/O-mode

UO\_CTT: CSRC Translation Table in U/O-mode

UO\_IETT: IPv6 Extension Header Translation Table in U/O-mode

Each element in R\_CTT and R\_IETT is defined in section 5.8.1.1.

Each element in UO\_CTT and UO\_IETT is defined in section 5.8.1.2.

4) ACKED: Indicates whether or not the decompressor has ever acked

5) CURR\_TIME: The current time value (used for context relocation when timer-based timestamp compression is used)

6) All the other flags and fields are defined elsewhere in the ROHC document.

#### 6.5.2. Decompressor context

The decompressor context consists of a static part and a dynamic part. The content of the static part is the same as the static chain defined in section 5.7.7. The dynamic part consists of multiple elements, one of which is the nonstatic reference header that includes all the nonstatic fields. These nonstatic fields are the fields in the dynamic chain defined in section 5.7.7, excluding UDP Checksum and TS\_Stride. All the remaining elements can be categorized into four types:

- a) Sliding Window (SW)
- b) Translation Table (TT)
- d) Flag
- e) Field

These elements may be mode specific or common to all modes. The following table summarizes all these elements.

	Common to all modes	Specific to R-mode	Specific to U/O-mode
SWs		R_CSW R_IESW	UO_CSW UO_IESW
TTs		R_CTT R_IETT	UO_CTT UO_IETT
Flags	UDP Checksum TSS, TIS RND, RND2 NBO, NBO2		ACKED
Fields	Profile D_MODE D_STATE D_TRANS TS_STRIDE (if TSS = 1) TS_OFFSET (if TSS = 1) TIME_STRIDE (if TIS = 1) PKT_ARR_TIME (if TIS = 1) LONGEST_LOSS_EVENT(O) CLOCK_RESOLUTION(O) MAX_JITTER(O)		CSRC_GEN_ID IPEH_GEN_ID PRE_SN_V_REF

1) ACKED: Indicates whether or not ACK has ever been sent.

2) PKT\_ARR\_TIME: The arrival time of the packet that most recently decompressed and verified using CRC.

PRE\_SN\_V\_REF: The sequence number of the packet verified before the most recently verified packet.

CSRC\_GEN\_ID: The CSRC gen\_id of the most recently received packet.

IPEH\_GEN\_ID: The IPv6 Extension Header gen\_id of the most recently received packet.

3) The remaining elements are as defined in the compressor context.

### 6.5.3. List compression: Sliding windows in R-mode and U/O-mode

In R-mode list compression (see section 5.8.2.1), each entry in the sliding window, both at the compressor side and at the decompressor side, has the following structure:

```
+-----+-----+-----+
| RTP Sequence Number | icount | index list |
+-----+-----+-----+
```

The table index list contains a list of index. Each of these index corresponds to the item in the original list carried in the packet identified by the RTP Sequence Number. The mapping between the index and the item is identified in the translation table. The icount field carries the number of index in the following index list.

In U/O-mode list compression, each entry in the sliding window at both the compressor side and decompressor side has the following structure.

```
+-----+-----+-----+
| Gen_id | icount | index list |
+-----+-----+-----+
```

The icount and index list fields are the same as defined in R-mode. Instead of using the RTP Sequence Number to identify each entry, the Gen\_id is included in the sliding window in U/O-mode.

## 7. Security Considerations

Because encryption eliminates the redundancy that header compression schemes try to exploit, there is some inducement to forego encryption of headers in order to enable operation over low-bandwidth links. However, for those cases where encryption of data (and not headers) is sufficient, RTP does specify an alternative encryption method in which only the RTP payload is encrypted and the headers are left in the clear. That would still allow header compression to be applied.

ROHC compression is transparent with regard to the RTP Sequence Number and RTP Timestamp fields, so the values of those fields can be used as the basis of payload encryption schemes (e.g., for computation of an initialization vector).

A malfunctioning or malicious header compressor could cause the header decompressor to reconstitute packets that do not match the original packets but still have valid IP, UDP and RTP headers and possibly also valid UDP checksums. Such corruption may be detected with end-to-end authentication and integrity mechanisms which will not be affected by the compression. Moreover, this header compression scheme uses an internal checksum for verification of reconstructed headers. This reduces the probability of producing decompressed headers not matching the original ones without this being noticed.

Denial-of-service attacks are possible if an intruder can introduce (for example) bogus STATIC, DYNAMIC or FEEDBACK packets onto the link and thereby cause compression efficiency to be reduced. However, an intruder having the ability to inject arbitrary packets at the link layer in this manner raises additional security issues that dwarf those related to the use of header compression.

## 8. IANA Considerations

The ROHC profile identifier is a non-negative integer. In many negotiation protocols, it will be represented as a 16-bit value. Due to the way the profile identifier is abbreviated in ROHC packets, the 8 least significant bits of the profile identifier have a special significance: Two profile identifiers with identical 8 LSBs should be assigned only if the higher-numbered one is intended to supersede the lower-numbered one. To highlight this relationship, profile identifiers should be given in hexadecimal (as in 0x1234, which would for example supersede 0x0A34).

Following the policies outlined in [IANA-CONSIDERATIONS], the IANA policy for assigning new values for the profile identifier shall be Specification Required: values and their meanings must be documented in an RFC or in some other permanent and readily available reference, in sufficient detail that interoperability between independent implementations is possible. In the 8 LSBs, the range 0 to 127 is reserved for IETF standard-track specifications; the range 128 to 254 is available for other specifications that meet this requirement (such as Informational RFCs). The LSB value 255 is reserved for future extensibility of the present specification.

The following profile identifiers are already allocated:

Profile identifier	Document	Usage
0x0000	RFCthis	ROHC uncompressed
0x0001	RFCthis	ROHC RTP
0x0002	RFCthis	ROHC UDP
0x0003	RFCthis	ROHC ESP



## 9. Acknowledgments

Earlier header compression schemes described in [CJHC], [IPHC], and [CRTP] have been important sources of ideas and knowledge.

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## 10. Intellectual Property Right Claim Considerations

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## Appendix A. Detailed classification of header fields

Header compression is possible thanks to the fact that most header fields do not vary randomly from packet to packet. Many of the fields exhibit static behavior or change in a more or less predictable way. When designing a header compression scheme, it is of fundamental importance to understand the behavior of the fields in detail.

In this appendix, all IP, UDP and RTP header fields are classified and analyzed in two steps. First, we have a general classification in A.1 where the fields are classified on the basis of stable knowledge and assumptions. The general classification does not take into account the change characteristics of changing fields because those will vary more or less depending on the implementation and on the application used. A less stable but more detailed analysis of the change characteristics is then done in A.2. Finally, A.3 summarizes this appendix with conclusions about how the various header fields should be handled by the header compression scheme to optimize compression and functionality.



## A.1. General classification

At a general level, the header fields are separated into 5 classes:

INFERRED	These fields contain values that can be inferred from other values, for example the size of the frame carrying the packet, and thus do not have to be handled at all by the compression scheme.
STATIC	These fields are expected to be constant throughout the lifetime of the packet stream. Static information must in some way be communicated once.
STATIC-DEF	STATIC fields whose values define a packet stream. They are in general handled as STATIC.
STATIC-KNOWN	These STATIC fields are expected to have well-known values and therefore do not need to be communicated at all.
CHANGING	These fields are expected to vary in some way: randomly, within a limited value set or range, or in some other manner.

In this section, each of the IP, UDP and RTP header fields is assigned to one of these classes. For all fields except those classified as CHANGING, the motives for the classification are also stated. In section A.2, CHANGING fields are further examined and classified on the basis of their expected change behavior.

## A.1.1. IPv6 header fields

Field	Size (bits)	Class
Version	4	STATIC
Traffic Class	8	CHANGING
Flow Label	20	STATIC-DEF
Payload Length	16	INFERRED
Next Header	8	STATIC
Hop Limit	8	CHANGING
Source Address	128	STATIC-DEF
Destination Address	128	STATIC-DEF

Version

The version field states which IP version is used. Packets with different values in this field must be handled by different IP stacks. All packets of a packet stream must therefore be of the same IP version. Accordingly, the field is classified as STATIC.

Flow Label

This field may be used to identify packets belonging to a specific packet stream. If not used, the value should be set to zero. Otherwise, all packets belonging to the same stream must have the same value in this field, it being one of the fields that define the stream. The field is therefore classified as STATIC-DEF.

Payload Length

Information about packet length (and, consequently, payload length) is expected to be provided by the link layer. The field is therefore classified as INFERRED.

Next Header

This field will usually have the same value in all packets of a packet stream. It encodes the type of the subsequent header. Only when extension headers are sometimes present and sometimes not, will the field change its value during the lifetime of the stream. The field is therefore classified as STATIC.

Source and Destination addresses

These fields are part of the definition of a stream and must thus be constant for all packets in the stream. The fields are therefore classified as STATIC-DEF.

Total size of the fields in each class:

Class	Size (octets)
INFERRED	2
STATIC	1.5
STATIC-DEF	34.5
CHANGING	2

## A.1.2. IPv4 header fields

Field	Size (bits)	Class
Version	4	STATIC
Header Length	4	STATIC-KNOWN
Type Of Service	8	CHANGING
Packet Length	16	INFERRED
Identification	16	CHANGING
Reserved flag	1	STATIC-KNOWN
Don't Fragment flag	1	STATIC
More Fragments flag	1	STATIC-KNOWN
Fragment Offset	13	STATIC-KNOWN
Time To Live	8	CHANGING
Protocol	8	STATIC
Header Checksum	16	INFERRED
Source Address	32	STATIC-DEF
Destination Address	32	STATIC-DEF

## Version

The version field states which IP version is used. Packets with different values in this field must be handled by different IP stacks. All packets of a packet stream must therefore be of the same IP version. Accordingly, the field is classified as STATIC.

## Header Length

As long no options are present in the IP header, the header length is constant and well known. If there are options, the fields would be STATIC, but it is assumed here that there are no options. The field is therefore classified as STATIC-KNOWN.

## Packet Length

Information about packet length is expected to be provided by the link layer. The field is therefore classified as INFERRED.

## Flags

The Reserved flag must be set to zero and is therefore classified as STATIC-KNOWN. The Don't Fragment (DF) flag will be constant for all packets in a stream and is therefore classified as STATIC.

Finally, the More Fragments (MF) flag is expected to be zero because fragmentation is NOT expected, due to the small packet size expected. The More Fragments flag is therefore classified as STATIC-KNOWN.

#### Fragment Offset

Under the assumption that no fragmentation occurs, the fragment offset is always zero. The field is therefore classified as STATIC-KNOWN.

#### Protocol

This field will usually have the same value in all packets of a packet stream. It encodes the type of the subsequent header. Only when extension headers are sometimes present and sometimes not, will the field change its value during the lifetime of a stream. The field is therefore classified as STATIC.

#### Header Checksum

The header checksum protects individual hops from processing a corrupted header. When almost all IP header information is compressed away, there is no point in having this additional checksum; instead it can be regenerated at the decompressor side. The field is therefore classified as INFERRED.

#### Source and Destination addresses

These fields are part of the definition of a stream and must thus be constant for all packets in the stream. The fields are therefore classified as STATIC-DEF.

#### Total size of the fields in each class:

Class	Size (octets)
INFERRED	4
STATIC	1 oct + 5 bits
STATIC-DEF	8
STATIC-KNOWN	2 oct + 3 bits
CHANGING	4

## A.1.3. UDP header fields

Field	Size (bits)	Class
Source Port	16	STATIC-DEF
Destination Port	16	STATIC-DEF
Length	16	INFERRED
Checksum	16	CHANGING

## Source and Destination ports

These fields are part of the definition of a stream and must thus be constant for all packets in the stream. The fields are therefore classified as STATIC-DEF.

## Length

This field is redundant and is therefore classified as INFERRED.

Total size of the fields in each class:

Class	Size (octets)
INFERRED	2
STATIC-DEF	4
CHANGING	2

## A.1.4. RTP header fields

Field	Size (bits)	Class
Version	2	STATIC-KNOWN
Padding	1	STATIC
Extension	1	STATIC
CSRC Counter	4	CHANGING
Marker	1	CHANGING
Payload Type	7	CHANGING
Sequence Number	16	CHANGING
Timestamp	32	CHANGING
SSRC	32	STATIC-DEF
CSRC	0(-480)	CHANGING

## Version

Only one working RTP version exists, namely version 2. The field is therefore classified as STATIC-KNOWN.

## Padding

The use of this field is application-dependent, but when payload padding is used it is likely to be present in all packets. The field is therefore classified as STATIC.

## Extension

If RTP extensions are used by the application, these extensions are likely to be present in all packets (but the use of extensions is very uncommon). However, for safety's sake this field is classified as STATIC and not STATIC-KNOWN.

## SSRC

This field is part of the definition of a stream and must thus be constant for all packets in the stream. The field is therefore classified as STATIC-DEF.

Total size of the fields in each class:

Class	Size (octets)
STATIC	2 bits
STATIC-DEF	4
STATIC-KNOWN	2 bits
CHANGING	7.5(-67.5)

## A.1.5. Summary for IP/UDP/RTP

Summarizing this for IP/UDP/RTP one obtains

Class \ IP ver	IPv6 (octets)	IPv4 (octets)
INFERRED	4	6
STATIC	1 oct + 6 bits	1 oct + 7 bits
STATIC-DEF	42.5	16
STATIC-KNOWN	2 bits	2 oct + 5 bits
CHANGING	11.5(-71.5)	13.5(-73.5)
Total	60(-120)	40(-100)

## A.2. Analysis of change patterns of header fields

To design suitable mechanisms for efficient compression of all header fields, their change patterns must be analyzed. For this reason, an extended classification is done based on the general classification in A.1, considering the fields which were labeled CHANGING in that classification. Different applications will use the fields in different ways, which may affect their behavior. For the fields whose behavior is variable, typical behavior for conversational audio and video will be discussed.

The CHANGING fields are separated into five different subclasses:

STATIC	These are fields that were classified as CHANGING on a general basis, but are classified as STATIC here due to certain additional assumptions.
SEMISTATIC	These fields are STATIC most of the time. However, occasionally the value changes but reverts to its original value after a known number of packets.
RARELY-CHANGING (RC)	These are fields that change their values occasionally and then keep their new values.
ALTERNATING	These fields alternate between a small number of different values.
IRREGULAR	These, finally, are the fields for which no useful change pattern can be identified.

To further expand the classification possibilities without increasing complexity, the classification can be done either according to the values of the field and/or according to the values of the deltas for the field.

When the classification is done, other details are also stated regarding possible additional knowledge about the field values and/or field deltas, according to the classification. For fields classified as STATIC or SEMISTATIC, the case could be that the value of the field is not only STATIC but also well KNOWN a priori (two states for SEMISTATIC fields). For fields with non-irregular change behavior, it could be known that changes usually are within a LIMITED range compared to the maximal change for the field. For other fields, the values are completely UNKNOWN.



Table A.1 classifies all the CHANGING fields on the basis of their expected change patterns, especially for conversational audio and video.

Field	Value/Delta	Class	Knowledge	
	Sequential	Delta	STATIC	KNOWN
IPv4 Id:	Seq. jump	Delta	RC	LIMITED
	Random	Value	IRREGULAR	UNKNOWN
IP TOS / Tr. Class	Value	RC	UNKNOWN	
IP TTL / Hop Limit	Value	ALTERNATING	LIMITED	
UDP Checksum:	Disabled	Value	STATIC	KNOWN
	Enabled	Value	IRREGULAR	UNKNOWN
RTP CSRC Count:	No mix	Value	STATIC	KNOWN
	Mixed	Value	RC	LIMITED
RTP Marker	Value	SEMISTATIC	KNOWN/KNOWN	
RTP Payload Type	Value	RC	UNKNOWN	
RTP Sequence Number	Delta	STATIC	KNOWN	
RTP Timestamp	Delta	RC	LIMITED	
RTP CSRC List:	No mix	-	-	-
	Mixed	Value	RC	UNKNOWN

Table A.1 : Classification of CHANGING header fields

The following subsections discuss the various header fields in detail. Note that table A.1 and the discussions below do not consider changes caused by loss or reordering before the compression point.

### A.2.1. IPv4 Identification

The Identification field (IP ID) of the IPv4 header is there to identify which fragments constitute a datagram when reassembling fragmented datagrams. The IPv4 specification does not specify exactly how this field is to be assigned values, only that each packet should get an IP ID that is unique for the source-destination pair and protocol for the time the datagram (or any of its fragments) could be alive in the network. This means that assignment of IP ID values can be done in various ways, which we have separated into three classes.

#### Sequential jump

This is the most common assignment policy in today's IP stacks. A single IP ID counter is used for all packet streams. When the sender is running more than one packet stream simultaneously, the IP ID can increase by more than one between packets in a stream. The IP ID values will be much more predictable and require less bits to transfer than random values, and the packet-to-packet increment (determined by the number of active outgoing packet streams and sending frequencies) will usually be limited.

#### Random

Some IP stacks assign IP ID values using a pseudo-random number generator. There is thus no correlation between the ID values of subsequent datagrams. Therefore there is no way to predict the IP ID value for the next datagram. For header compression purposes, this means that the IP ID field needs to be sent uncompressed with each datagram, resulting in two extra octets of header. IP stacks in cellular terminals SHOULD NOT use this IP ID assignment policy.

#### Sequential

This assignment policy keeps a separate counter for each outgoing packet stream and thus the IP ID value will increment by one for each packet in the stream, except at wrap around. Therefore, the delta value of the field is constant and well known a priori. When RTP is used on top of UDP and IP, the IP ID value follows the RTP Sequence Number. This assignment policy is the most desirable for header compression purposes. However, its usage is not as common as it perhaps should be. The reason may be that it can be realized only when UDP and IP are implemented together so that UDP, which separates packet streams by the Port identification fields, can make IP use separate ID counters for each packet stream.

In order to avoid violating [IPv4], packets sharing the same IP address pair and IP protocol number cannot use the same IP ID values. Therefore, implementations of sequential policies must make the ID number spaces disjoint for packet streams of the same IP protocol going between the same pair of nodes. This can be done in a number of ways, all of which introduce occasional jumps, and thus makes the policy less than perfectly sequential. For header compression purposes less frequent jumps are preferred.

It should be noted that the ID is an IPv4 mechanism and is therefore not a problem for IPv6. For IPv4 the ID could be handled in three different ways. First, we have the inefficient but reliable solution where the ID field is sent as-is in all packets, increasing the compressed headers by two octets. This is the best way to handle the ID field if the sender uses random assignment of the ID field. Second, there can be solutions with more flexible mechanisms requiring less bits for the ID handling as long as sequential jump assignment is used. Such solutions will probably require even more bits if random assignment is used by the sender. Knowledge about the sender's assignment policy could therefore be useful when choosing between the two solutions above. Finally, even for IPv4, header compression could be designed without any additional information for the ID field included in compressed headers. To use such schemes, it must be known which assignment policy for the ID field is being used by the sender. That might not be possible to know, which implies that the applicability of such solutions is very uncertain. However, designers of IPv4 stacks for cellular terminals SHOULD use an assignment policy close to sequential.

#### A.2.2. IP Traffic-Class / Type-Of-Service

The Traffic-Class (IPv6) or Type-Of-Service (IPv4) field is expected to be constant during the lifetime of a packet stream or to change relatively seldom.

#### A.2.3. IP Hop-Limit / Time-To-Live

The Hop-Limit (IPv6) or Time-To-Live (IPv4) field is expected to be constant during the lifetime of a packet stream or to alternate between a limited number of values due to route changes.

#### A.2.4. UDP Checksum

The UDP checksum is optional. If disabled, its value is constantly zero and could be compressed away. If enabled, its value depends on the payload, which for compression purposes is equivalent to it changing randomly with every packet.

#### A.2.5. RTP CSRC Counter

This is a counter indicating the number of CSRC items present in the CSRC list. This number is expected to be almost constant on a packet- to-packet basis and change by small amounts. As long as no RTP mixer is used, the value of this field is zero.

#### A.2.6. RTP Marker

For audio the marker bit should be set only in the first packet of a talkspurt, while for video it should be set in the last packet of every picture. This means that in both cases the RTP marker is classified as SEMISTATIC with well-known values for both states.

#### A.2.7. RTP Payload Type

Changes of the RTP payload type within a packet stream are expected to be rare. Applications could adapt to congestion by changing payload type and/or frame sizes, but that is not expected to happen frequently.

#### A.2.8. RTP Sequence Number

The RTP Sequence Number will be incremented by one for each packet sent.

#### A.2.9. RTP Timestamp

In the audio case:

As long as there are no pauses in the audio stream, the RTP Timestamp will be incremented by a constant delta, corresponding to the number of samples in the speech frame. It will thus mostly follow the RTP Sequence Number. When there has been a silent period and a new talkspurt begins, the timestamp will jump in proportion to the length of the silent period. However, the increment will probably be within a relatively limited range.

In the video case:

Between two consecutive packets, the timestamp will either be unchanged or increase by a multiple of a fixed value corresponding to the picture clock frequency. The timestamp can also decrease by a multiple of the fixed value if B-pictures are used. The delta interval, expressed as a multiple of the picture clock frequency, is in most cases very limited.

#### A.2.10. RTP Contributing Sources (CSRC)

The participants in a session, which are identified by the CSRC fields, are expected to be almost the same on a packet-to-packet basis with relatively few additions and removals. As long as RTP mixers are not used, no CSRC fields are present at all.

#### A.3. Header compression strategies

This section elaborates on what has been done in previous sections. On the basis of the classifications, recommendations are given on how to handle the various fields in the header compression process. Seven different actions are possible; these are listed together with the fields to which each action applies.

##### A.3.1. Do not send at all

The fields that have well known values a priori do not have to be sent at all. These are:

- IPv6 Payload Length
- IPv4 Header Length
- IPv4 Reserved Flag
- IPv4 Last Fragment Flag
- IPv4 Fragment Offset
  
- UDP Checksum (if disabled)
- RTP Version

##### A.3.2. Transmit only initially

The fields that are constant throughout the lifetime of the packet stream have to be transmitted and correctly delivered to the decompressor only once. These are:

- IP Version
- IP Source Address
- IP Destination Address
- IPv6 Flow Label
- IPv4 May Fragment Flag
- UDP Source Port
- UDP Destination Port
- RTP Padding Flag
- RTP Extension Flag
- RTP SSRC

#### A.3.3. Transmit initially, but be prepared to update

The fields that are changing only occasionally must be transmitted initially but there must also be a way to update these fields with new values if they change. These fields are:

- IPv6 Next Header
- IPv6 Traffic Class
- IPv6 Hop Limit
- IPv4 Protocol
- IPv4 Type Of Service (TOS)
- IPv4 Time To Live (TTL)
- RTP CSRC Counter
- RTP Payload Type
- RTP CSRC List

Since the values of the IPv4 Protocol and the IPv6 Next Header fields are in effect linked to the type of the subsequent header, they deserve special treatment when subheaders are inserted or removed.

#### A.3.4. Be prepared to update or send as-is frequently

For fields that normally either are constant or have values deducible from some other field, but that frequently diverge from that behavior, there must be an efficient way to update the field value or send it as-is in some packets. These fields are:

- IPv4 Identification (if not sequentially assigned)
- RTP Marker
- RTP Timestamp

#### A.3.5. Guarantee continuous robustness

For fields that behave like a counter with a fixed delta for ALL packets, the only requirement on the transmission encoding is that packet losses between compressor and decompressor must be tolerable. If several such fields exist, all these can be communicated together. Such fields can also be used to interpret the values for fields listed in the previous section. Fields that have this counter behavior are:

- IPv4 Identification (if sequentially assigned)
- RTP Sequence Number

#### A.3.6. Transmit as-is in all packets

Fields that have completely random values for each packet must be included as-is in all compressed headers. Those fields are:

- IPv4 Identification (if randomly assigned)
- UDP Checksum (if enabled)

#### A.3.7. Establish and be prepared to update delta

Finally, there is a field that is usually increasing by a fixed delta and is correlated to another field. For this field it would make sense to make that delta part of the context state. The delta must then be initiated and updated in the same way as the fields listed in A.3.3. The field to which this applies is:

- RTP Timestamp

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